



CANHEIT 2011
The Nature of Technology

DNS Security

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<O>

**BUYER
BEWARE**

Caveat Emptor



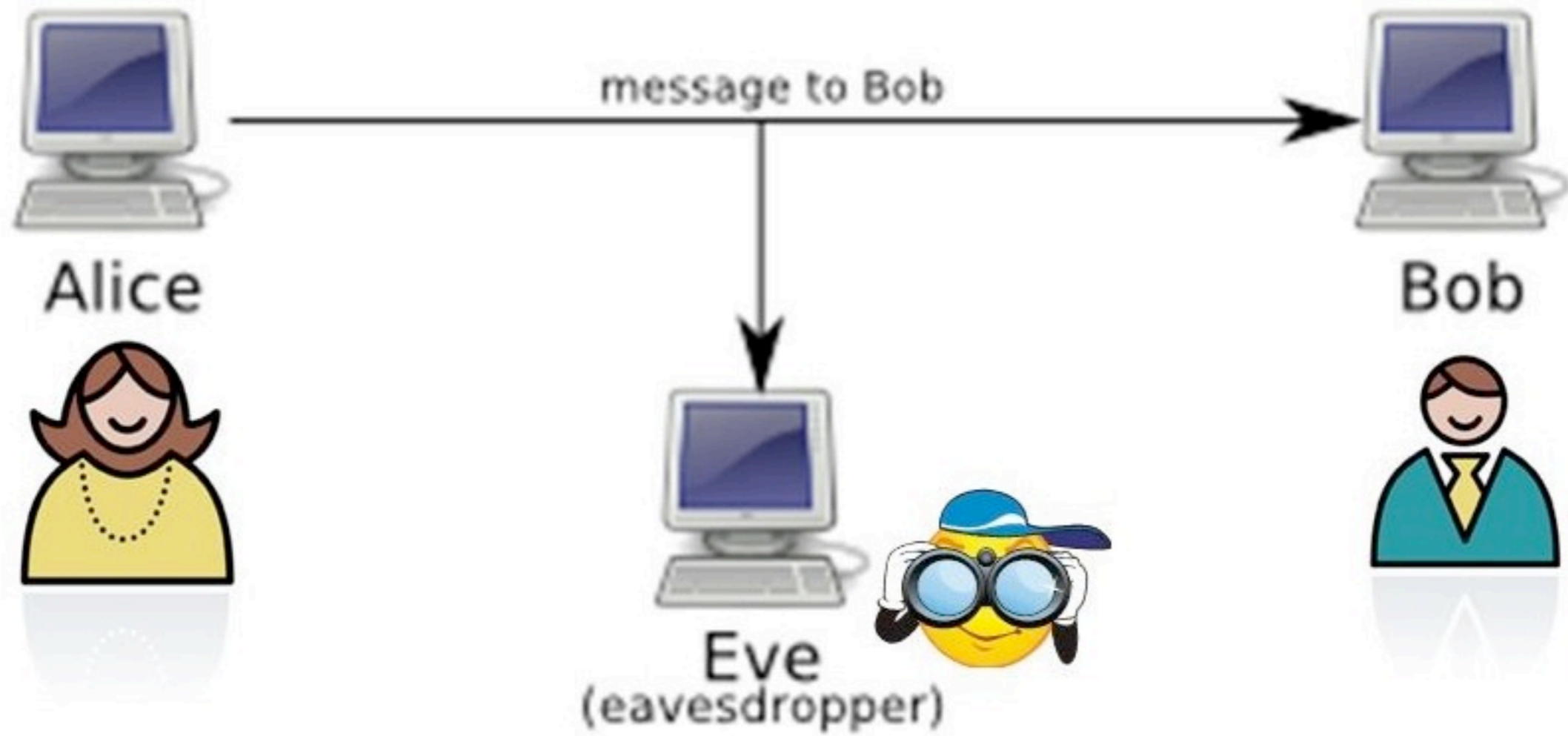
Computer security
usually means...



Communication
between two parties



But it's complicated



Eve **is** the Wo(Man) in the Middle

3 goals of good computer security



Confidentiality...



despite espionage...



e.g. Eve wants to steal data



But to Eve it's all gibberish



)@*#)@*#*\$@)



Integrity...



despite corruption...



e.g. Eve wants to change data



Alice and Bob are NOT
getting the wrong data



Forged data is detected



e.g. Not knowing that data has been corrupted is a violation of integrity



Availability...



despite sabotage...



e.g. Eve wants to destroy data



Bob and Alice are
getting the right data



e.g. DoS attack is a
violation of availability



e.g. Blocking data is a violation of availability



Computer security
usually means...

Cryptography

Cryptography usually
means...

Mathematics

Mathematics usually
means...

Rigorous proofs

Rigorous proof
means ...

The masses* trust the
few elite who “know”

* the intelligent, educated ones are included here

၉.၄.

Pierre de Fermat



Famous “Last” Theorem

Andrew Wiles



1637

$$a^n + b^n \neq c^n$$

$$\{n \geq 3, a, b, c > 0\} \quad a, b, c, n \in \mathbb{Z}$$

1994/95



It takes a very great deal of mathematical knowledge to understand Wile's proof; the original paper is **hundreds** of pages and when I tried to read it I couldn't get past page **one**; this highlights the fact that proof really has to do with authority and trust, not whether a computer can verify the mathematician's steps

Allan Reeve Wilks Ph.D., Statistician, AT&T Bell Labs
allan@research.att.com

Therefore ...

it's good to remember

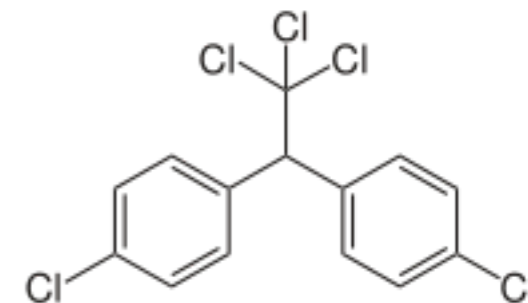
when it comes to security
and cryptography in
particular ...

There are professionals*,
stuffed shirts, plumbers
and actors

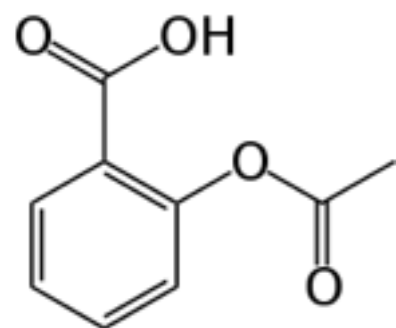
* Even the professionals
do not always agree.

* This is certainly the case
in DNS security.

<1>



Nomenclature



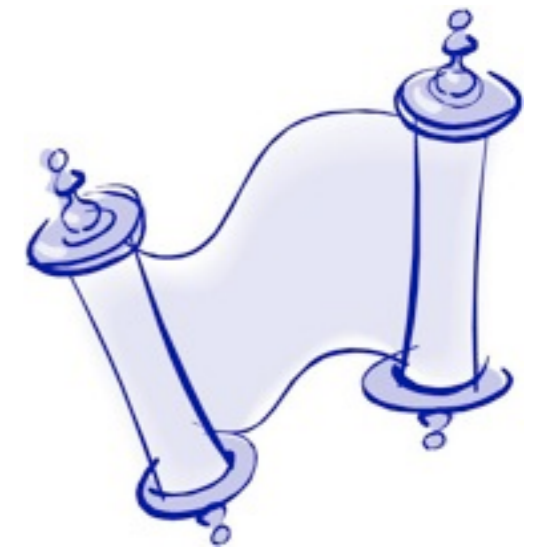
DNS == Domain Name System

Stub Resolver



Caching Name Server





Authoritative Name Server





Function of the DNS

$$f(x) = \sqrt{a^2 - 4ab + b^2}$$

Name/Number Lookup Service



Map/Translate Names to Addresses*

* there are other important data /resource records

E.g.

www.mcgill.ca

www.mcgill.ca
132.216.177.140

A record

worldbank.int

worldbank.int

dns1.worldbank.org

dns2.worldbank.org

dns3.worldbank.org

dns4.worldbank.org

NS record

president@whitehouse.gov

president@whitehouse.gov

mail1.eop.gov.

mail2.eop.gov.

mail3.eop.gov.

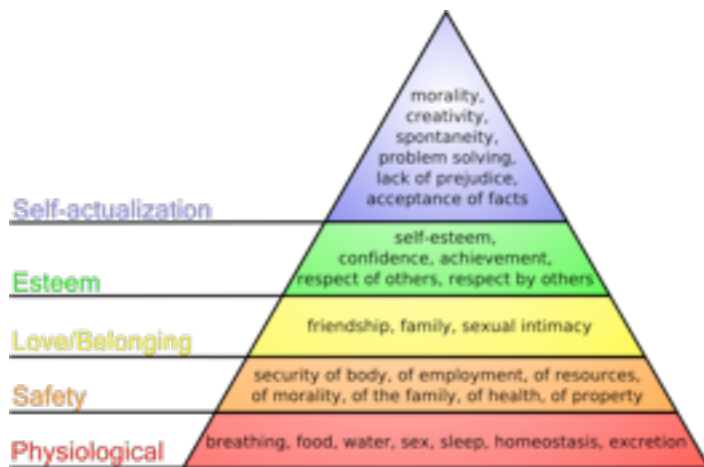
mail4.eop.gov.

MX record

<3>

DNS Implementation

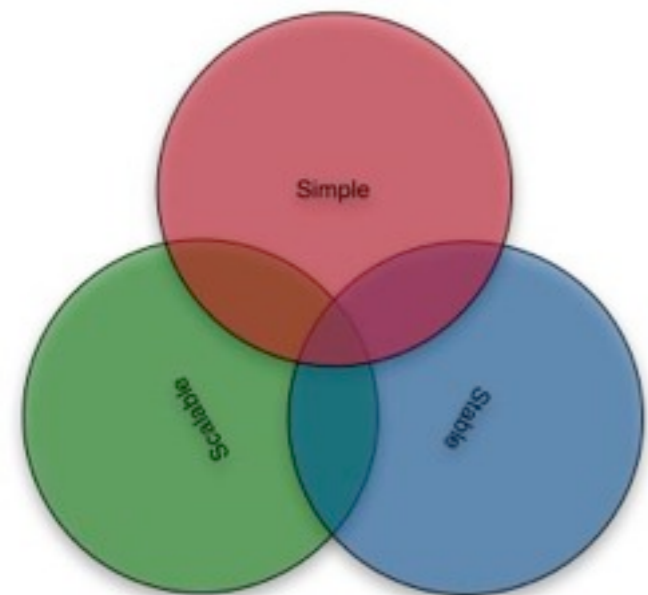
Hierarchical



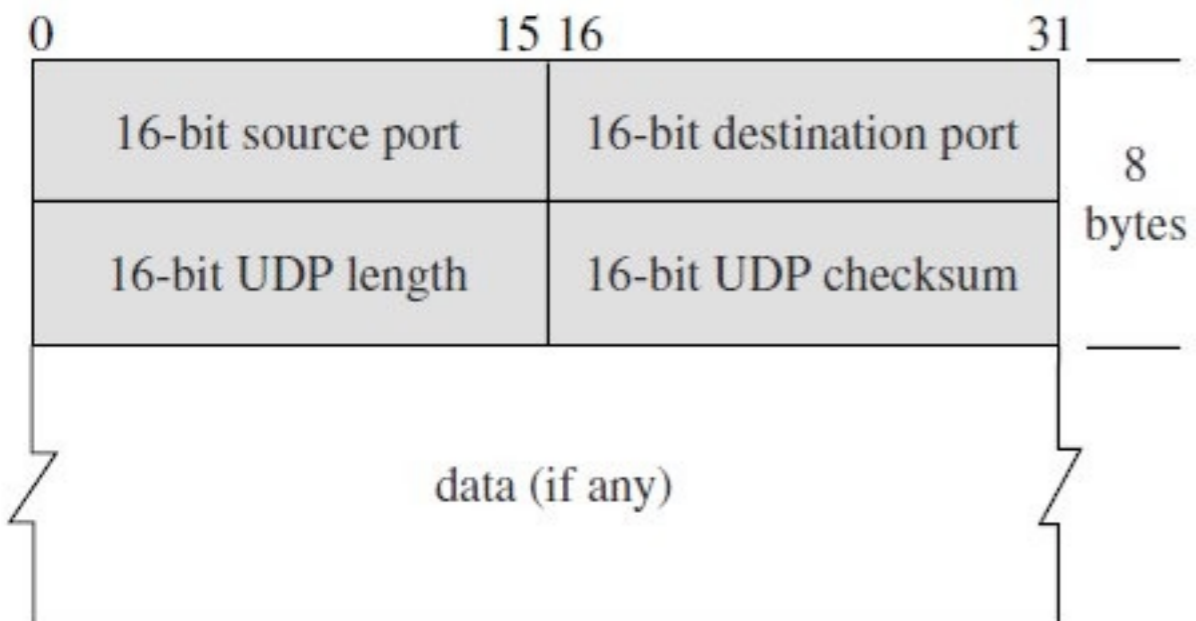
Globally Distributed



Scalable



Standard IP Protocol



Database



Query



Response

<4>

DNS Centrality & Criticality

DNS is a central anchor point of trust for the entire Internet's infrastructure

DNS is a **central anchor point of trust** for the entire Internet's infrastructure



Almost every user interaction deals
with names.

Almost every user interaction deals
with **names**.

president@utoronto.ca

<https://www.royalbank.com/>

Almost every Internet protocol interaction deals with numbers and addresses.

128.100.103.10

2001:beef:0666::2

2007: 24 Billion DNS Queries A Day



2008: 48 Billion DNS Queries A Day



Q: How secure is the Internet?



Q: Can Internet Mail be stolen?



A: Yes. And it's not too hard.



Mail client uses DNS, gets the wrong address for the remote mail server



Mail client uses DNS, attackers see
and change the packets en-route



Q: Can Web Pages be Forged?



A: Yes. And it's not too hard.

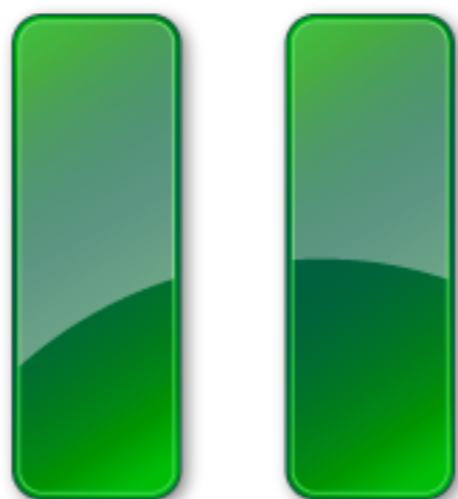


Browser sends DNS request, gets the wrong address and makes a HTTP connection to the bad guys server



Browser sends DNS request, makes a HTTP connection, attackers see and change the packets.





Q: How do we protect the DNS?



Q: Does cryptography solve the problem?



A: In theory... Yes!



A: In practice... ???



Q: Am I using cryptography?



Q: Are you using cryptography?



A: Sometimes yes; Normally no.



Q: Why is this so?



A: Most Internet Protocols do not support cryptography



Q: Why is this so?



A: Integration of cryptography is hard for protocol designers.



N.B. Some popular IP protocols do have
cryptographic options!



e.g. HTTPS (RFC 2818)



Q: Why do some implementations of these protocols do not support cryptography?



A: It's hard for software authors to implement the cryptography. Non cryptographic options are much easier.

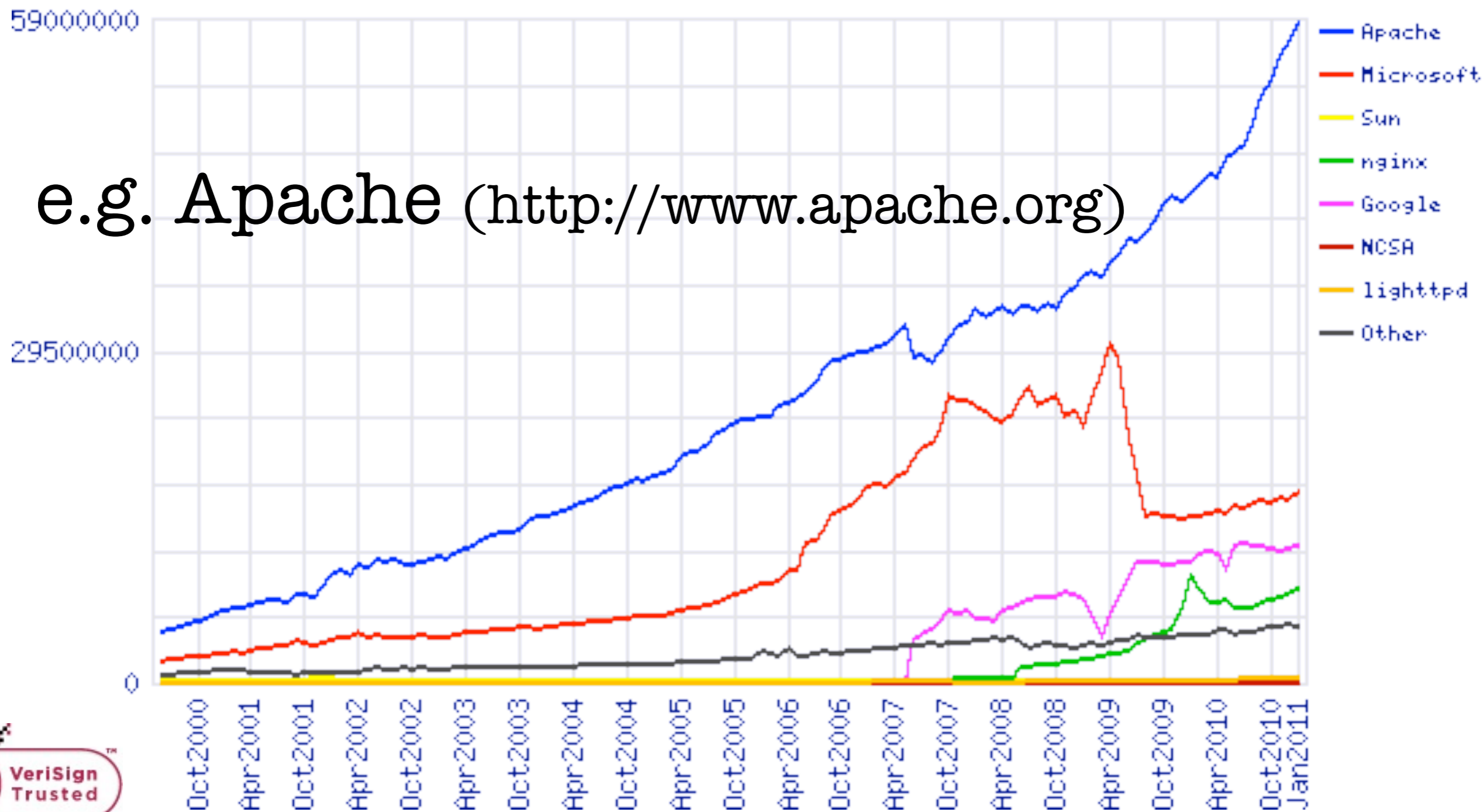


N.B.

Some popular implementations do support cryptography!



e.g. Apache (<http://www.apache.org>)



A: It's harder and more costly for site administrators to turn it on and to keep it on.





Q: How secure are SSL certificates?



A: You tell me ...

Domain Access Verification

ra@godaddy.com <ra@godaddy.com>

To: Russell Sutherland <russell.sutherland@utoronto.ca>

Dear Secure Certificate Customer,

We have received a Certificate Signing Request for the following domain: theta.utoronto.ca.

Our query of the Whois database returned your name as the administrator for the domain in the certificate request.

In order to verify the validity of this request and that it was submitted by the entity to which the domain in the request is registered, please signify your final approval or disapproval of the certificate request by clicking the link below.

https://certs.godaddy.com/anonymous/domainapproval.seam?vk=3126963_dbb95e2dcbe3

Approval of the request will enable us to continue processing your request. Failure to approve the certificate request will lead to denial of the request.

If the above address does not appear as a clickable link, cut/copy and paste it into your browser's address bar.

If the Verification Page requests it, please use the following Verification Key: **3126963_dbb95e2dcbe3**

This part of our authentication process serves to ensure that only the entity/individual that controls the domain in the request can obtain a certificate for that domain.

If you encounter any problems or have any questions, our Customer Support department is ready to help, around-the-clock, seven days a week.

Customer Support:

E-Mail: ra@godaddy.com

Phone: 480.505.8852

Fax: 480.393.5009

For further information, log in to your account at <https://certs.godaddy.com>.





Q: How many certificate authorities does your browser trust?



A: ~ 1400 if your browser is Firefox or IE



N.B. Some important installations do support cryptography!



e.g. SourceForge has an SSL certificate
and has set up SSL servers:
<https://sourceforge.net/account/>



N.B. Cryptography is not enabled
everywhere on the site!



e.g.

<https://sourceforge.net/community> gets
redirected to

<http://sourceforge.net/community>



Q: Why does SourceForge turn off SSL/
cryptographic protection?

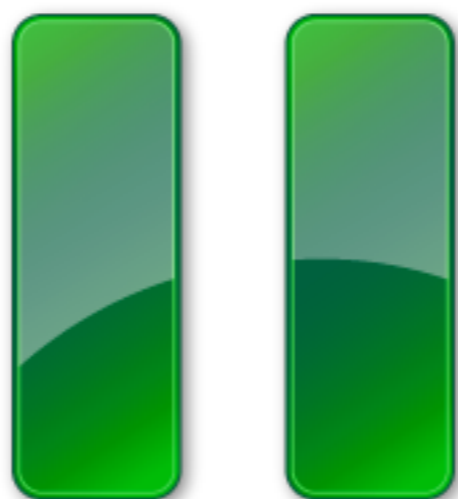


A: Enabling SSL for all transactions is costly in terms of CPU cycles/load.



A: SSL-acceleration is available but again costly (\$\$\$).





Q: Why are cryptographic operations so expensive?



Q: Can cryptographic operations be made faster and still be correct?



Q: Can cryptographic operations be made fast enough to handle all of a www site's operation?



Q: Can cryptographic operations be made fast enough to handle all Internet transactions?



Q: Can Internet cryptography be easy to implement and manage?



Q: Can Internet cryptography be done in software?



Q: Can Internet cryptography be easy to add to Internet protocols?





Q: Will governments be afraid of universal Internet cryptography ?



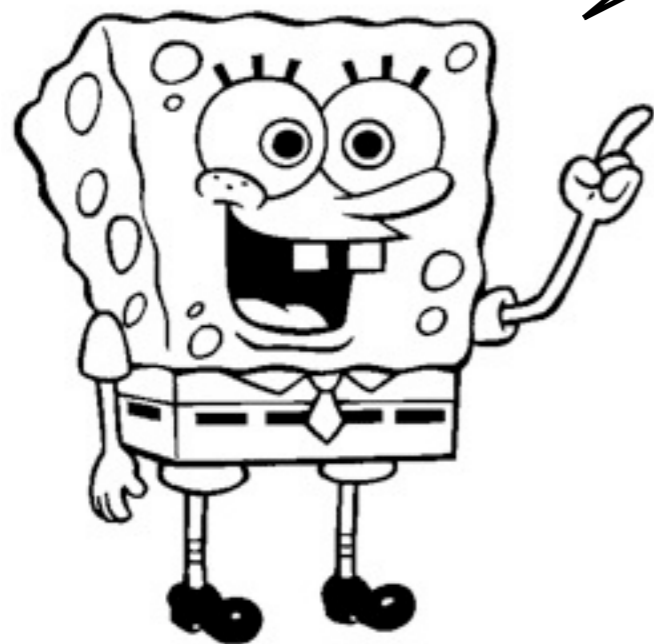


Q: Given gangsters and bad guys get to use cryptography can the average user have access to the same stuff?

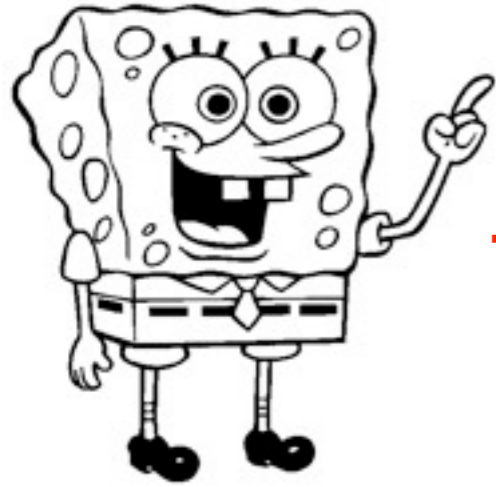


<5>

A normal DNS transaction



Time to check my
bank balance!!



www.cibc.ca





www.cibc.ca ?





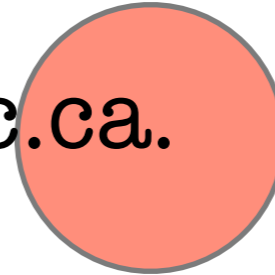
I am going to assume
Foxy meant:

www.cibc.ca.



I am going to assume
Foxy meant:

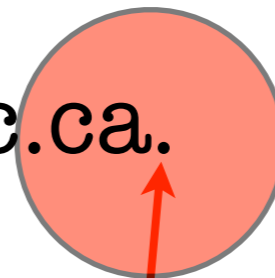
www.cibc.ca





I am going to assume
Foxy meant:

www.cibc.ca.



That last dot is actually significant.



www.cibc.ca ?



www.cibc.ca?

Hmmm. Have I not
seen that name before?



www.cibc.ca. ? Yes!

Here it is in my stash:

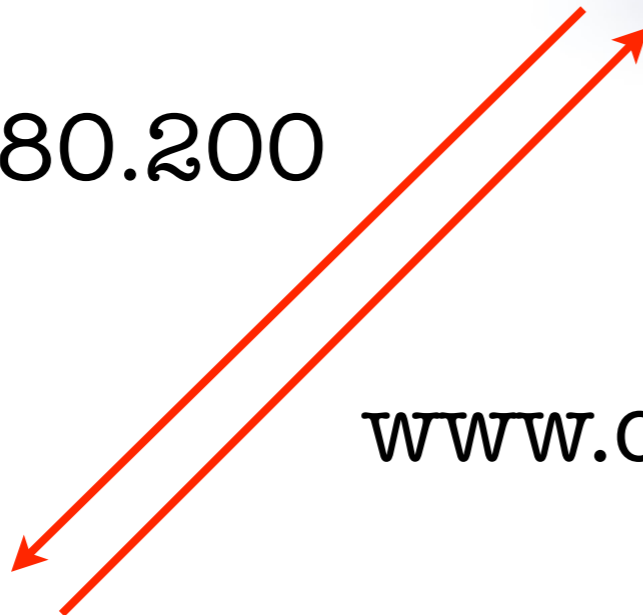
159.231.80.200



159.231.80.200



www.cibc.ca. ?



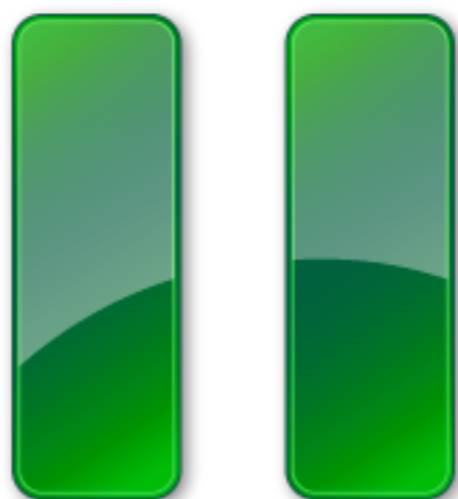


www.cibc.ca ?

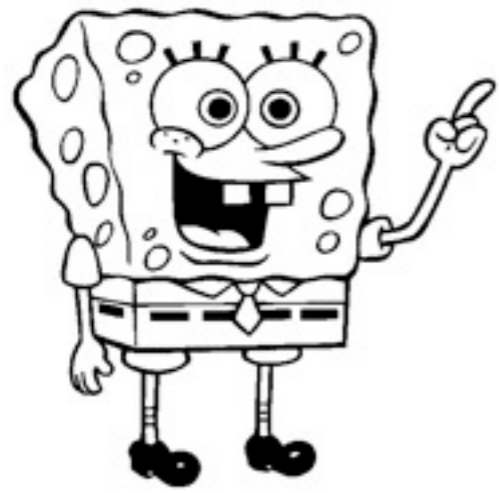


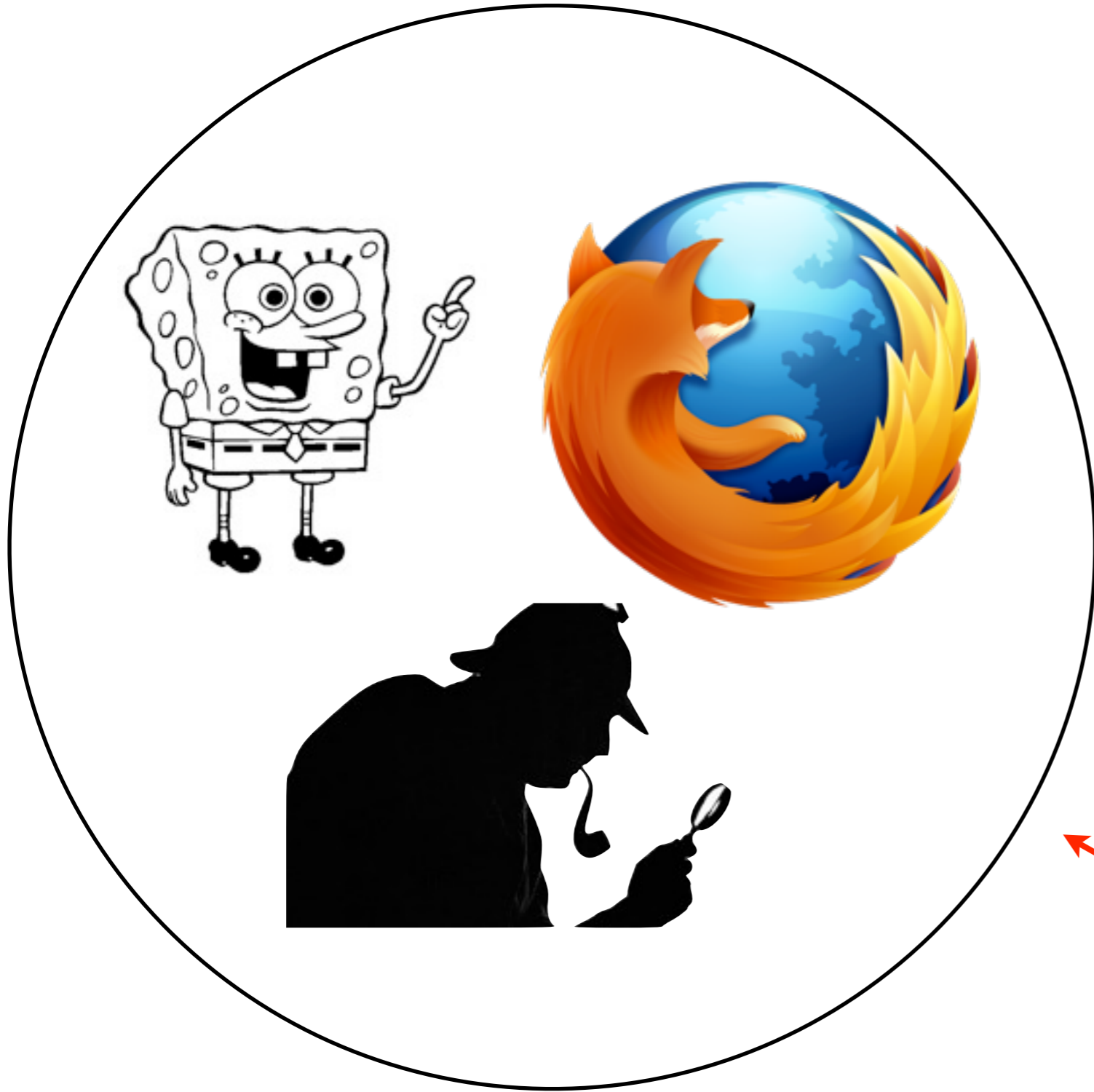
159.231.80.200





Security Audit time ...





Confidentiality ?

Integrity ?

Availability ?

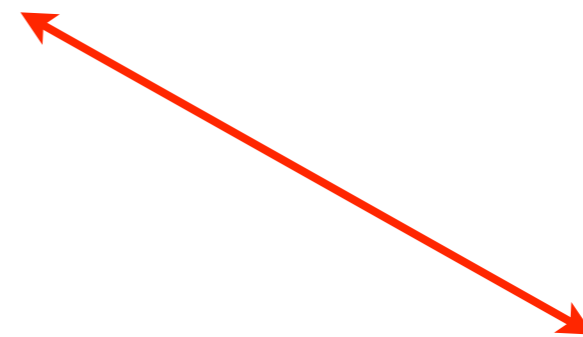




Confidentiality ?

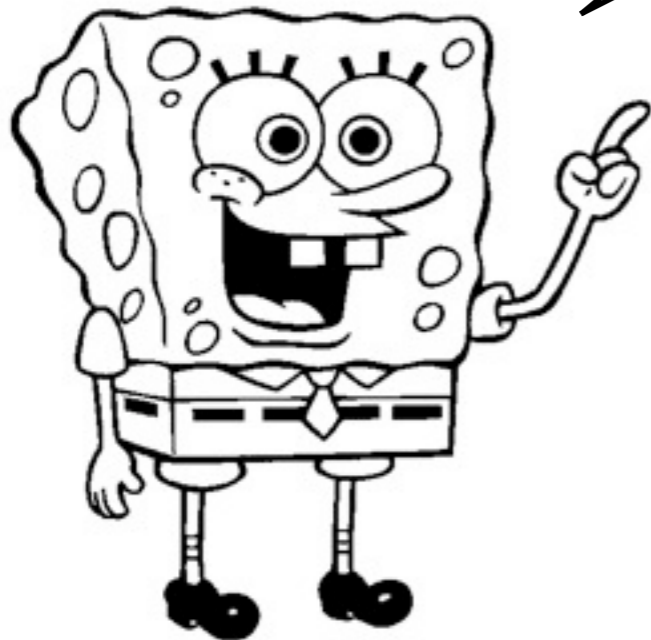
Integrity ?

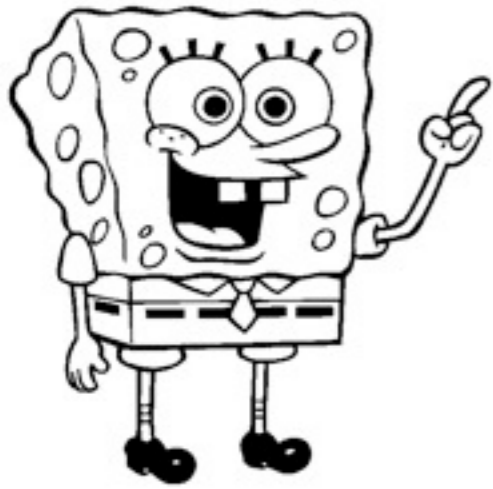
Availability ?





**Time to check my
credit card balance!!**





www.americanexpress.com





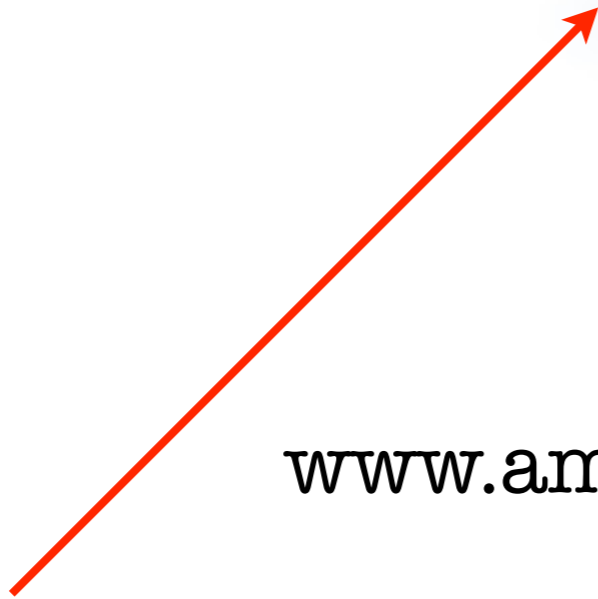
www.americanexpress.com



People always forget
that final dot. Sigh ...

[www.americanexpress.com.](http://www.americanexpress.com)





www.americanexpress.com. ?



[www.americanexpress.com.](http://www.americanexpress.com)

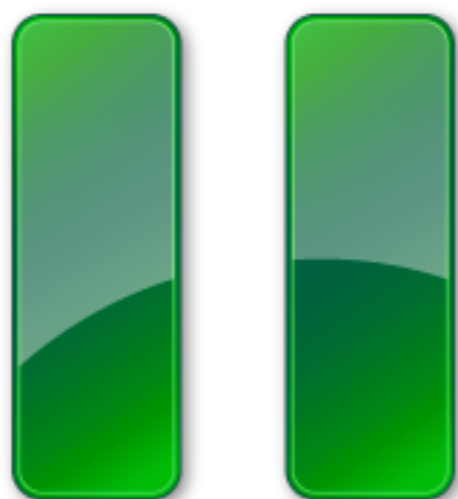
Hmmm. That looks like
a new one. Let's check.



www.americanexpress.com.

Rats! I do not have that
one stored. Time to go fish!





\$ dig -t ns .

; <<>> DiG 9.7.0-P1 <<>> -t ns .

;; global options: +cmd

;; Got answer:

;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 47571

;; flags: qr rd ra; QUERY: 1, ANSWER: 13, AUTHORITY: 0, ADDITIONAL: 0

;; QUESTION SECTION:

;. IN NS

;; ANSWER SECTION:

.	518375	IN	NS	g.root-servers.net.
.	518375	IN	NS	a.root-servers.net.
.	518375	IN	NS	e.root-servers.net.
.	518375	IN	NS	b.root-servers.net.
.	518375	IN	NS	k.root-servers.net.
.	518375	IN	NS	l.root-servers.net.
.	518375	IN	NS	c.root-servers.net.
.	518375	IN	NS	i.root-servers.net.
.	518375	IN	NS	d.root-servers.net.
.	518375	IN	NS	j.root-servers.net.
.	518375	IN	NS	f.root-servers.net.
.	518375	IN	NS	m.root-servers.net.
.	518375	IN	NS	h.root-servers.net.

```
$ dig +short -t ns .  
g.root-servers.net.  
a.root-servers.net.  
e.root-servers.net.  
b.root-servers.net.  
k.root-servers.net.  
l.root-servers.net.  
c.root-servers.net.  
i.root-servers.net.  
d.root-servers.net.  
j.root-servers.net.  
f.root-servers.net.  
m.root-servers.net.  
h.root-servers.net.  
  
$ dig +short -t a g.root-servers.net.  
192.112.36.4
```

```
$ cat @  
198.41.0.4  
128.9.0.107  
192.33.4.12  
128.8.10.90  
192.203.230.10  
192.5.5.241  
192.112.36.4  
128.63.2.53  
192.36.148.17  
192.58.128.30  
193.0.14.129  
198.32.64.12  
202.12.27.33
```



Query: A record for
www.americanexpress.com.

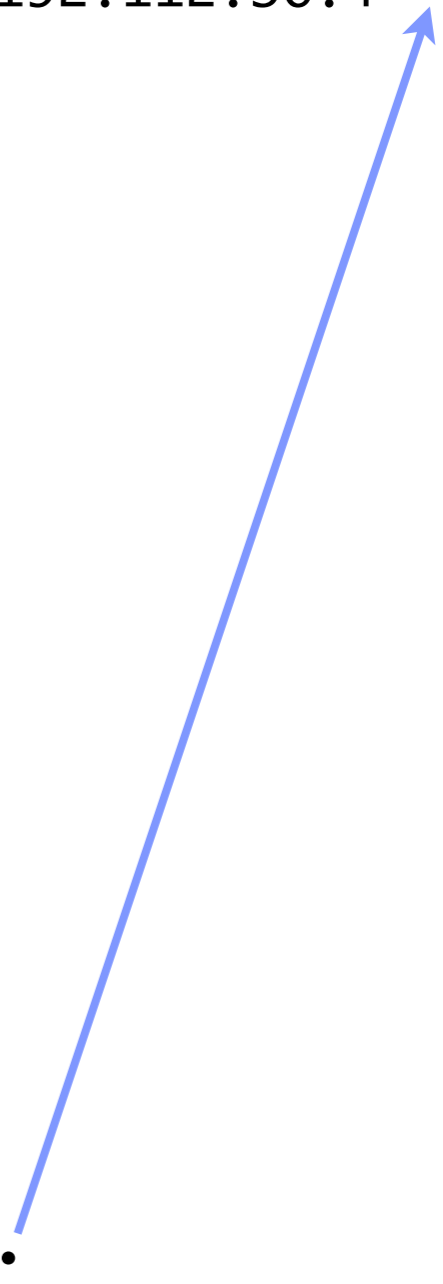
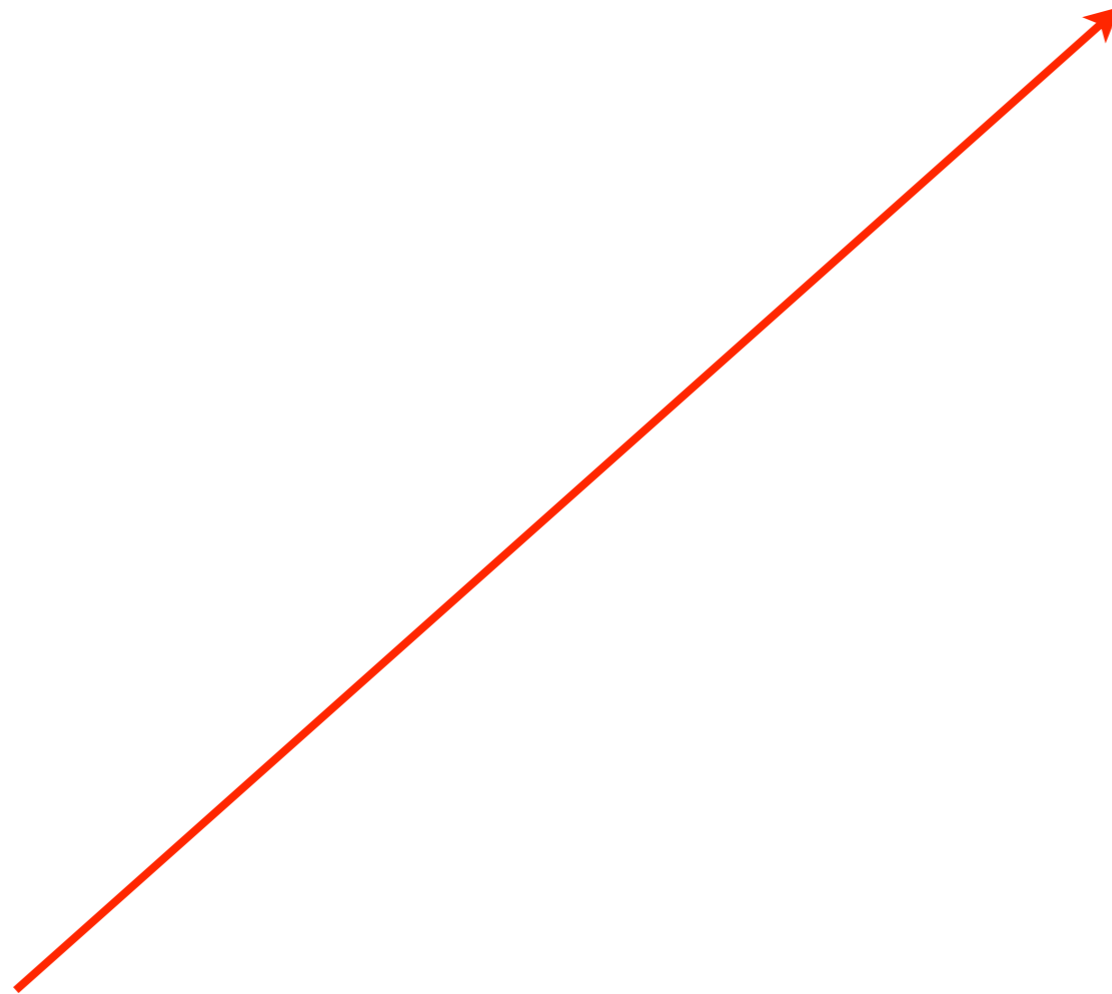


Root

g.root-servers.net.
192.112.36.4



www.americanexpress.com.



Query: A record for
www.americanexpress.com.

Response: NS records for com.



Root

g.root-servers.net.
192.112.36.4

d.gtld-servers.net.
a.gtld-servers.net.
k.gtld-servers.net.
c.gtld-servers.net.
m.gtld-servers.net.
i.gtld-servers.net.
l.gtld-servers.net.
f.gtld-servers.net.
e.gtld-servers.net.
h.gtld-servers.net.
g.gtld-servers.net.
b.gtld-servers.net.
j.gtld-servers.net.



www.americanexpress.com.

Query: A record for
www.americanexpress.com.



Root

g.root-servers.net.
192.112.36.4

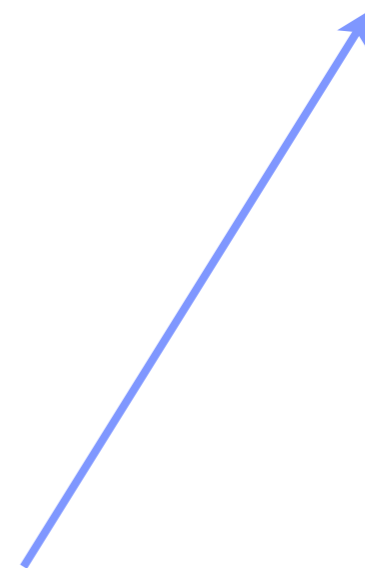
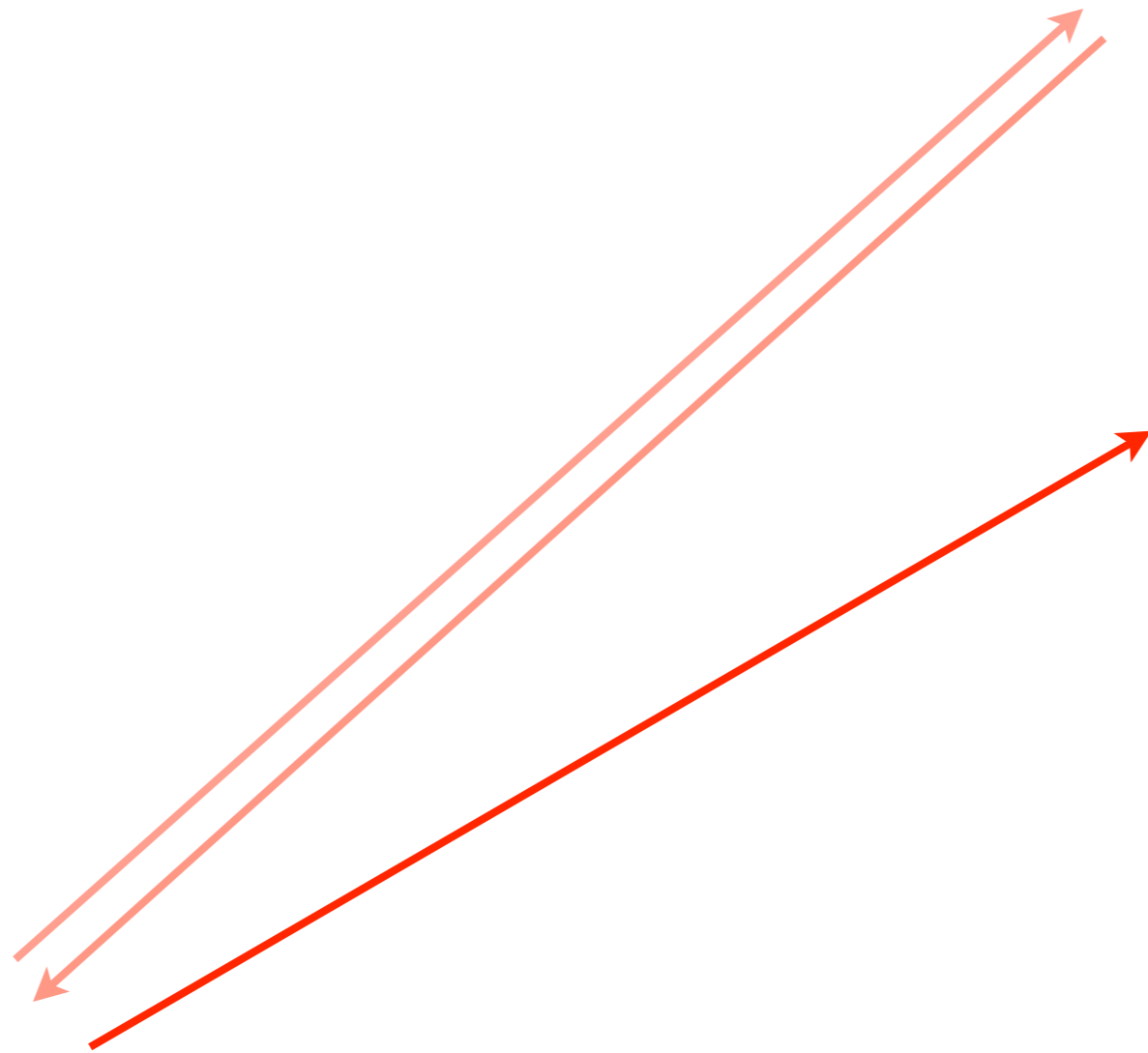


.com

f.gtld-servers.net.
192.35.51.30



www.americanexpress.com.



Query: A record for
www.americanexpress.com.

Response: NS records for
americanexpress.com.



Root

g.root-servers.net.
192.112.36.4



.com

f.gtld-servers.net.
192.35.51.30

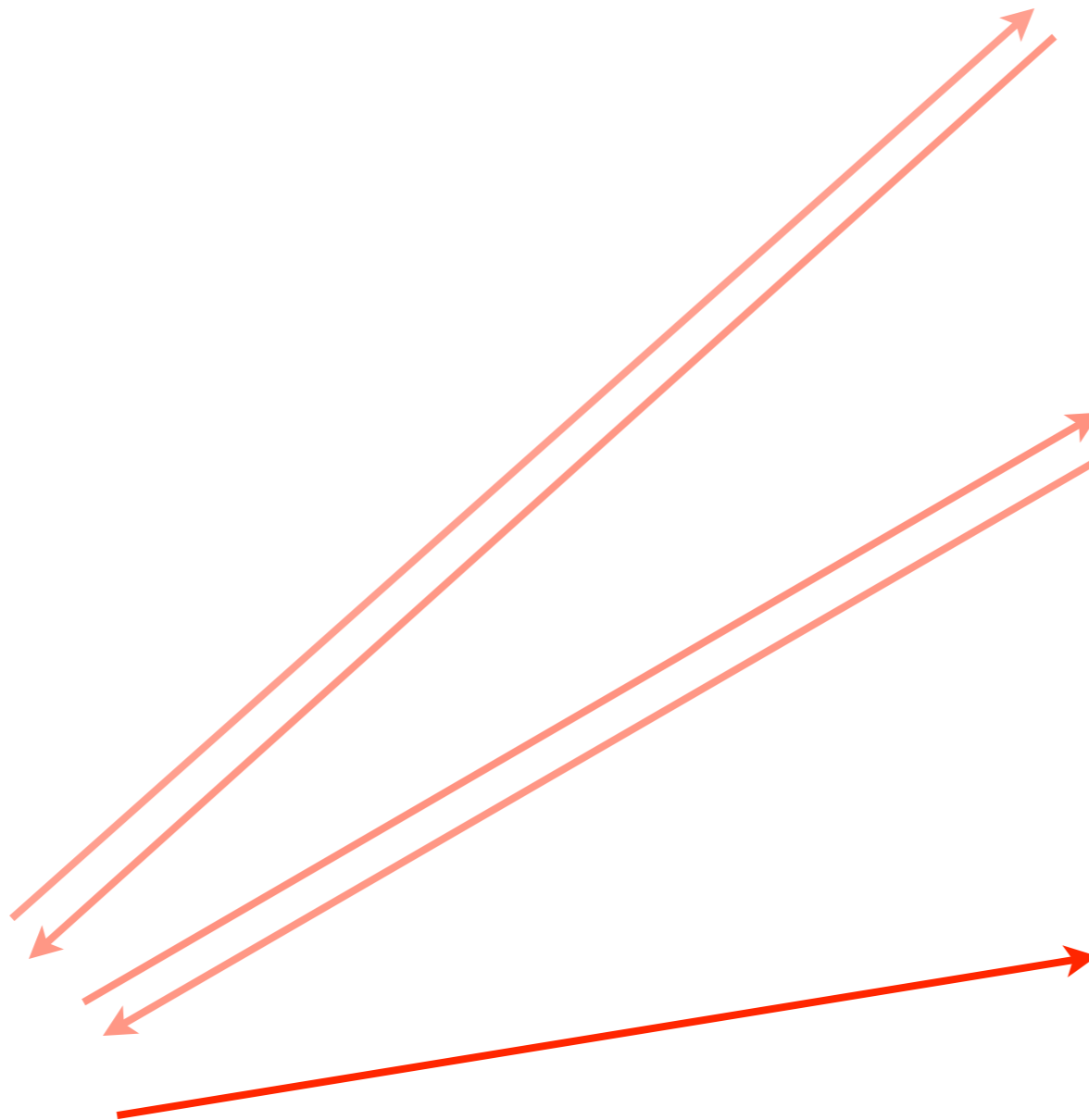
gw4.aexp.com.
gw5.aexp.com.
gw.aexp.com.
gw2.aexp.com.
gw3.aexp.com.



www.americanexpress.com.



Query: A record for
www.americanexpress.com.



Root

g.root-servers.net.
192.112.36.4



.com

f.gtld-servers.net.
192.35.51.30



americanexpress.com

gw5.aexp.com.
192.102.253.16

www.americanexpress.com.

Query: A record for
www.americanexpress.com.

Response: A record for
www.americanexpress.com.
12.29.100.148



Root

g.root-servers.net.
192.112.36.4



.com

f.gtld-servers.net.
192.35.51.30



americanexpress.com

gw5.aexp.com.
192.102.253.16



12.29.100.148

www.americanexpress.com.



www.americanexpress.com.
12.29.100.148

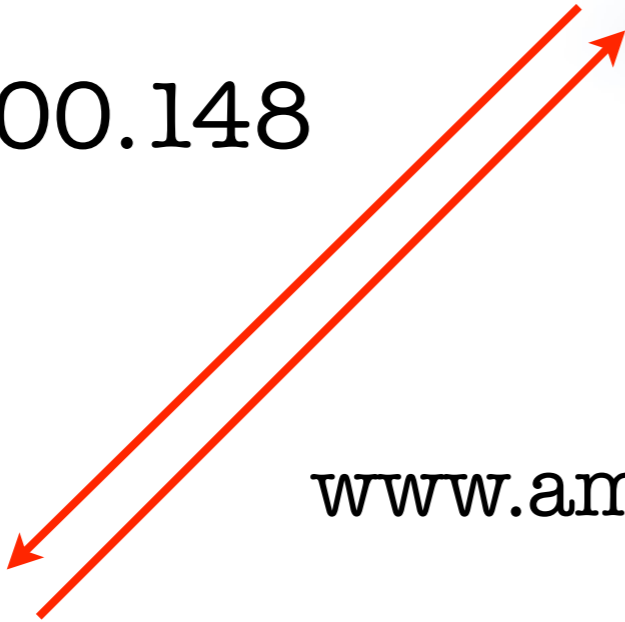
That took a bit of work. Let's
store it in our local stash.





12.29.100.148

www.americanexpress.com ?



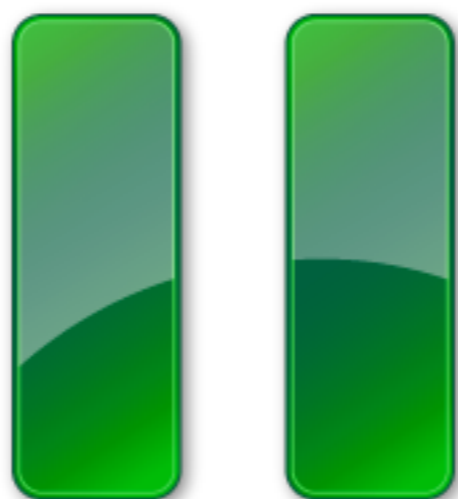


www.americanexpress.com



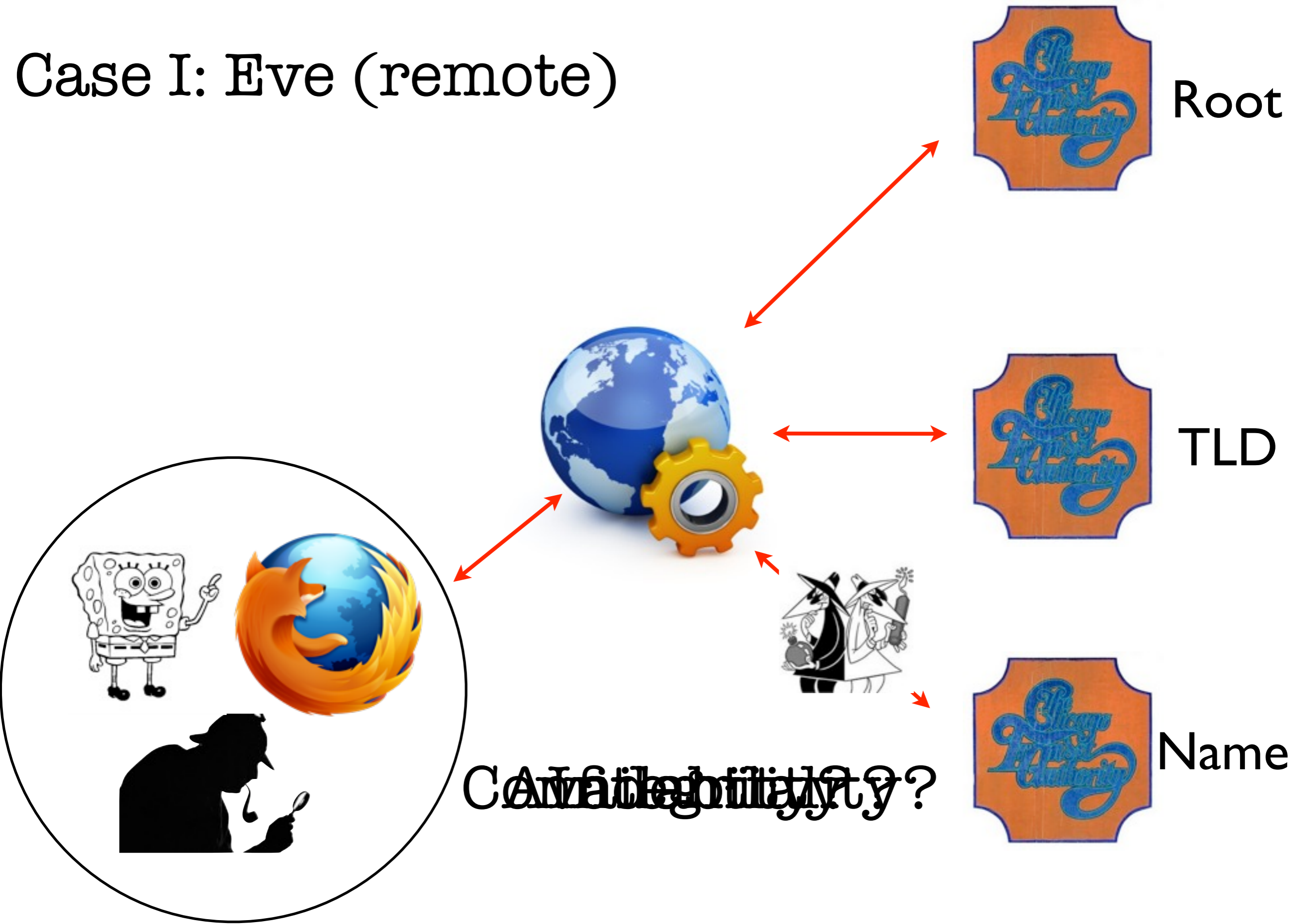
12.29.100.148





Security Audit time ...

Case I: Eve (remote)



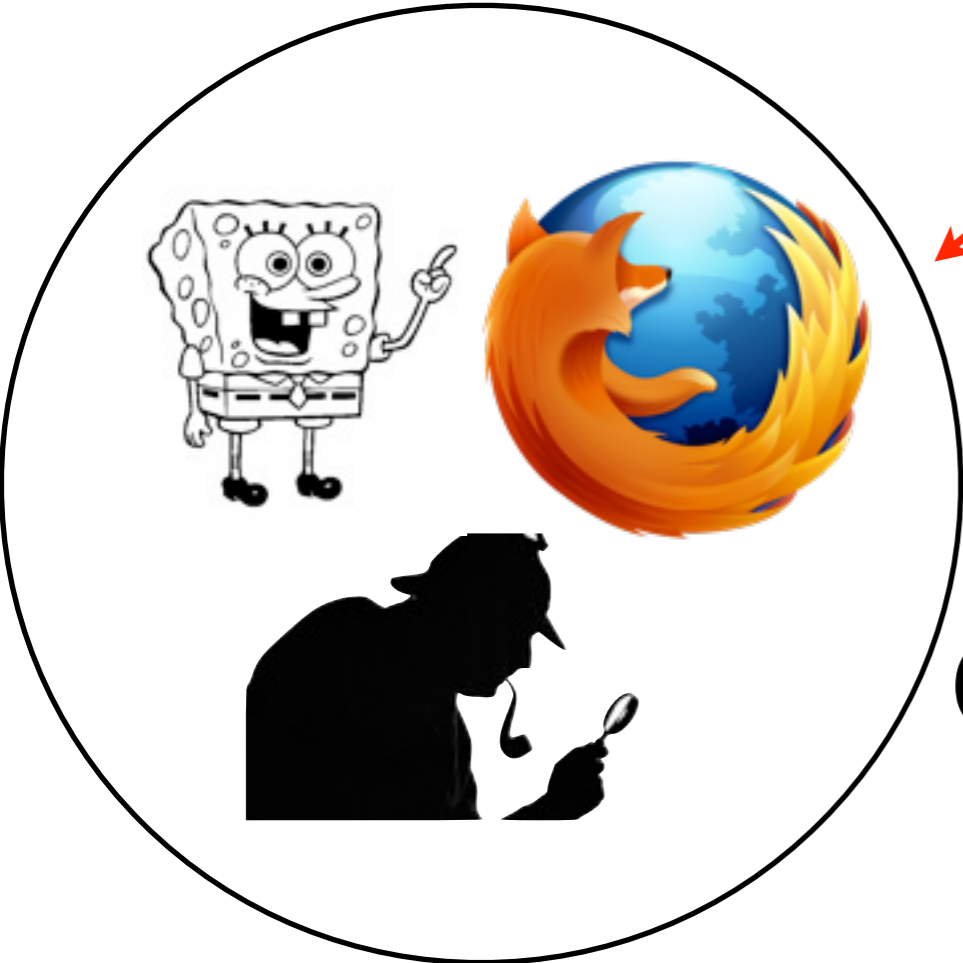
Root



TLD



Name



Can it be legitimized?



Case II: Eve (local)

Dan Kaminsky 2008-07-21

Cache Poisoning Attack



Can it be legitimized?



Root



TLD

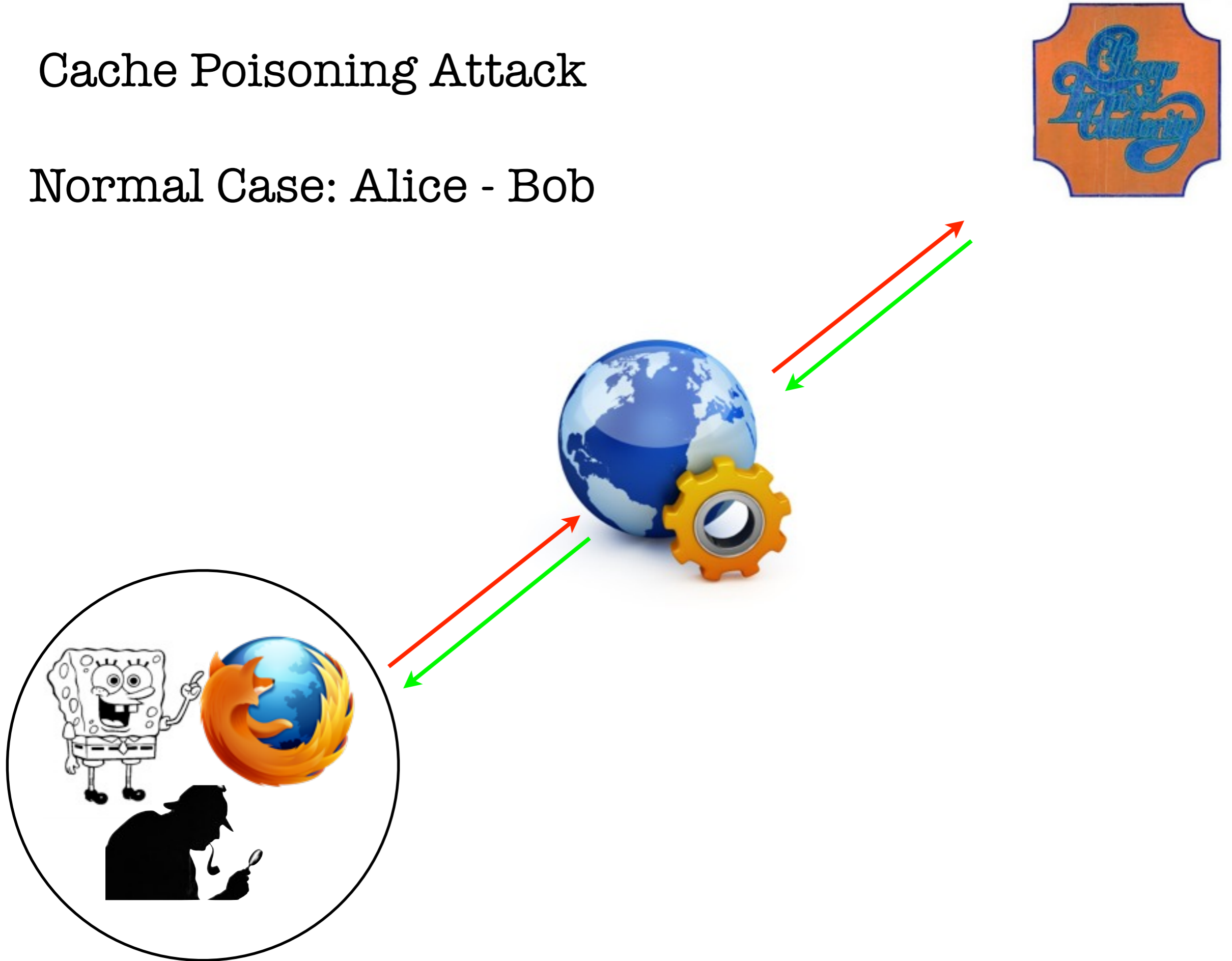


Name



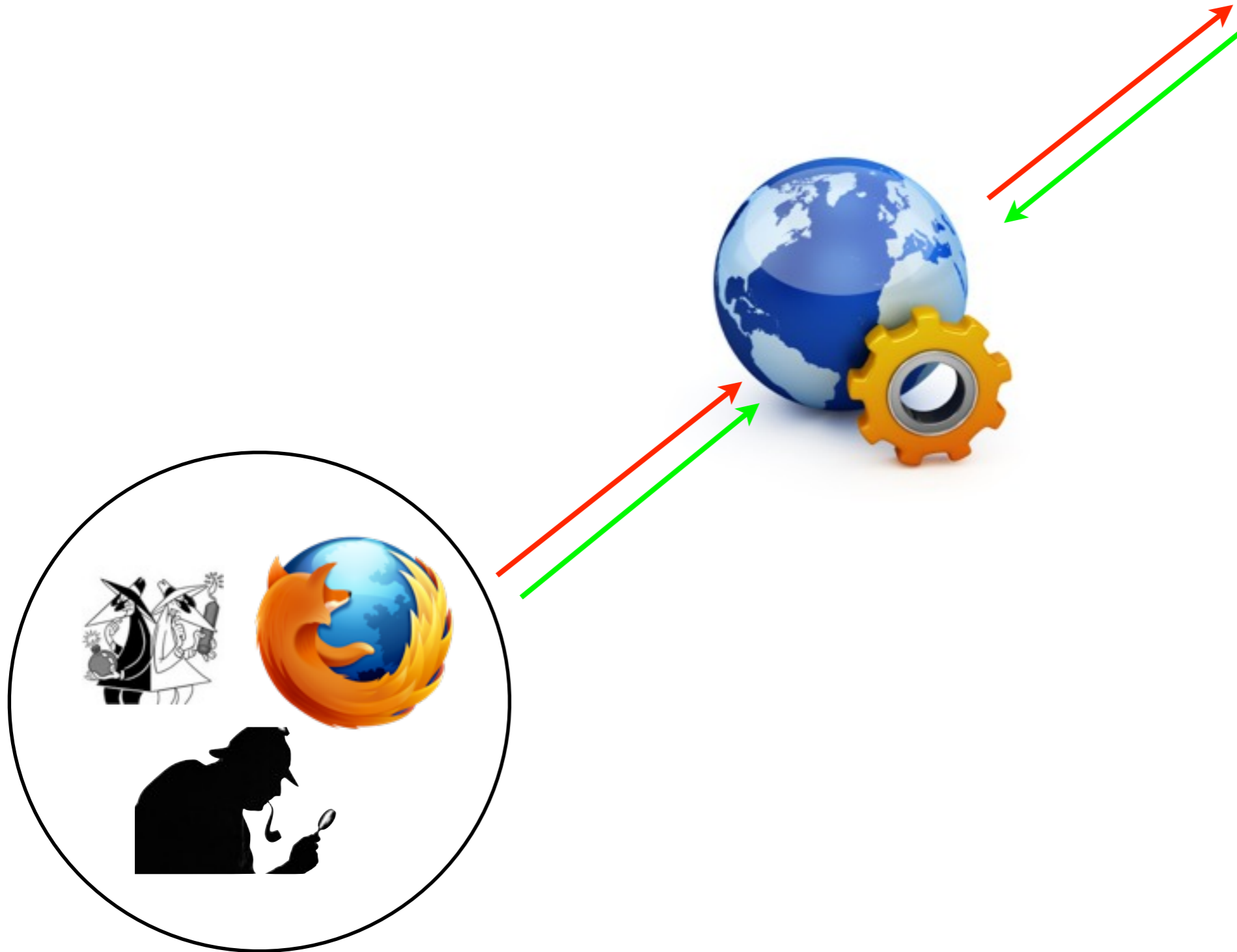
Cache Poisoning Attack

Normal Case: Alice - Bob



Cache Poisoning Attack

Imposter Case: Eve - Bob





Houston, we really do have a problem!



There be many DNS dragons...

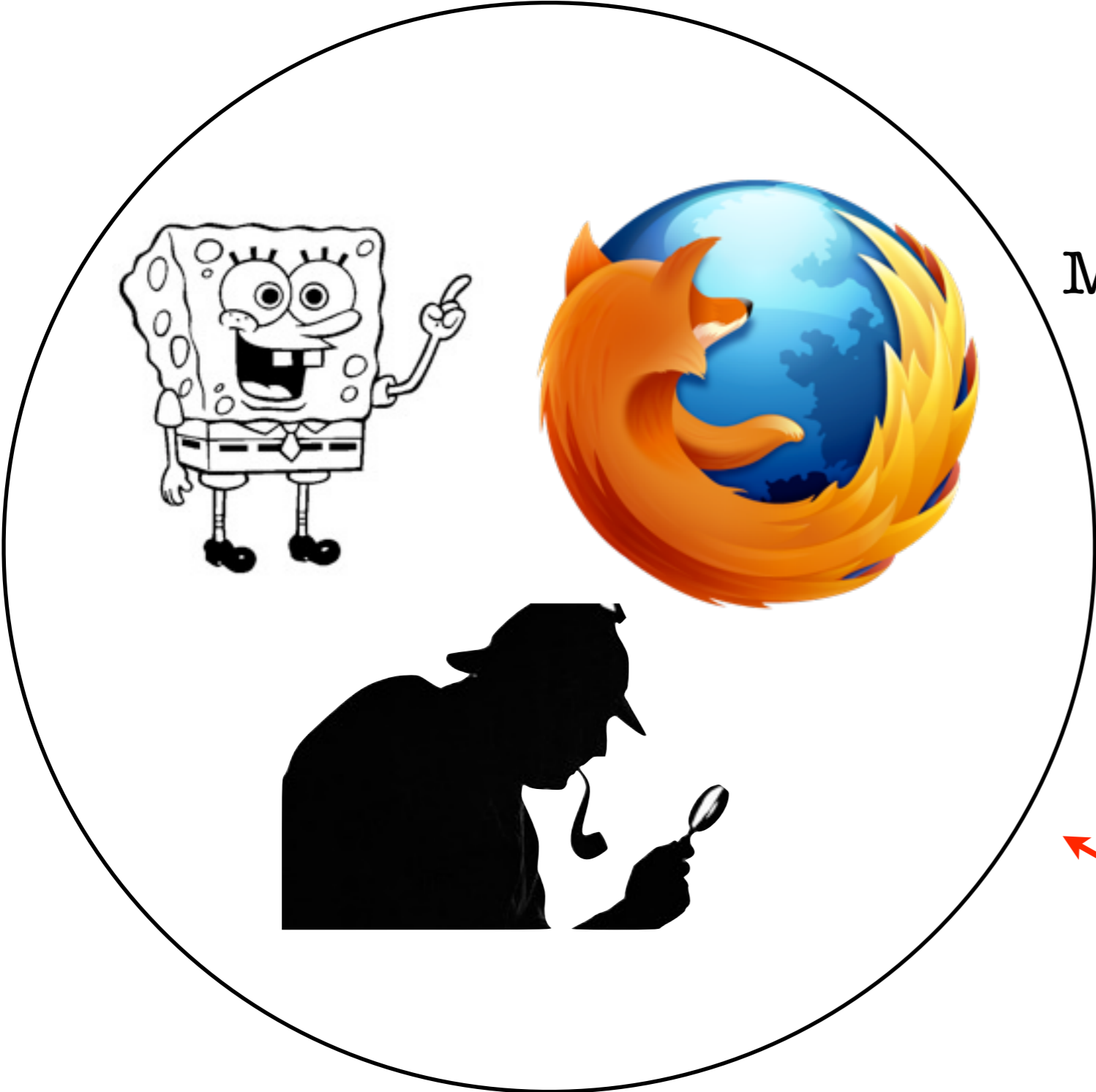
Passive Attacks



Man in Middle collects data regarding users queries

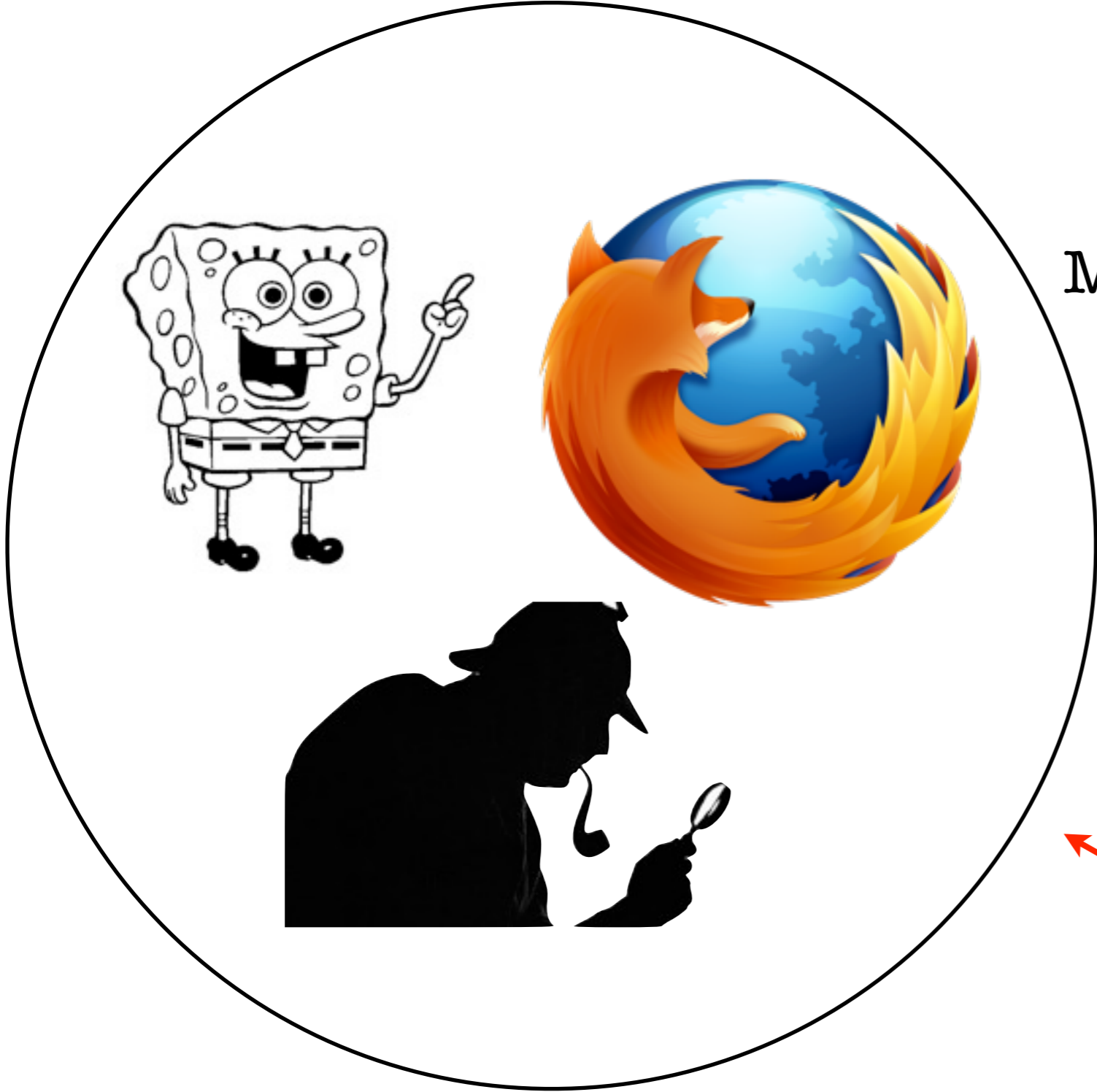
Active Attacks

Man in Middle changes data



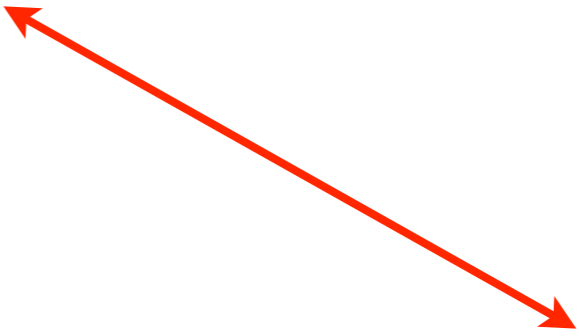
Active Attacks

Man in Middle changes data



Wikipedia is a free encyclopedia

Passive Attacks



Devious users can poison Caching Name servers

Solutions

For passive cache poisoning attacks

$$2^{16} = 65536$$

Random Query IDs makes guessing difficult



$$2^{16} * 2^{11} = 2^{27} = 134217728$$

Random QIDS & SRC ports minimizes risk





Modify DNS to use Cryptographic Tools



Confidentiality...



Would thwart passive man in the middle attacks



Integrity...



Would thwart all of the spoofing attacks



<6>

DNS Security Extensions

aka DNSSEC

What is it?

“The DNS Security extensions provide origin authentication and integrity protection for DNS data, as well as a means of public key distribution. These extensions do not provide confidentiality.”



“ It is a set of extensions to DNS, which provide:
a. origin authentication of DNS data
b. data integrity
c. authenticated denial of existence ”



History

1987: Regular DNS standardized [RFC 1034,1035]

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1990: DNS vulnerabilities come to light
[Steve Bellovin, Bell Laboratories]

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1995: IETF strikes a DNSEXT working group

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**1997: IETF Domain Name Security Extensions
[RFC 2065]**

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**1999: RFC 2535 supercedes RFC 2065
implementation problems**

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1999: end of year, ISC ships bind with RFC2535 exts.

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1999: RFC 2535 supercedes RFC 2065

1999: end of year, ISC ships bind with RFC2535 exts.

**2001: RFC2535 key handling operational problems.
Restart! Writing, Drafting, Publishing**

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- 2010: All root name servers are DNSSEC ready
- 2012: .ca name servers are DNSSEC ready**

Recent Uptake

End of 2009: ~ 1000

End of 2010: ~2500

Number of registered domains: ~200,000,000

DNSSEC: Objectives

<0>

N.B:

DNSSEC is designed to **detect** attacks and not necessarily to **prevent** them.

<1>

Origin Authentication of DNS Data

Client can trust that the authoritative name server really is the authority for a certain zone.

Authenticity is a case of Integrity

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<2>

Data Integrity

Either end can detect if query or response has been modified by an unauthorized third party

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<3>

Authenticated Denial of Existence

The client can be sure not only of the response from data which exists but also that certain data does NOT exist.

Authenticated Denial of Existence

The client can be sure not only of the response from data which exists but also that certain data does **NOT** exist.

This is necessary to prevent certain forms of attacks

<4>

Backward Compatibility with Regular DNS

DNSSEC and non DNSSEC
environments both need to
interoperate.

DNSSEC: Non Objectives

<1>

~~Confidentiality~~



All DNSSEC traffic is plaintext

... that all data in the DNS is thus visible. Accordingly, DNSSEC is not designed to provide confidentiality, access control lists, or other means of differentiating between inquirers.

RFC 4033

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RFC 4033

This design decision has security implications



Availability



DNSSEC provides no protection against denial of service attacks. Security-aware resolvers and security-aware name servers are vulnerable to an additional class of denial of service attacks based on cryptographic operations.

RFC 4033

DNSSEC provides **no protection against denial of service attacks**. Security-aware resolvers and security-aware name servers are **vulnerable** to an **additional class of denial of service attacks** based on cryptographic operations.

RFC 4033

There are no extra features, over and above regular DNS, to prevent DoS or buffer overflow attacks

DNSSEC Specifications

New Resource Records

Regular Resource Records

NS - name server delegation

A - IP address

MX - mail server name & priority

CNAME - name alias

TXT - text description

SOA - start of authority

New Resource Records

DNSSEC uses public key cryptography to sign and authenticate DNS resource record sets (RRsets).

RFC 4034

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RFC 4034

DNSKEY - for storing public keys

A zone signs its authoritative RRsets by using a private key and stores the corresponding public key in a DNSKEY RR.

RFC 4034

A zone **signs** its authoritative RRsets by using a **private key** and stores the corresponding **public key** in a **DNSKEY RR**.

RFC 4034

Note:

The DNSKEY RR is only intended to store DNS related public keys. It **MUST NOT** be used to store generic public keys and certificates

RFC 4034

DNSKEY - algorithms are also stored

```
# dig +multiline vix.com DNSKEY
```

```
;; Got answer:
```

```
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 11667
```

```
;; flags: qr rd ra; QUERY: 1, ANSWER: 2, AUTHORITY: 0, ADDITIONAL: 0
```

```
;; QUESTION SECTION:
```

```
;vix.com.                IN DNSKEY
```

```
;; ANSWER SECTION:
```

```
vix.com.                2894 IN DNSKEY 257 3 5 (
                          AwEAAbKW5zsYMBUX4MS0yq3MNM4312c7WEF1Af2Iy20/
                          A+U+h7F3Etb1BDJVs/LgtdjsE3JHak51iRaELL0oEvVe
                          RIIa1UjNvXieia+QV1nlSas8LcXya0X0YA2Jfzez0pEW
                          ArN1QLhkgVDPAsEwKLzYfVjW78CF10ZnYxbBWXwKgb4z
                          ) ; key id = 26437
vix.com.                2894 IN DNSKEY 256 3 5 (
                          BEAAAA06wBt1U39U8meHca3JBCWixBi8BvZLMJZp51/5
                          vViM2+fh93XF1SqJaAaqgX6PszTPU1E1vuTV2xTV4uQj
                          UTaFv8qDnsjbfXVusE1v+0aQpSVuP8GjI28cGi9PgA0c
                          z2ACdiD2XVbYKUDTJb+pqoE/o3Z6FjKf6ByTkJUI5x9D
                          lw==
                          ) ; key id = 63066
```



```
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```
;; QUESTION SECTION:
```

```
vix.com. IN DNSKEY
```

```
;; ANSWER SECTION:
```

```
vix.com. 2894 IN DNSKEY 257 3 5 (
    AwEAAbKW5zsYMBUX4MS0yq3MNM4312c7WEF1Af2Iy20/
    A+U+h7F3Etb1BDJVs/LgtdjsE3JHak51iRaELL0oEvVe
    RIIa1UjNvXieia+QV1nlSas8LcXya0X0YA2Jfzez0pEW
    ArN1QLhkgVDPAsEwKLzYfVjW78CF10ZnYxbBWXwKgb4z
    ) ; key id = 26437
vix.com. 2894 IN DNSKEY 256 3 5 (
    BEAAAA06wBt1U39U8meHca3JBCWixBi8BvZLMJZp51/5
    vViM2+fh93XF1SqJaAaqgX6PszTPU1E1vuTV2xTV4uQj
    UTaFv8qDnsjbfXVusE1v+0aQpSVuP8GjI28cGi9PgA0c
    z2ACdiD2XVbYKUDTJb+pqoE/o3Z6FjKf6ByTkJUI5x9D
    lw==
    ) ; key id = 63066
```

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```
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```

```
;vix.com. IN DNSKEY
```

```
;; ANSWER SECTION:
```

```
vix.com.
```

```
2894 IN DNSKEY 257 3 5 (
```

```
AwEAAbKW5zsYMBUX4MS0yq3MNM4312c7WEF1Af2Iy20/  
A+U+h7F3Etb1BDJVs/LgtdjsE3JHak51iRaELL0oEvVe  
RIIa1UjNvXieia+QV1nlSas8LcXya0X0YA2Jfvez0pEW  
ArN1QLhkgVDPAsEwKLzYfvjW78CF10ZnYxbBWXwKgb4z
```

```
) ; key id = 26437
```

```
vix.com.
```

```
2894 IN DNSKEY 256 3 5 (
```

```
BEAAAA06wBt1U39U8meHca3JBCWixBi8BvZLMJZp51/5  
vViM2+fh93XF1SqJaAaqgX6PszTPU1E1vuTV2xTV4uQj  
UTaFv8qDnsjbfXVusE1v+0aQpSVuP8GjI28cGi9PgA0c  
z2ACdiD2XVbYKUDTJb+pqoE/o3Z6FjKf6ByTkJUI5x9D
```

```
lw==
```

```
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vix.com.
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2894 IN DNSKEY 257 3 5 (
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```
AwEAAbKW5zsYMBUX4MS0yq3MNM4312c7WEF1Af2Iy20/  
A+U+h7F3Etb1BDJVs/LgtdjsE3JHak51iRaELL0oEvVe  
RIIa1UjNvXieia+QV1nlSas8LcXya0X0YA2Jfvez0pEW  
ArN1QLhkgVDPAsEwKLzYfVjW78CF10ZnYxbBWXwKgb4z  
) ; key id = 26437
```

```
vix.com.
```

```
2894 IN DNSKEY 256 3 5 (
```

```
BEAAAA06wBt1U39U8meHca3JBCWixBi8BvZLMJZp51/5  
vViM2+fh93XF1SqJaAaqgX6PszTPU1E1vuTV2xTV4uQj  
UTaFv8qDnsjbfXVusE1v+0aQpSVuP8GjI28cGi9PgA0c  
z2ACdiD2XVbYKUDTJb+pqoE/o3Z6FjKf6ByTkJUI5x9D  
lw==  
) ; key id = 63066
```

RRSIG - signatures are stored here

Signature: DNS data + owners private key

RRSIG - other fields as well

RRSIG - sig. inception time

RRSIG - sig. expiration time

RRSIG - type covered

RRSIG - algorithm

RRSIG - key tag

RRSIG - signers name

RRSIG - Data signature

```
# drill -D vix.com @ns.sql1.vix.com. rrsig
```

```
;; QUESTION SECTION:
```

```
;; vix.com.      IN      RRSIG
```

```
;; ANSWER SECTION:
```

```
vix.com.      3600      IN      RRSIG   A 5 2 3600 20110629030401 20110331030401 63066 vix.com.
```

```
CsxzLHeqDLi2XXKqGALXYn4lbmZrqkDzCYegv6EiZQFpPHG8oVdxvqJDCczpVHF3mykB05uHntpyo0S4om34l8fkIuVKViE6c/3b  
+j3jiJIIfXbFYPqM501NChRf/SwkBqsmKRj4jbTp3jCicUG6M3lyNWe5B2CjVd9hEUmzrbjY= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   MX 5 2 3600 20110629030401 20110331030401 63066 vix.com.
```

```
urjAd1NVJKNfUOI/l0aJRNEQJJfexjnwRTcyzcZmVvxvV5FlqLT904aIzcKMPnM2L3FWpf  
+F0Tzfjr9Cb46pUHRj9LApaKxAH7RT0GKz7t2kVd8bD62LbhkFiVVlvqVTBBIhHinzAx8wPSCaU2saAt4fYc  
+0w86it8IKBuwZyjE= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   DNSKEY 5 2 3600 20110629030401 20110331030401 26437 vix.com.
```

```
QvyoIbB1fTtge9aBTj88oBBFUnfLdGxGoyABG3bkPDAiDB5TUgJa68UDcF5k9c5fQEHZA6rd52QRxkPKy0hb5Reh64cZMjzMBZxZ  
aivxX+W+hmkEk9ztSgWaotNBw2RHechItBI4/IPZWRXGNPr1IIduI8KC+dm96tf404BraAU= ;{id =  
26437}
```

```
# drill -D vix.com @ns.sql1.vix.com. rrsig
```

```
;; QUESTION SECTION:
```

```
;; vix.com.      IN      RRSIG
```

```
;; ANSWER SECTION:
```

```
vix.com.      3600      IN      RRSIG  A 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
CsxzLHeqDLi2XXKqGALXYn4lbmZrqkDzCYegv6EiZQFpPHG8oVdxvqJDCczpVHF3mykB05uHntpyo0S4om34l8fkIuVKViE6c  
/3b+j3jiJIIfXbFYPqM501NChRf/SwkBqsmKRj4jbTp3jCicUG6M3lyNWe5B2CjVd9hEUmzrbjY= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG  MX 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
urjAd1NVJKNfUOI/l0aJRNEQJJfexjnwRTcyzcZmVvxvV5FlqLT904aIzcKMPnM2L3FWpf  
+F0Tzfjr9Cb46pUhrj9LApaKxAH7RT0GKz7t2kVd8bD62LbhkFiVVlvqVTBBIhHinzAx8wPSCaU2saAt4fYc  
+0w86it8IKBuwZyjE= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG  DNSKEY 5 2 3600 20110629030401 20110331030401 26437
```

```
vix.com.  
QvyoIbB1fTtge9aBTj88oBBFUnfLdGxGoyABG3bkPDAiDB5TUgJa68UDcF5k9c5fQEHZA6rd52QRxkPKy0hb5Reh64cZMjzMB  
ZxZaivxX+W+hmkEk9ztSgWaotNBw2RHechItBI4/IPZWRXGNPr1IIduI8KC+dm96tf404BraAU= ;{id =  
26437}
```

```
# drill -D vix.com @ns.sql1.vix.com. rrsig
```

```
;; QUESTION SECTION:
```

```
;; vix.com.      IN      RRSIG
```

```
;; ANSWER SECTION:
```

```
vix.com.      3600      IN      RRSIG  A 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
CsxzLHeqDLi2XXKqGALXyn4lbmZrqkDzCYegv6EiZQFpPHG8oVdxvqJDCczpVHF3mykB05uHntpyo0S4om34l8fkIuVKViE6c  
/3b+j3jiJIIfXbFYPqM501NChRf/SwkBqsmKRj4jbTp3jCicUG6M3lyNWe5B2CjVd9hEUmzrbjY= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG  MX 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
urjAd1NVJKNfUOI/l0aJRNEQJJfexjnwRTcyzcZmVvxvV5FlqLT904aIzcKMPnM2L3FWpf  
+F0Tzfjr9Cb46pUhrj9LApaKxAH7RT0GKz7t2kVd8bD62LbhkFiVVlvqVTBBIhHinzAx8wPSCaU2saAt4fYc  
+0w86it8IKBuwZyjE= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG  DNSKEY 5 2 3600 20110629030401 20110331030401 26437
```

```
vix.com.  
QvyoIbB1fTtge9aBTj88oBBFUnfLdGxGoyABG3bkPDAiDB5TUgJa68UDcF5k9c5fQEHZA6rd52QRxkPKy0hb5Reh64cZMjzMB  
ZxZaivxX+W+hmkEk9ztSgWaotNBw2RHechItBI4/IPZWRXGNPr1IIduI8KC+dm96tf404BraAU= ;{id =  
26437}
```



```
# drill -D vix.com @ns.sql1.vix.com. rrsig
```

```
;; QUESTION SECTION:
```

```
;; vix.com.      IN      RRSIG
```

```
;; ANSWER SECTION:
```

```
vix.com.      3600      IN      RRSIG   A 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
CsxzLHeqDLi2XXKqGALXyn4lbmZrqkDzCYegv6EiZQFpPHG8oVdxvqJDCczpVHF3mykB05uHntpyo0S4om34l8fkIuVKViE6c  
/3b+j3jiJIIfXbFYPqM501NChRf/SwkBqsmKRj4jbTp3jCicUG6M3lyNWe5B2CjVd9hEUmzrbjY= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   MX 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
urjAd1NVJKNfUOI/l0aJRNEQJJfexjnwRTcyzcZmVvxvV5FlqLT904aIzcKMPnM2L3FWpf  
+F0Tzfjr9Cb46pUhrj9LApaKxAH7RT0GKz7t2kVd8bD62LbhkFiVVlvqVTBBIhHinzAx8wPSCaU2saAt4fYc  
+0w86it8IKBuwZyjE= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   DNSKEY 5 2 3600 20110629030401 20110331030401 26437
```

```
vix.com.  
QvyoIbB1fTtge9aBTj88oBBFUnfLdGxGoyABG3bkPDAiDB5TUgJa68UDcF5k9c5fQEHZA6rd52QRxkPKy0hb5Reh64cZMjzMB  
ZxZaivxX+W+hmkEk9ztSgWaotNBw2RHechItBI4/IPZWRXGNPr1IIduI8KC+dm96tf404BraAU= ;{id =  
26437}
```

```
# drill -D vix.com @ns.sql1.vix.com. rrsig
```

```
;; QUESTION SECTION:
```

```
;; vix.com.      IN      RRSIG
```

```
;; ANSWER SECTION:
```

```
vix.com.      3600      IN      RRSIG   A 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
CsxzLHeqDLi2XXKqGALXYn4lbmZrqkDzCYegv6EiZQFpPHG8oVdxvqJDCczpVHF3mykB05uHntpyo0S4om34l8fkIuVKViE6c  
/3b+j3jiJIIfXbFYPqM501NChRf/SwkBqsmKRj4jbTp3jCicUG6M3lyNWe5B2CjVd9hEUmzrbjY= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   MX 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
urjAd1NVJKNfUOI/l0aJRNEQJJfexjnwRTcyzcZmVvxvV5FlqLT904aIzcKMPnM2L3FWpf  
+F0Tzfjr9Cb46pUhrj9LApaKxAH7RT0GKz7t2kVd8bD62LbhkFiVVlvqVTBBIhHinzAx8wPSCaU2saAt4fYc  
+0w86it8IKBuwZyjE= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   DNSKEY 5 2 3600 20110629030401 20110331030401 26437
```

```
vix.com.  
QvyoIbB1fTtge9aBTj88oBBFUnfLdGxGoyABG3bkPDAiDB5TUgJa68UDcF5k9c5fQEHZA6rd52QRxkPKy0hb5Reh64cZMjzMB  
ZxZaivxX+W+hmkEk9ztSgWaotNBw2RHechItBI4/IPZWRXGNPr1IIduI8KC+dm96tf404BraAU= ;{id =  
26437}
```

```
# drill -D vix.com @ns.sql1.vix.com. rrsig
```

```
;; QUESTION SECTION:
```

```
;; vix.com.      IN      RRSIG
```

```
;; ANSWER SECTION:
```

```
vix.com.      3600      IN      RRSIG   A 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
CsxzLHeqDLi2XXKqGALXYn4lbmZrqkDzCYegv6EiZQFpPHG8oVdxvqJDCczpVHF3mykB05uHntpyo0S4om34l8fkIuVKViE6c  
/3b+j3jiJIIfXbFYPqM501NChRf/SwkBqsmKRj4jbTp3jCicUG6M3lyNWe5B2CjVd9hEUmzrbjY= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   MX 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
urjAd1NVJKNfUOI/l0aJRNEQJJfexjnwRTcyzcZmVvxvV5FlqLT904aIzcKMPnM2L3FWpf  
+F0Tzfjr9Cb46pUhrj9LApaKxAH7RT0GKz7t2kVd8bD62LbhkFiVVlvqVTBBIhHinzAx8wPSCaU2saAt4fYc  
+0w86it8IKBuwZyjE= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   DNSKEY 5 2 3600 20110629030401 20110331030401 26437
```

```
vix.com.  
QvyoIbB1fTtge9aBTj88oBBFUnfLdGxGoyABG3bkPDAiDB5TUgJa68UDcF5k9c5fQEHZA6rd52QRxkPKy0hb5Reh64cZMjzMB  
ZxZaivxX+W+hmkEk9ztSgWaotNBw2RHechItBI4/IPZWRXGNPr1IIduI8KC+dm96tf404BraAU= ;{id =  
26437}
```

```
# drill -D vix.com @ns.sql1.vix.com. rrsig
```

```
;; QUESTION SECTION:
```

```
;; vix.com.      IN      RRSIG
```

```
;; ANSWER SECTION:
```

```
vix.com.      3600      IN      RRSIG   A 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
CsxzLHeqDLi2XXKqGALXYn4lBmZrqkDzCYegv6EiZQFpPHG8oVdxvqJDCczpVHF3mykB05uHntpyo0S4om34l8fkIuVKViE6c  
/3b+j3jiJIIfXbFYPqM501NChRf/SwkBqsmKRj4jbTp3jCicUG6M3lyNWe5B2CjVd9hEumzrbjY= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   MX 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
urjAd1NVJKNfUOI/l0aJRNEQJJfexjnwRTcyzcZmVvxv5FlqLT904aIzcKMPnM2L3FWpf  
+F0Tzfjr9Cb46pUhrj9LApaKxAH7RT0GKz7t2kVd8bD62LbhkFiVVlvqVTBBIhHinzAx8wPSCaU2saAt4fYc  
+0w86it8IKBuwZyjE= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   DNSKEY 5 2 3600 20110629030401 20110331030401 26437
```

```
vix.com.  
QvyoIbB1fTtge9aBTj88oBBFUnfLdGxGoyABG3bkPDAiDB5TUgJa68UDcF5k9c5fQEHZA6rd52QRxkPKy0hb5Reh64cZMjzMB  
ZxZaivxX+W+hmkEk9ztSgWaotNBw2RHechItBI4/IPZWRXGNPr1IIduI8KC+dm96tf404BraAU= ;{id =  
26437}
```

```
# drill -D vix.com @ns.sql1.vix.com. rrsig
```

```
;; QUESTION SECTION:
```

```
;; vix.com.      IN      RRSIG
```

```
;; ANSWER SECTION:
```

```
vix.com.      3600      IN      RRSIG   A 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
CsxzLHeqDLi2XXKqGALXYn4lbmZrqkDzCYegv6EiZQFpPHG8oVdxvqJDCczpVHF3mykB05uHntpyo0S4om34l8fkIuVKViE6c  
/3b+j3jiJIIfXbFYPqM501NChRf/SwkBqsmKRj4jbTp3jCicUG6M3lyNWe5B2CjVd9hEUmzrbjY= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   MX 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
urjAd1NVJKNfUOI/l0aJRNEQJJfexjnwRTcyzcZmVvxvV5FlqLT904aIzcKMPnM2L3FWpf  
+F0Tzfjr9Cb46pUhrj9LApaKxAH7RT0GKz7t2kVd8bD62LbhkFiVVlvqVTBBihHinzAx8wPSCaU2saAt4fYc  
+0w86it8IKBuwZyjE= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG   DNSKEY 5 2 3600 20110629030401 20110331030401 26437
```

```
vix.com.  
QvyoIbB1fTtge9aBTj88oBBFUnfLdGxGoyABG3bkPDAiDB5TUgJa68UDcF5k9c5fQEHZA6rd52QRxkPKy0hb5Reh64cZMjzMB  
ZxZaivxX+W+hmkEk9ztSgWaotNBw2RHechItBI4/IPZWRXGNPr1IIduI8KC+dm96tf404BraAU= ;{id =  
26437}
```

```
# drill -D vix.com @ns.sql1.vix.com. rrsig
```

```
;; QUESTION SECTION:
```

```
;; vix.com.      IN      RRSIG
```

```
;; ANSWER SECTION:
```

```
vix.com.      3600      IN      RRSIG      A 5 2 3600 20110629030401 20110331030401 63066 vix.com.
```

```
CsxzLHeqDLi2XXKqGALXyn4lbmZrqkDzCYegv6EiZQFpPHG8oVdxvqJDCczpVHF3mykB05uHntpyo0S4om34l8fkIuVKViE6c  
/3b+j3jiJIIfXbFYPqM501NChRf/SwkBqsmKRj4jbTp3jCicUG6M3LyNWe5B2CjVd9hEumzrbjY=
```

```
 ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG      MX 5 2 3600 20110629030401 20110331030401 63066 vix.com.
```

```
urjAd1NVJKNfUOI/l0aJRNEQJJfexjnwRTcyzcZmVvxvV5FlqLT904aIzcKMPnM2L3FWpf
```

```
+F0Tzfjr9Cb46pUHrj9LApaKxAH7RT0GKz7t2kVd8bD62LbhkFiVVlvqVTBBihHinzAx8wPSCaU2saAt4fYc
```

```
+0w86it8IKBuwZyjE= ;{id = 63066}
```

```
vix.com.      3600      IN      RRSIG      DNSKEY 5 2 3600 20110629030401 20110331030401 26437
```

```
vix.com.
```

```
QvyoIbB1fTtge9aBTj88oBBFUnfLdGxGoyABG3bkPDAiDB5TUgJa68UDcF5k9c5fQEHZA6rd52QRxkPKy0hb5Reh64cZMjzMB
```

```
ZxZaivxX+W+hmkEk9ztSgWaotNBw2RHechItBI4/IPZWRXGNPr1IIduI8KC+dm96tf404BraAU= ;{id =
```

```
26437}
```

NSEC - Next SECure Record

NSEC - to accommodate negative authentication requests

NSEC - indicates all zones for which the name server is authoritative

NSEC - assume the following zones:

alpha.org

charlie.org

delta.org

NSEC - A query for beta.org will yield:

alpha.org NSEC charlie.org

NSEC3 : RFC5155 : 2008-02

NSEC3 - like NSEC but to prevent tree walking

NSEC3 - instead of actual names, the
hash of the name is given

```
# drill -D rps.vix.com @ns.sql1.vix.com A
```

```
;; QUESTION SECTION:
```

```
;; rps.vix.com. IN A
```

```
;; ANSWER SECTION:
```

```
;; AUTHORITY SECTION:
```

```
vix.com. 3600 IN SOA ns.lah1.vix.com. hostmaster.vix.com. 2011033116 3600 1800  
604800 3600
```

```
vix.com. 3600 IN RRSIG SOA 5 2 3600 20110629030401 20110331030401 63066 vix.com.
```

```
fPpFeE/Y/1HfFtKTAjfWLBQafC2i4qf5gYewmr0fQHzH7xIYmvx
```

```
+rpenaKfr4By2R01Dh5q6kKgB3DR7G9swmAXcAVB5TzvQ6UcmjXcGGZPw+HUwUSIAAt6q559YMKxSN6DTeh7/
```

```
kNLLPtoPZqSmz7rxIr0USe2VwAYDznGtlzdQ= ;{id = 63066}
```

```
vix.com. 3600 IN NSEC ns-lah1._meta.vix.com. A NS SOA MX TXT AAAA RRSIG NSEC  
DNSKEY
```

```
vix.com. 3600 IN RRSIG NSEC 5 2 3600 20110629030401 20110331030401 63066
```

```
vix.com. Hm3dfubDRTtF8BrztQ3X2tCc5IJ7+J03cB8F5rQhDAyzBz7Xc0ESJrwyUCk8YL/
```

```
w3i360fUuhN3Mah0dTzrzoAMxWp90yM5MRbRSzUQwQ
```

```
+73cRbq2C2YefsYPatPiL9vHnc5Wvo9xtrFEjiWK7qcHgBw03SrXPsUYzn8seB8DtA= ;{id = 63066}
```

```
relay.vix.com. 3600 IN NSEC server99.vix.com. CNAME RRSIG NSEC
```

```
relay.vix.com. 3600 IN RRSIG NSEC 5 3 3600 20110629030401 20110331030401 63066
```

```
vix.com. SrJ3NLPntXcN+SpT9igyoEyQYznsomsbAxqfXutF5o0VDfaeHZvB2LZC7+HjCAwwH6F7YWItdRVFt4PVQ6/
```

```
ouZ2K7r2RLoaMyaaHAJzhq4EN503AoFD70NcMyVV4BIzI6vsquORSPW8H03ym/0kZ0aQG0Bw2UFW/Q6Pwx4xVbKA= ;{id =  
63066}
```

```
# drill -D rps.vix.com @ns.sql1.vix.com
```

```
;; QUESTION SECTION:
```

```
;; rps.vix.com. IN A
```

```
;; ANSWER SECTION:
```

```
;; AUTHORITY SECTION:
```

```
vix.com. 3600 IN SOA ns.lah1.vix.com. hostmaster.vix.com. 2011033116 3600 1800  
604800 3600
```

```
vix.com. 3600 IN RRSIG SOA 5 2 3600 20110629030401 20110331030401 63066 vix.com.  
fPpFeE/Y/1HfFtKTAjfWLBQafC2i4qf5gYewmr0fQHzH7xIYmvx
```

```
+rpenaKfr4By2R01Dh5q6kKgB3DR7G9swmAXcAVB5TzvQ6UcmjXcGGZPw+HUwUSIAAt6q559YMKxSN6DTeh7/  
kNLLPtoPZqSmz7rxIr0USe2VwAYDznGtlzdQ= ;{id = 63066}
```

```
vix.com. 3600 IN NSEC ns-lah1._meta.vix.com. A NS SOA MX TXT AAAA RRSIG NSEC  
DNSKEY
```

```
vix.com. 3600 IN RRSIG NSEC 5 2 3600 20110629030401 20110331030401 63066
```

```
vix.com. Hm3dfubDRTtF8BrztQ3X2tCc5IJ7+J03cB8F5rQhDAyzBz7Xc0ESJrwyUCk8YL/  
w3i360fUuhN3Mah0dTzrzoAMxWp90yM5MRbRSzUQwQ
```

```
+73cRbq2C2YefsYPatPiL9vHnc5Wvo9xtrFEjiWK7qcHgBw03SrXPsUYzn8seB8DtA= ;{id = 63066}
```

```
relay.vix.com. 3600 IN NSEC server99.vix.com. CNAME RRSIG NSEC
```

```
relay.vix.com. 3600 IN RRSIG NSEC 5 3 3600 20110629030401 20110331030401 63066
```

```
vix.com. SrJ3NLPntXcN+SpT9igyoEyQYznsomsbAxqfXutF5o0VDfaeHZvB2LZC7+HjCAwwH6F7YWItdRVFt4PVQ6/  
ouZ2K7r2RLoaMyaaHAJzhq4EN503AoFD70NcMyVV4BIzI6vsquORSPW8H03ym/0kZ0aQG0Bw2UFW/Q6Pwx4xVbKA= ;{id =  
63066}
```



```
# drill -D rps.vix.com @ns.sql1.vix.com
```

```
;; QUESTION SECTION:
```

```
;; rps.vix.com. IN A
```

```
;; ANSWER SECTION:
```

```
;; AUTHORITY SECTION:
```

```
vix.com. 3600 IN SOA ns.lah1.vix.com. hostmaster.vix.com. 2011033116 3600 1800  
604800 3600
```

```
vix.com. 3600 IN RRSIG SOA 5 2 3600 20110629030401 20110331030401 63066 vix.com.
```

```
fPpFeE/Y/1HfFtKTAjfWLBQafC2i4qf5gYewmr0fQHzH7xIYmvx
```

```
+rpenaKfr4By2R01Dh5q6kKgB3DR7G9swmAXcAVB5TzvQ6UcmjXcGGZPw+HUwUSIAt6q559YMKxSN6DTeh7/
```

```
kNLLPtoPZqSmz7rxIr0USe2VwAYDznGtlzdQ= ;{id = 63066}
```

```
vix.com. 3600 IN NSEC ns-lah1._meta.vix.com. A NS SOA MX TXT AAAA RRSIG NSEC  
DNSKEY
```

```
vix.com. 3600 IN RRSIG NSEC 5 2 3600 20110629030401 20110331030401 63066
```

```
vix.com. Hm3dfubDRTtF8BrztQ3X2tCc5IJ7+J03cB8F5rQhDAyzBz7Xc0ESJrwyUCk8YL/
```

```
w3i360fUuhN3Mah0dTzrzoAMxWp90yM5MRbRSzUQwQ
```

```
+73cRbq2C2YefsYPatPiL9vHnc5Wvo9xtrFEjiWK7qcHgBw03SrXPsUYzn8seB8DtA= ;{id = 63066}
```

```
relay.vix.com. 3600 IN NSEC server99.vix.com. CNAME RRSIG NSEC
```

```
relay.vix.com. 3600 IN RRSIG NSEC 5 3 3600 20110629030401 20110331030401 63066
```

```
vix.com. SrJ3NLPntXcN+SpT9igyoEyQYznsomsbAxqfXutF5o0VDfaeHZvB2LZC7+HjCAwwH6F7YWItdRVFt4PVQ6/
```

```
ouZ2K7r2RLoaMyaaHAJzhq4EN503AoFD70NcMyVV4BIzI6vsquORSPW8H03ym/0kZ0aQG0Bw2UFW/Q6Pwx4xVbKA=
```

```
;{id = 63066}
```

```
# drill -D rps.vix.com @ns.sql1.vix.com
;; QUESTION SECTION:
;; rps.vix.com. IN      A

;; ANSWER SECTION:

;; AUTHORITY SECTION:
vix.com.      3600    IN      SOA      ns.lah1.vix.com. hostmaster.vix.com. 2011033116 3600 1800
604800 3600
vix.com.      3600    IN      RRSIG    SOA 5 2 3600 20110629030401 20110331030401 63066 vix.com.
fPpFeE/Y/1HfFtKTAjfWLBQafC2i4qf5gYewmr0fQHzH7xIYmvx
+rpenaKfr4By2R01Dh5q6kKgB3DR7G9swmAXcAVB5TzvQ6UcmjXcGGZPw+HUwUSIAAt6q559YMKxSN6DTeh7/
kNLLPtoPZqSmz7rxIr0USe2VwAYDznGtlzdQ= ;{id = 63066}
vix.com.      3600    IN      NSEC     ns-lah1._meta.vix.com. A NS SOA MX TXT AAAA RRSIG NSEC
DNSKEY
vix.com.      3600    IN      RRSIG    NSEC 5 2 3600 20110629030401 20110331030401 63066
vix.com. Hm3dfubDRTtF8BrztQ3X2tCc5IJ7+J03cB8F5rQhDAyzBz7Xc0ESJrwyUCk8YL/
w3i360fUuhN3Mah0dTzrzoAMxWp90yM5MRbRSzUQwQ
+73cRbq2C2YEfsYPatPiL9vHnc5Wvo9xtrFEjiWK7qcHgBw03SrXPsUYzn8seB8DtA= ;{id = 63066}
relay.vix.com. 3600    IN      NSEC     server99.vix.com. CNAME RRSIG NSEC
relay.vix.com. 3600    IN      RRSIG    NSEC 5 3 3600 20110629030401 20110331030401 63066
vix.com. SrJ3NLPntXcN+SpT9igyoEyQYznsombsAxqfXutF5o0VDfaeHZvB2LZC7+HjCAwwH6F7YWItdRVFt4PVQ6/
ouZ2K7r2RLoaMyaaHAJzhq4EN503AoFD70NcMyVV4BIzI6vsquORSPW8H03ym/0kZ0aQG0Bw2UFW/Q6Pwx4xVbKA= ;{id =
63066}
```

DS - Delegation Signer

DS:

A key signed by the parent zone to indicate stuff down in the hierarchy can be trusted

DS:

The parent zone stores the key tag, algorithm and a digest/hash of the `DNSKEY` in the child zone.

DS:

The parent zone then signs the DS record and creates a corresponding RRSIG record

```
# drill -s -D vix.com DNSKEY
```

```
;; ANSWER SECTION:
```

```
vix.com.          3600      IN        DNSKEY   257 3 5  
AwEAAbKW5zsYMBUX4MS0yq3MNM4312c7WEF1Af2Iy20/A+U+h7F3Etb1BDJVs/  
LgtdjsE3JHak51iRaELL0oEvVeRIIa1UjNvXIei+QV1n1Sas8LcXya0X0YA2Jfxyz0p  
EWArN1QLhkgVDPAsEwKLzYfvjW78CF10ZnYxbBWXwKgb4z  
;{id = 26437 (ksk), size = 1024b}
```

```
vix.com.          3600      IN        DNSKEY   256 3 5  
BEAAAA06wBt1U39U8meHca3JBCWixBi8BvZLMJZp51/5vViM2+fh93XF1SqJaAaq  
gX6PszTPULE1vuTV2xTV4uQjUTaFv8qDnsjbfXVusE1v+0aQpSVuP8GjI28cGi9P  
gA0cz2ACdiD2XVbYKUDTJb+pqoE/o3Z6FjKf6ByTkJUI5x9Dlw==  
;{id = 63066 (zsk), size = 1024b}
```

```
; vix.com.        3600      IN        DS        26437 5 1  
483cca94fd7e2aa30f4fca34ccf0db4ddc601388  
; xidaf-sidan-gazol-vupap-fofag-zudif-gufyz-bikag-talyk-begom-muxox
```

```
; vix.com.        3600      IN        DS        63066 5 1  
8229b0484396f12015de1cb7e3267ed1f109060a  
; xobid-nusog-mebon-kusid-bihyt-velyr-lomid-kizut-ceseb-nacab-puxux
```

} vix.com.

} com.

```
# drill -s -D vix.com DNSKEY
;; ANSWER SECTION:
```

```
vix.com.          3600      IN         DNSKEY    257 3 5
AwEAAbKW5zsYMBUX4MS0yq3MNM4312c7WEF1Af2Iy20/A+U+h7F3Etb1BDJVs/
LgtdjsE3JHak51iRaELL0oEvVeRIIa1UjNvXIei+QV1n1Sas8LcXya0X0YA2Jfxyz0p
EWArN1QLhkgVDPAsEwKLzYfVjW78CF10ZnYxbBWXwKgb4z
;{id = 26437 (ksk), size = 1024b}
```

```
vix.com.          3600      IN         DNSKEY    256 3 5
BEAAAA06wBt1U39U8meHca3JBCWixBi8BvZLMJZp51/5vViM2+fh93XF1SqJaAaq
gX6PszTPULE1vuTV2xTV4uQjUTaFv8qDnsjbfXVusE1v+0aQpSVuP8GjI28cGi9P
gA0cz2ACdiD2XVbYKUDTJb+pqoE/o3Z6FjKf6ByTkJUI5x9Dlw==
;{id = 63066 (zsk), size = 1024b}
```

```
; vix.com.        3600      IN         DS         26437 5 1
483cca94fd7e2aa30f4fca34ccf0db4ddc601388
; xidaf-sidan-gazol-vupap-fofag-zudif-gufyz-bikag-talyk-begom-muxox
```

```
; vix.com.        3600      IN         DS         63066 5 1
8229b0484396f12015de1cb7e3267ed1f109060a
; xobid-nusog-mebon-kusid-bihyt-velyr-lomid-kizut-ceseb-nacab-puxux
```



```
# drill -s -D vix.com DNSKEY
```

```
;; ANSWER SECTION:
```

```
vix.com.          3600      IN        DNSKEY   257 3 5  
AwEAAbKW5zsYMBUX4MS0yq3MNM4312c7WEF1Af2Iy20/A+U+h7F3Etb1BDJVs/  
LgtdjsE3JHak51iRaELL0oEvVeRIIa1UjNvXIei+QV1n1Sas8LcXya0X0YA2Jfzez0p  
EWArN1QLhkgVDPAsEwKLzYfvjW78CF10ZnYxbBWXwKgb4z  
;{id = 26437 (ksk), size = 1024b}
```

```
vix.com.          3600      IN        DNSKEY   256 3 5  
BEAAAA06wBt1U39U8meHca3JBCWixBi8BvZLMJZp51/5vViM2+fh93XF1SqJaAaq  
gX6PszTPULE1vuTV2xTV4uQjUTaFv8qDnsjbfXVusE1v+0aQpSVuP8GjI28cGi9P  
gA0cz2ACdiD2XVbYKUDTJb+pqoE/o3Z6FjKf6ByTkJUI5x9Dlw==  
;{id = 63066 (zsk), size = 1024b}
```

```
; vix.com.        3600      IN        DS        26437 5 1  
483cca94fd7e2aa30f4fca34ccf0db4ddc601388  
; xidaf-sidan-gazol-vupap-fofag-zudif-gufyz-bikag-talyk-begom-muxox
```

```
; vix.com.        3600      IN        DS        63066 5 1  
8229b0484396f12015de1cb7e3267ed1f109060a  
; xobid-nusog-mebon-kusid-bihyt-velyr-lomid-kizut-ceseb-nacab-puxux
```

```
# drill -s -D vix.com DNSKEY
```

```
;; ANSWER SECTION:
```

```
vix.com.          3600      IN        DNSKEY   257 3 5  
AwEAAbKW5zsYMBUX4MS0yq3MNM4312c7WEF1Af2Iy20/A+U+h7F3Etb1BDJVs/  
LgtdjsE3JHak51iRaELL0oEvVeRIIa1UjNvXIei+QV1n1Sas8LcXya0X0YA2Jfzez0p  
EWArN1QLhkgVDPAsEwKLzYfvjW78CF10ZnYxbBWXwKgb4z  
;{id = 26437 (ksk), size = 1024b}
```

```
vix.com.          3600      IN        DNSKEY   256 3 5  
BEAAAA06wBt1U39U8meHca3JBCWixBi8BvZLMJZp51/5vViM2+fh93XF1SqJaAaq  
gX6PszTPULE1vuTV2xTV4uQjUTaFv8qDnsjbfXVusE1v+0aQpSVuP8GjI28cGi9P  
gA0cz2ACdiD2XVbYKUDTJb+pqoE/o3Z6FjKf6ByTkJUI5x9Dlw==  
;{id = 63066 (zsk), size = 1024b}
```

```
; vix.com.        3600      IN        DS        26437 5 1  
483cca94fd7e2aa30f4fca34ccf0db4ddc601388  
; xidaf-sidan-gazol-vupap-fofag-zudif-gufyz-bikag-talyk-begom-muxox
```

```
; vix.com.        3600      IN        DS        63066 5 1  
8229b0484396f12015de1cb7e3267ed1f109060a  
; xobid-nusog-mebon-kusid-bihyt-velyr-lomid-kizut-ceseb-nacab-puxux
```

.CA Recommended DNSSEC Key Parameters

Parameter	Value
Key Signing Key (KSK)	2048 RSA
KSK Rollover Schedule	once per year
KSK Algorithm	RSA/SHA/256
Zone Signing Key (ZSK)	1024-bit RSA
ZSK Rollover Schedule	once per month
ZSK Signature Algorithm	RSA/SHA/256
Authenticated Proof of Non Existence	NSEC3 with opt-out

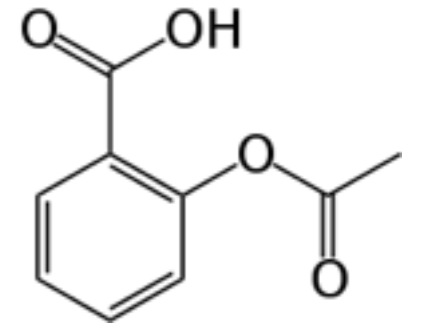
DNSSEC - Traversal

Caching Name Server



DNSSEC aware

RFC 4035



client (aka)
stub resolver

recursive
name
server



caching

resolver



authority

DNSSEC - Bootstrapping

Assumption - Root Zone has been signed

Assumption - Public Keys are available

FACT - Root keys available since: 2010-07-15

DNSSEC - Fetch and Validate Root Keys

This is a manual process.

Guide: <http://data.iana.org/root-anchors/draft-icann-dnssec-trust-anchor.txt>

Complete trust anchor: <https://data.iana.org/root-anchors/root-anchors.xml>

PGP signature: <https://data.iana.org/root-anchors/root-anchors.asc>

DNSSEC - Traversal

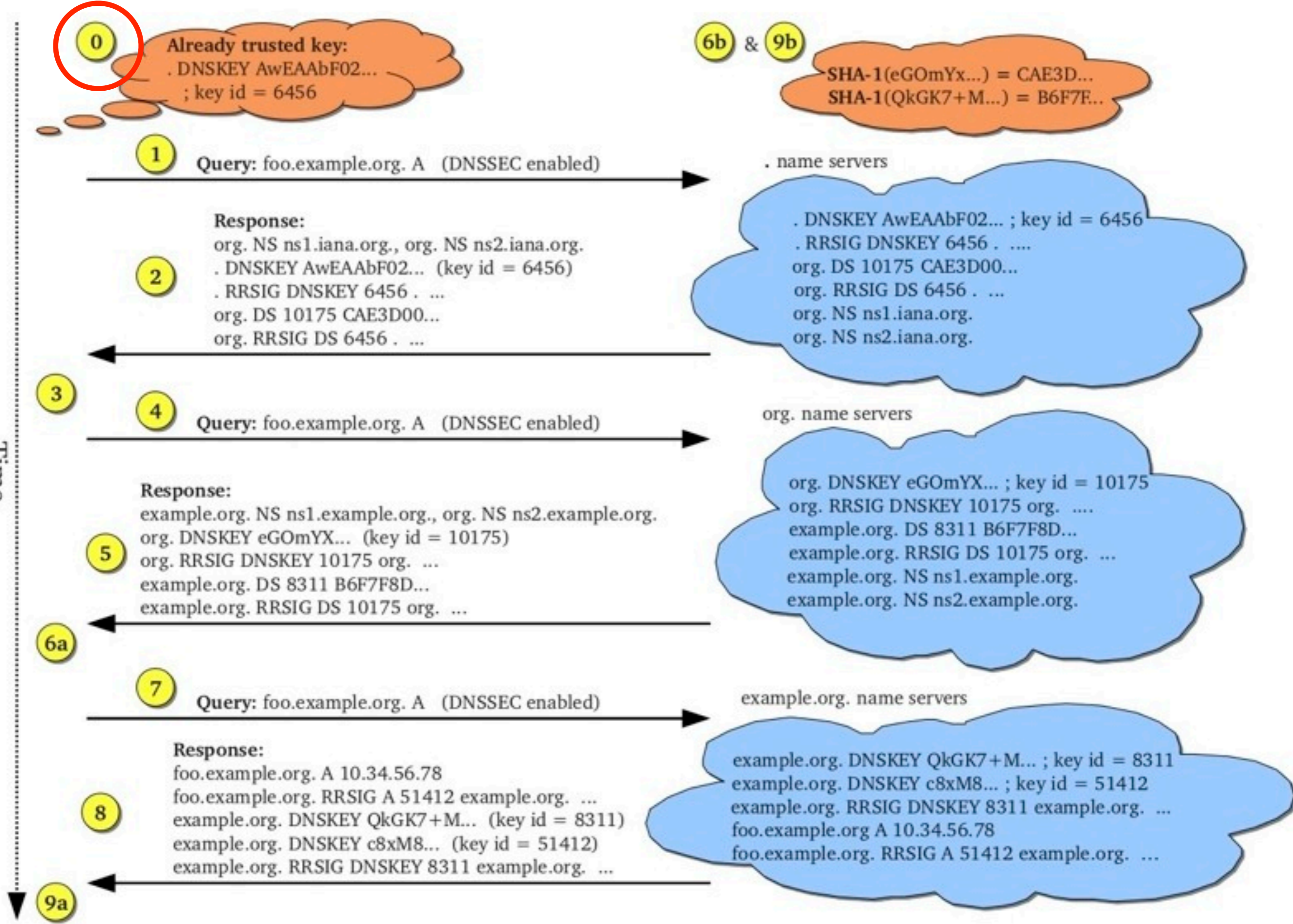


Figure 3.4: Graphical representation of a DNSSEC traversal

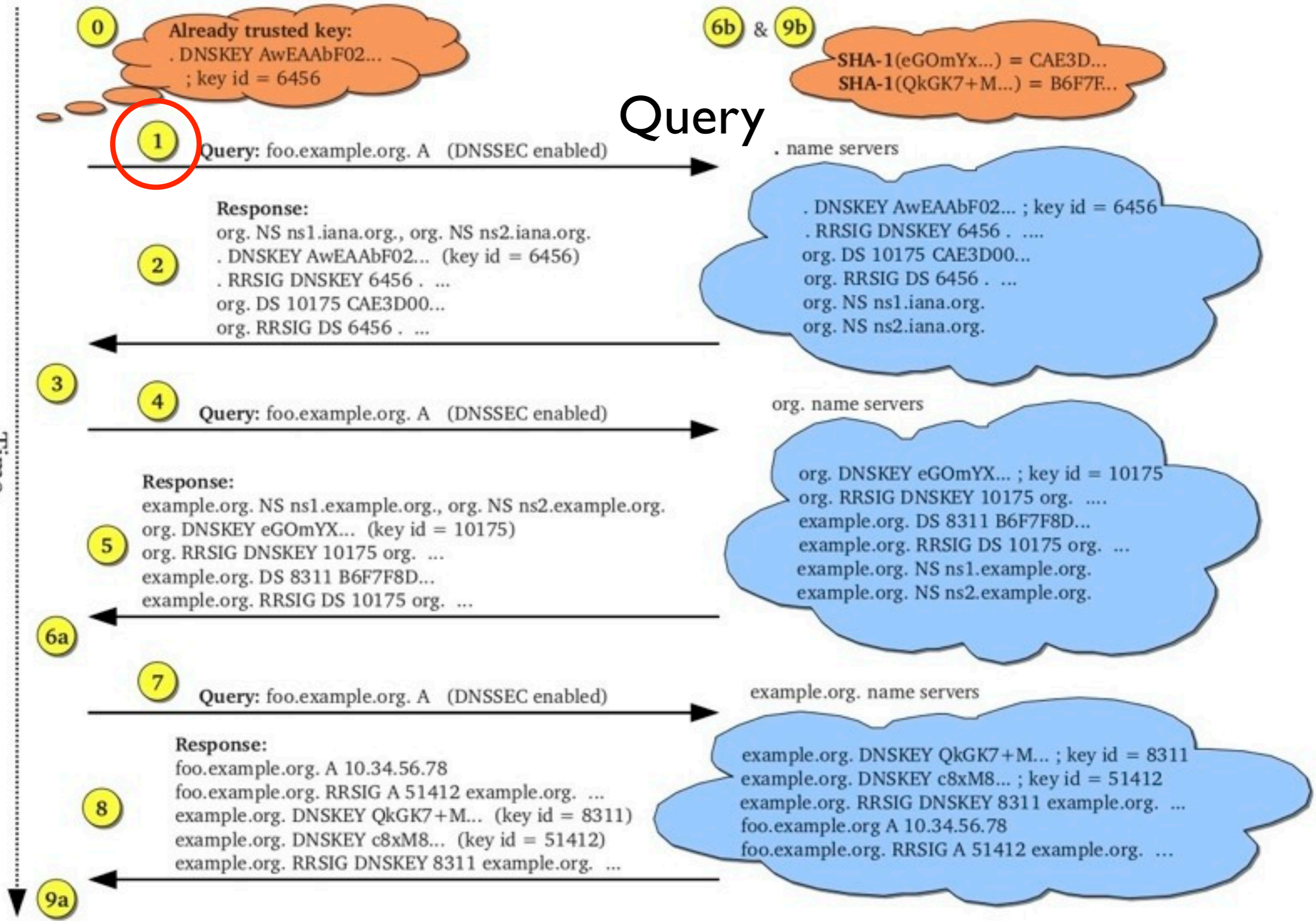


Figure 3.4: Graphical representation of a DNSSEC traversal

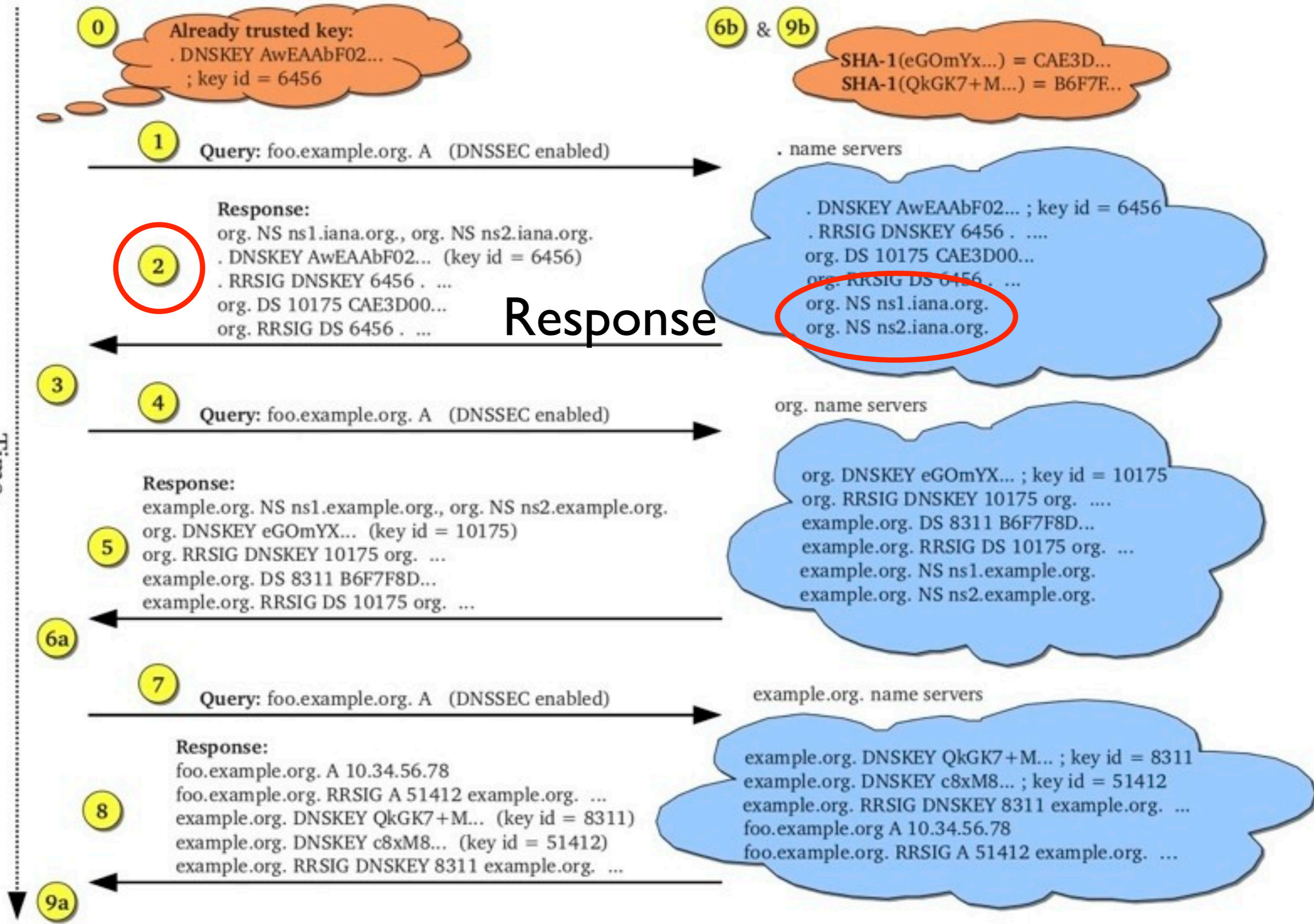


Figure 3.4: Graphical representation of a DNSSEC traversal

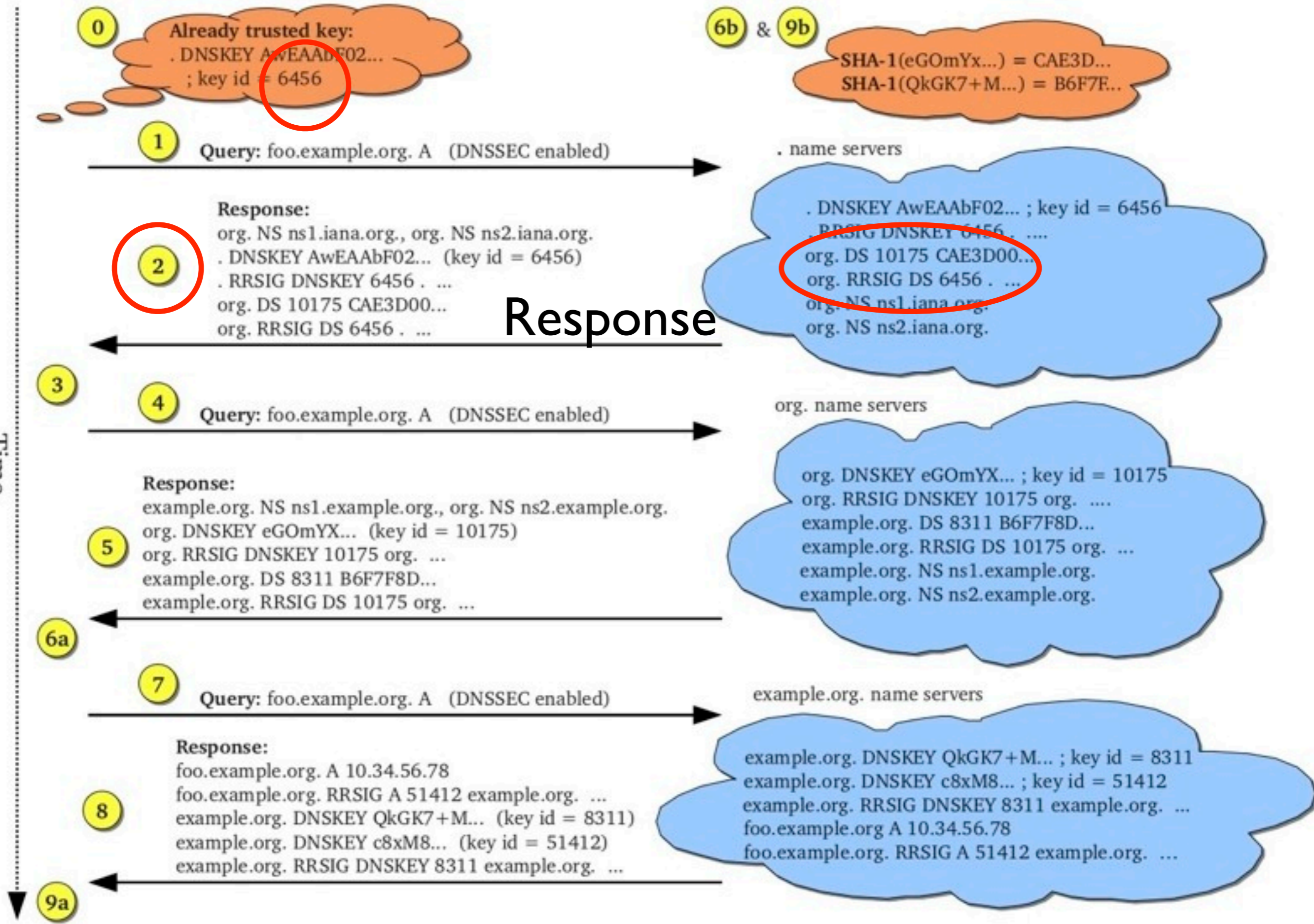


Figure 3.4: Graphical representation of a DNSSEC traversal

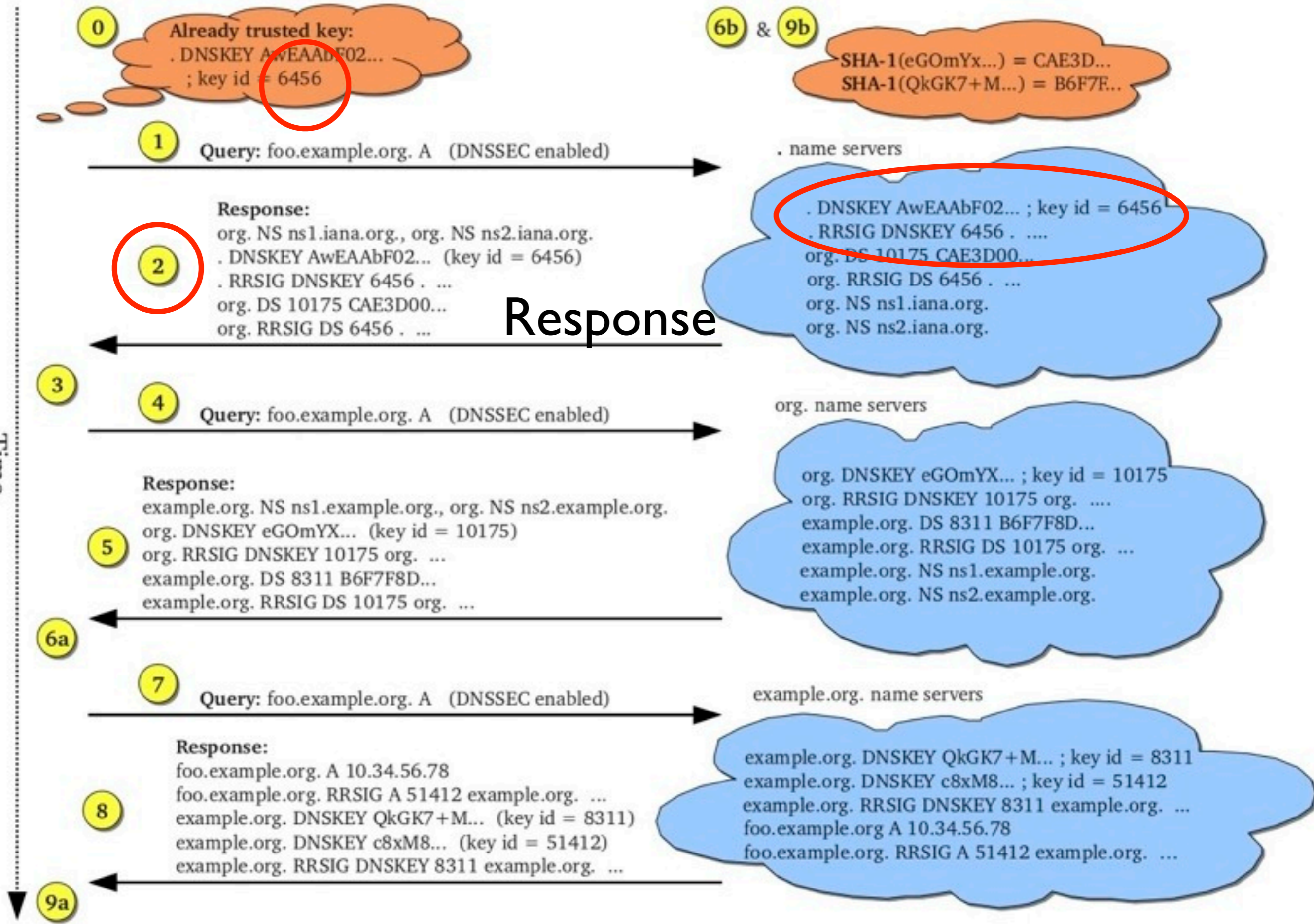


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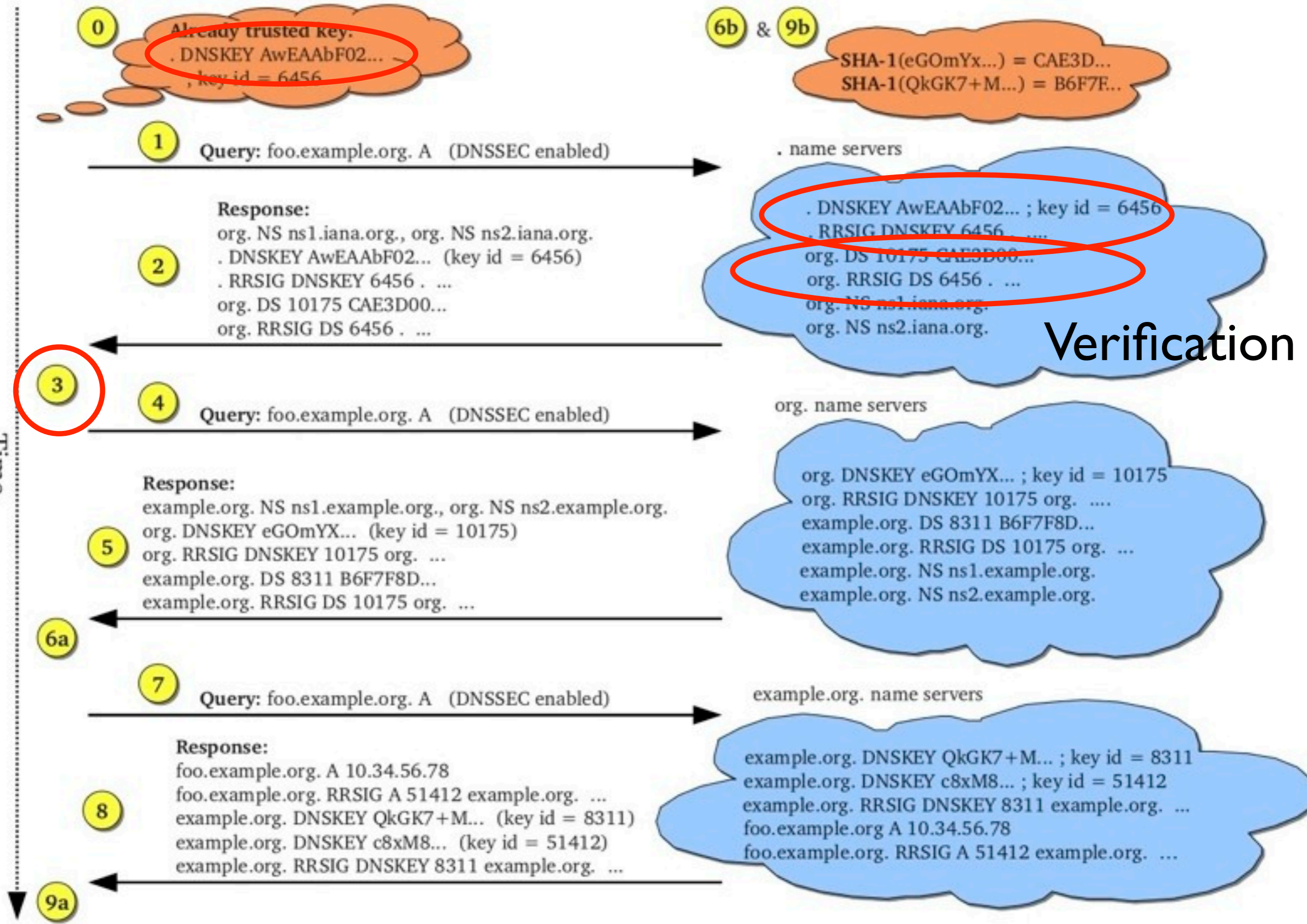


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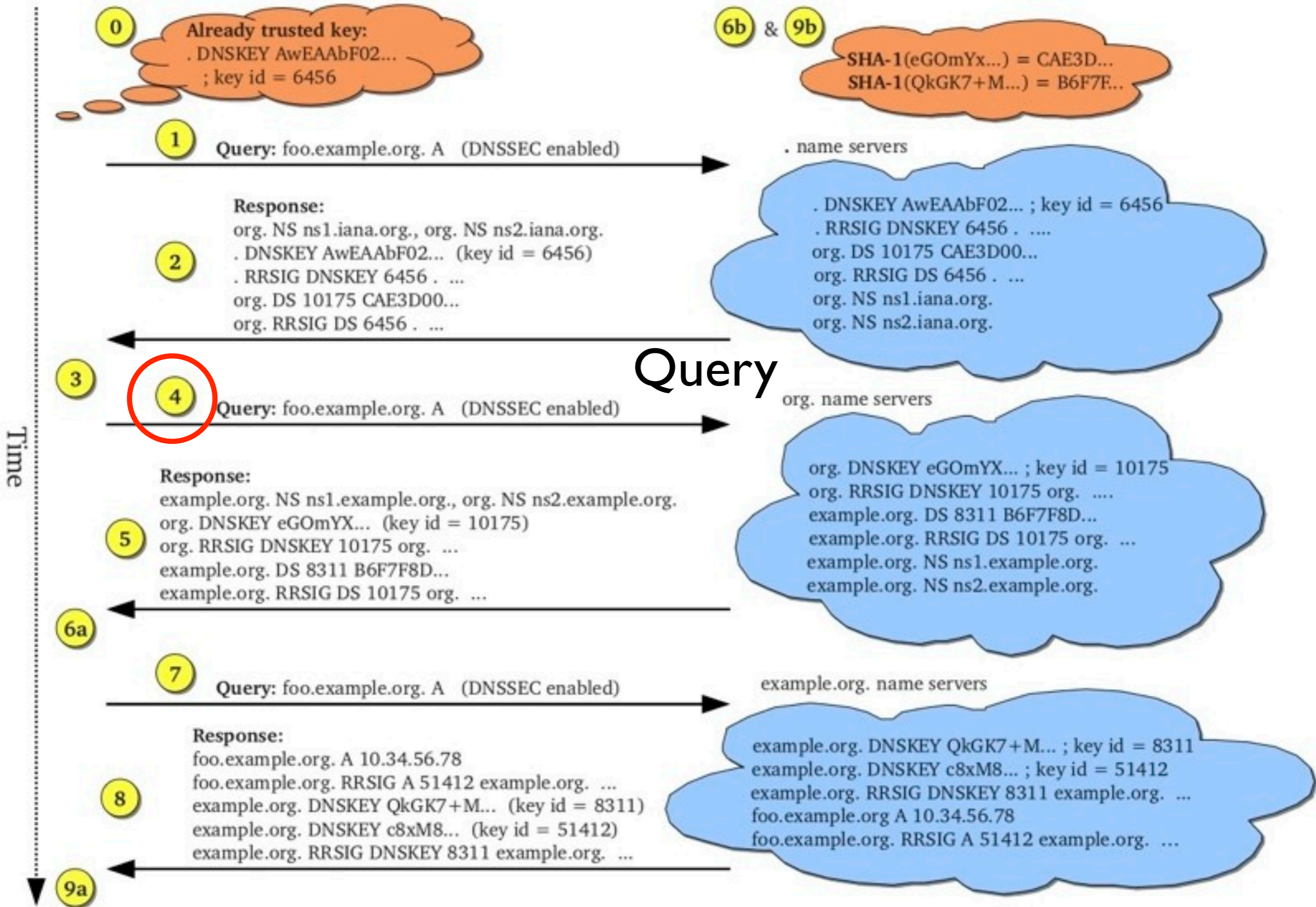


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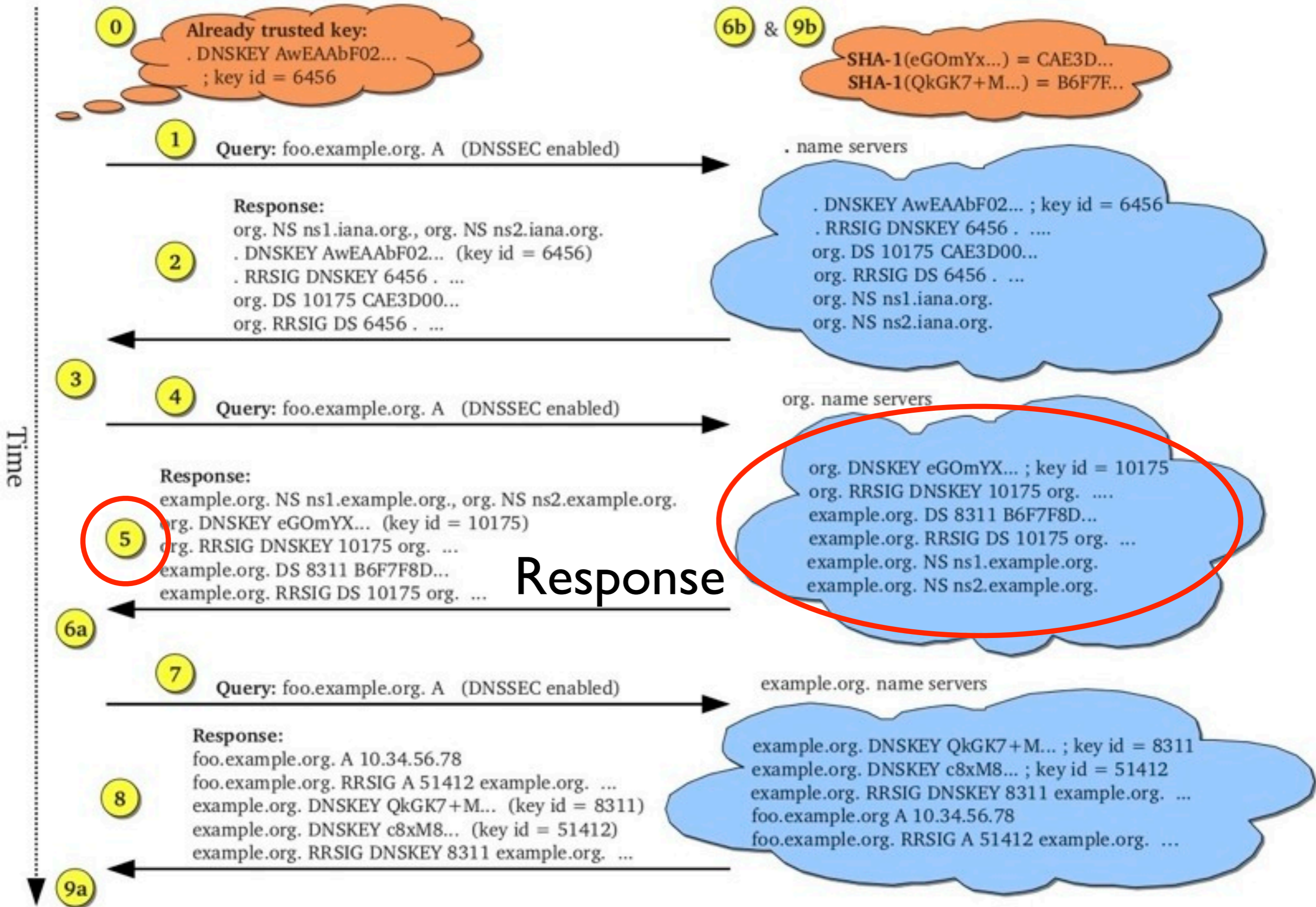


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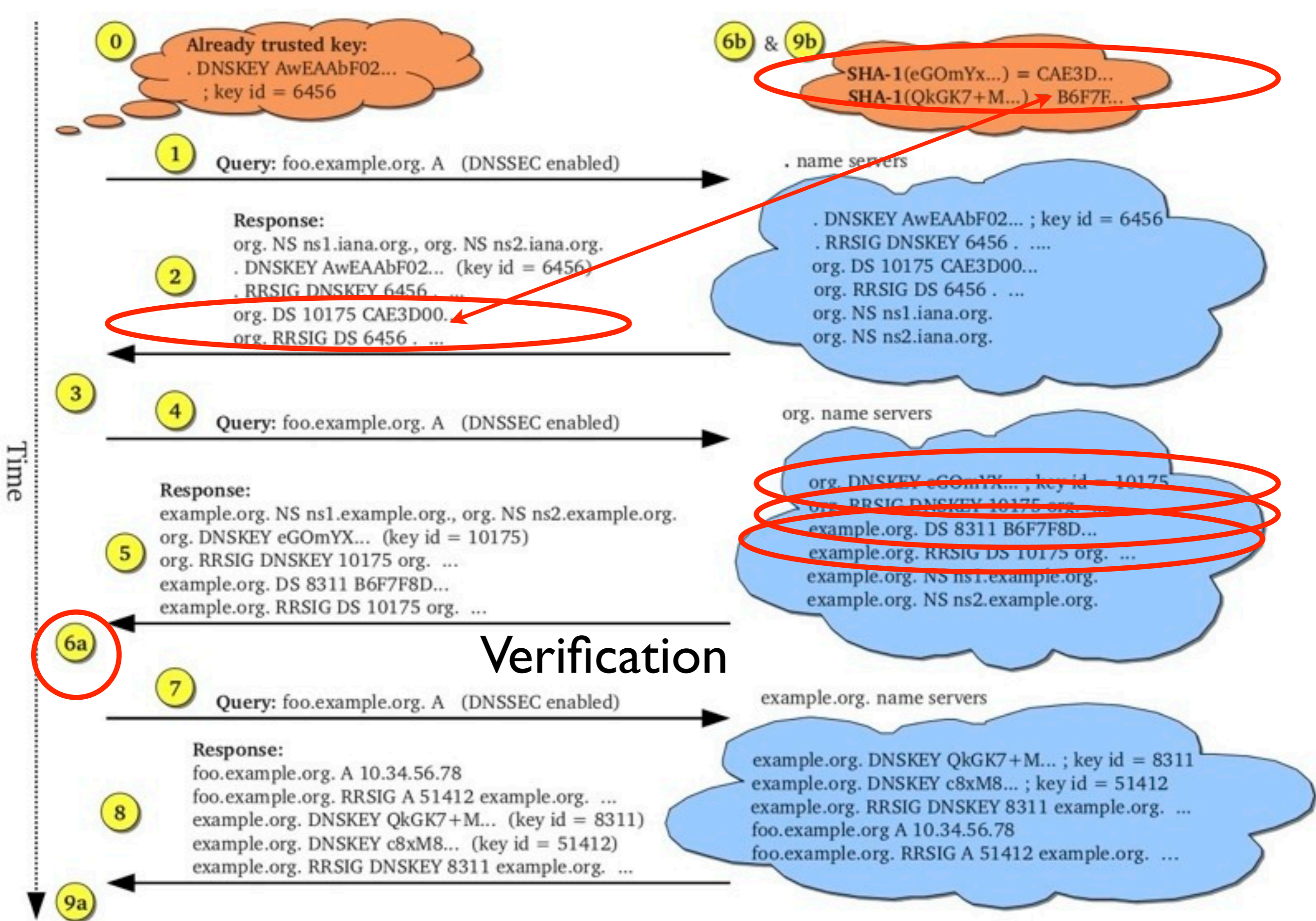


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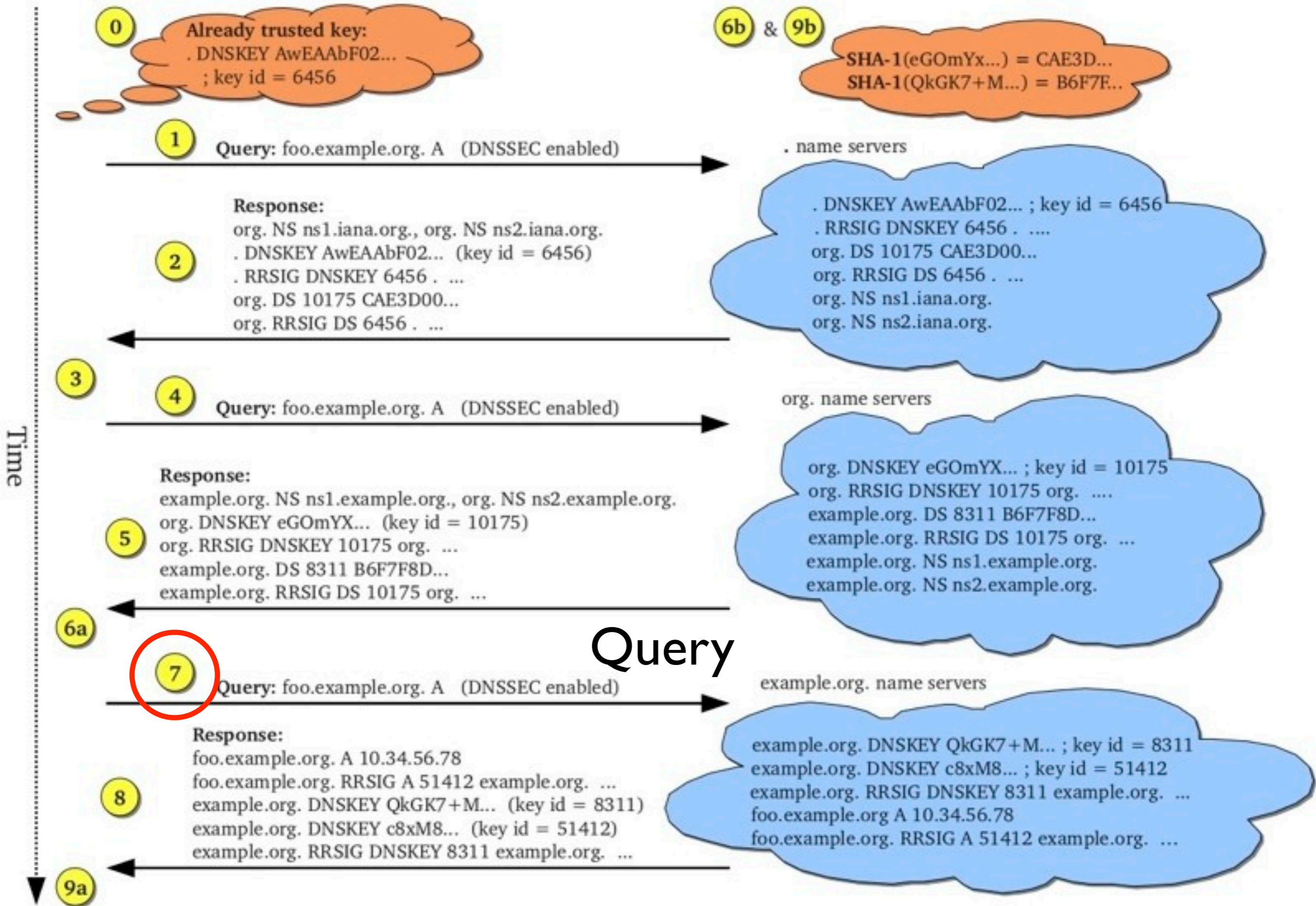


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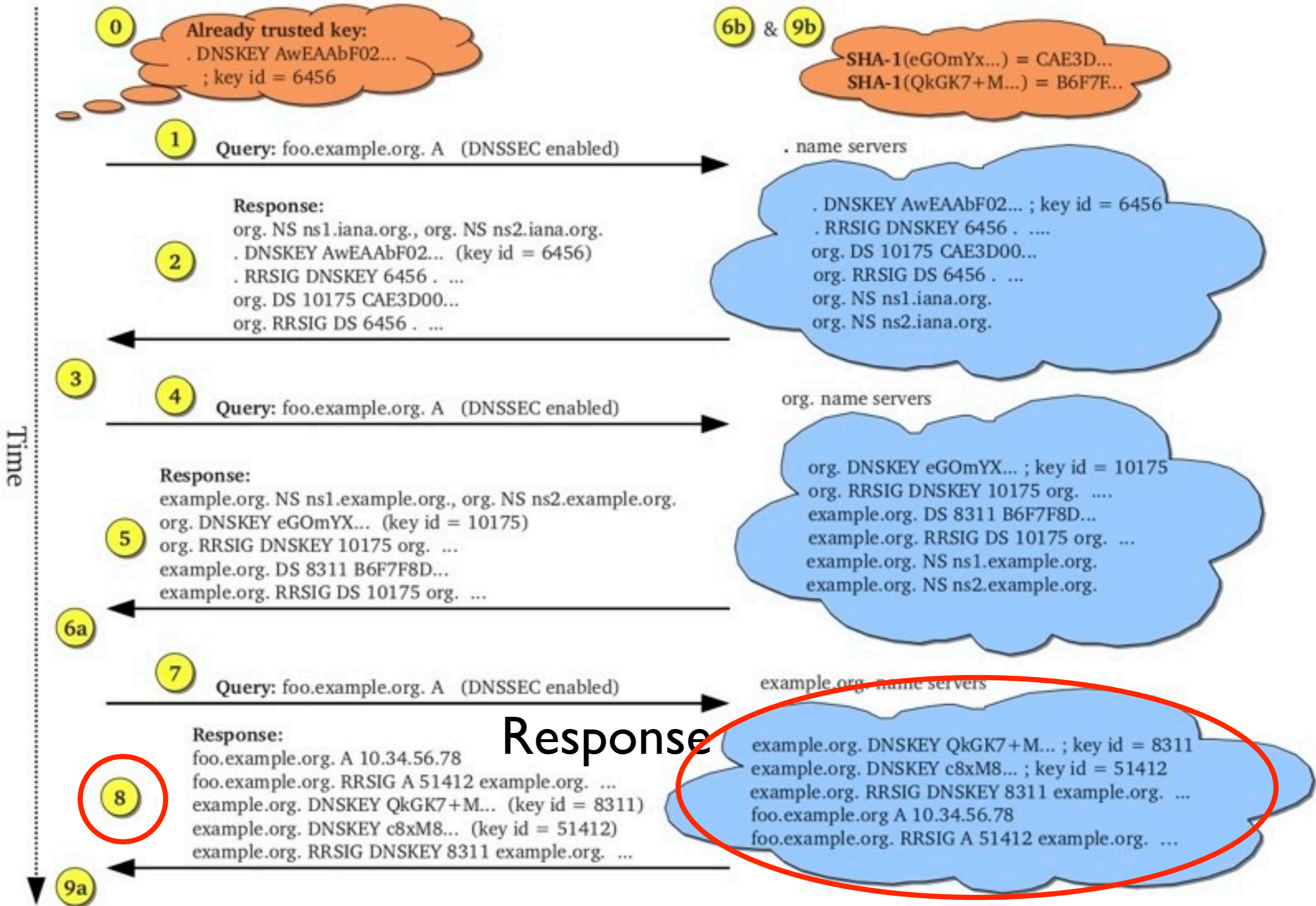


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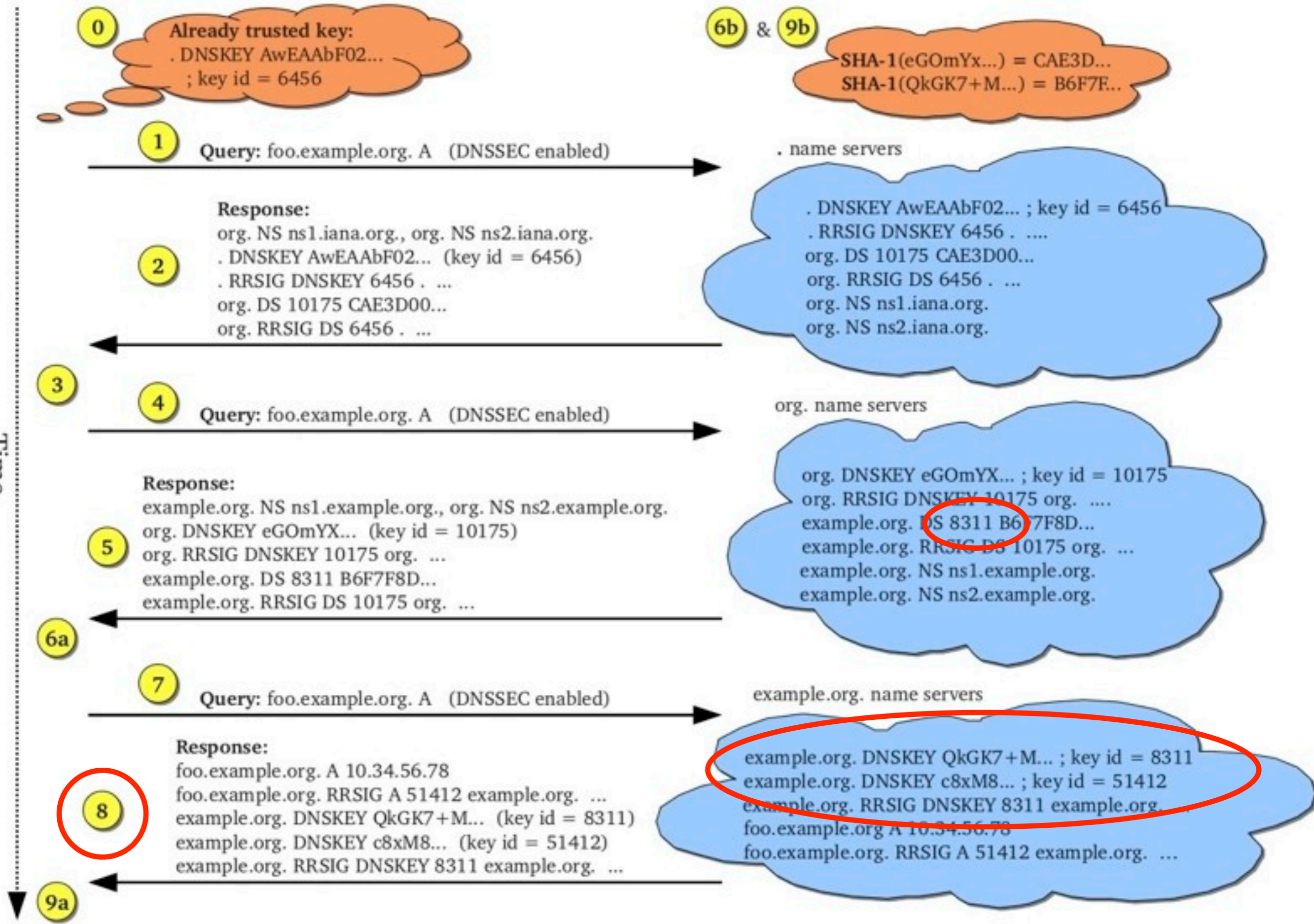


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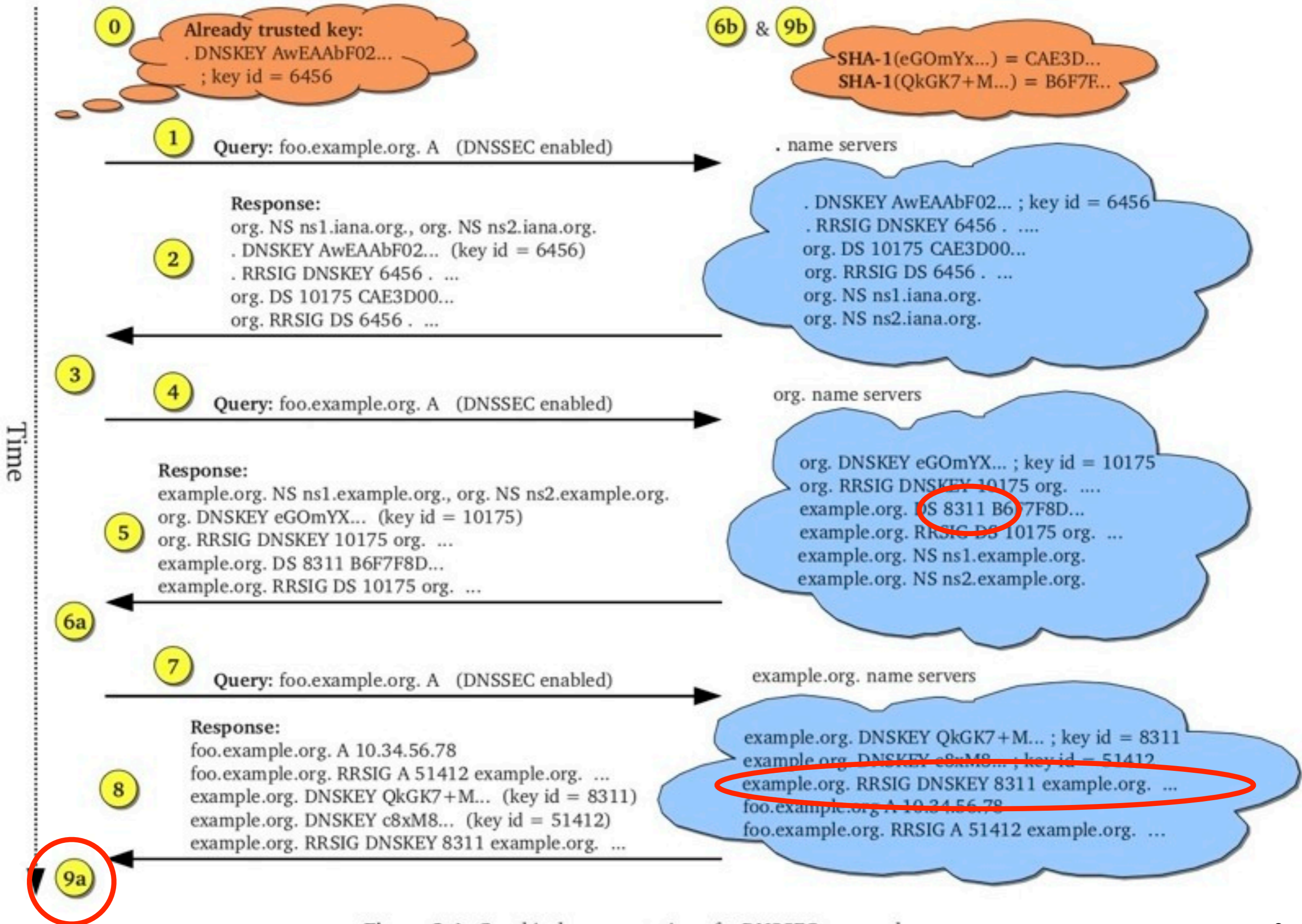


Figure 3.4: Graphical representation of a DNSSEC traversal

DNSSEC - Management

<1>

Generate a Key Signing Pair*

* KSK: Larger is better as the life time may be long

<2>

Create a MD of the KSK for the Boss*

* Whoever administers the zone above in the hierarchy

<3>

Generate a Zone Signing Pair*

* ZSK is authenticated by signing with the KSK.

<4>

Create the NSEC/NSEC3 RRs

<5>

Create RRSIGs for all RRs using the ZSK

<6>

Recreate RRSIGs when editing RRs or upon expiry *

* NTP is important for proper DNSSEC

<7>

Recreate ZSK/KSK from time to time and resign

<8>

time, disk, bandwidth, memory requirements

time, disk, bandwidth, memory requirements

hours, 4-12x, 90% -> 400%, 10-200%

DNSSEC - Security

3 goals of good computer security



Confidentiality



Integrity



Availability



Confidentiality?



FAIL



Integrity?



Mostly



WINNING



DNSSEC - Prevents Active MiM attacks *

* given all links in the DNS query-response are DNSSEC aware

DNSSEC - Allows for replay attacks.

DNSSEC - Crypto brute force attacks

DNSSEC - How safe is RSA/SHA-1?

Availability



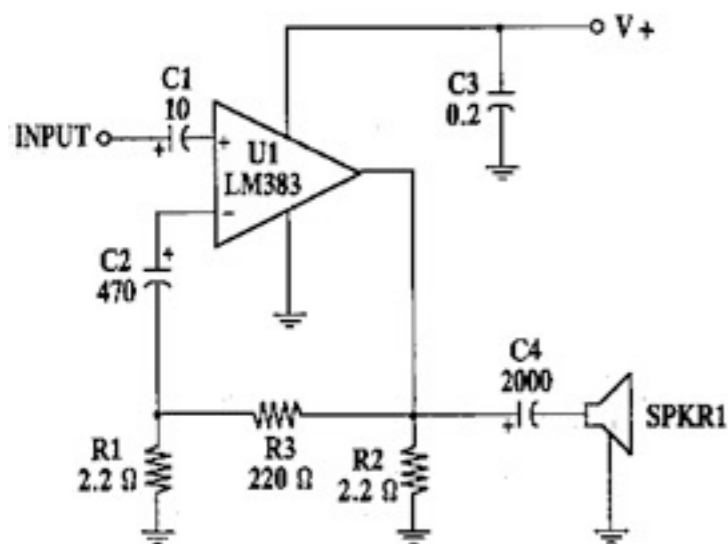
Potential and Risk of



DNSSEC - Creates other types of attacks

DNSSEC - Enhances DoS/Availability attacks

DNSSEC - Amplification Attack



```
# drill -D se. @a.ns.se. DNSKEY
```

```
;; ->>HEADER<<- opcode: QUERY, rcode: NOERROR, id: 37009  
;; flags: qr aa rd ; QUERY: 1, ANSWER: 6, AUTHORITY: 0, ADDITIONAL: 0  
;; QUESTION SECTION:  
;; se. IN DNSKEY
```

```
;; ANSWER SECTION:
```

```
se. 3600 IN DNSKEY 256 3 5 AwEAAYUDNvoT6yfrNra/7d2c7ZoSBphCxjs5xrpPbAPi8F6aP/  
oC2W9xPaXa5BXcEhneMwiabUBErF4LWFkSher5z2mrPN/3/YH/78IGwMMWV2wxOAtzbLkhuXWTh7cjH2u5sR8xvVeAQgAV0fGCE/ykvj6bF2pvj1r/  
5KRY0izG7PAZ ;{id = 57240 (zsk), size = 1024b}  
se. 3600 IN DNSKEY 256 3 5 AwEAAdZgee2x1z9yHDWafJ3oqAnwoU/v/awp0bk6lCnxtoZ6ukq6+0xgY0dusS7qCHF  
+LcBRCsAehpQJAWzL7c4xRrs2PT4/z4jZMtfA1EX6hN+s0ZXjxYwR7WdqVje4/Jtn2krpUvE+jjIyegQ+DKFkxawGJ5pG3EgU3B0MSvEonMT ;{id =  
12973 (zsk), size = 1024b}  
se. 3600 IN DNSKEY 257 3 5 AwEAAbaxTum9L7z1DmPiXPk0QZ2/qUM3to210Caey/ycZuvQ8Mh/  
dgGpwBmyZB9xZSkaCLa2Mw6pmDLrjk9hW0ffq5PXRvm9RrcA/eIEBEvbQzkY5sFkWAczNAs580scxi+/  
Gd5KfuVi3lJpYgJwwa2JB4doZ00IXywcCn0VTz0Hsl/lqpA2Bqj+e  
+ATzA5hWyiNyHPjiYvyMckSXTiGgFVVuG8H3N6Us8uSABu02UoFQeQi6YikIiCbf1FfCzr4vBIRXW6MaDs8kqAAadKjLk3i39dviL/  
YeyGUvq9Dan9PsvkwQejKN/7J0yCr2nYXfwGGCHkcBkkagv79EaRlZigUCp8= ;{id = 39547 (ksk), size = 2048b}  
se. 3600 IN DNSKEY 257 3 5 AwEAAb6IEZ2ETrgngbj0NAC10b4dRs/jD0MYPcMXRzQlo/eqo5AHXvqPaav+rgA3q  
+I6zvWYFTMUPxNT2wdJwV4R7VbXb3pBfYPBzeacqPaWSbw4W1BFdYy0WKe0sw3gvwD62dLGbykQAqx5gUYZ8gBtFXDsJe/x+JvenC/  
wmz7yW6mxpn3Tzd60vE6wjXhnBs62905xck0BskVx6dI  
+dMLoXNG2p5tpXfT4dGrA10SFwVb2C9QTaFww9fP60QqVoYz1xU1Z5BXa5ZB108I4rHYGtDYU36n0UhCG4nWnTJgUbNRsN3CeeTp1kZ94JS8jMsda0x983  
VIn5stGU1W4juyDqF0= ;{id = 7649 (ksk), size = 2048b}  
se. 3600 IN RRSIG DNSKEY 5 1 3600 20110409220345 20110403210548 7649 se. U8IIYExbbCcxYeTfQQGB/  
jEYuKnVf1G2c8bojkt3U7fNx7l7Z3IudEuLxATR4+xw3aKmGdfioC6EXnt5Umc0oUNxyf6t4zhEMmV9/LDXxULASDIBmk/e5RWTCFeiY  
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7649}  
se. 3600 IN RRSIG DNSKEY 5 1 3600 20110410160345 20110404150548 39547 se.  
E3uA3bUS0hBNTWoqARj6fSrFdxvaGDjsQRipT5Em+HUX4N09TR2/02tweeA8QKII3bKBRfZ1r56b1LK9nq0eLv3UhPgEbwHwmdpC9fHbRi9FCX  
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JsMeKLIbgEvQIsu6zpaZvc1l8ng== ;{id = 39547}
```



```
# drill -D se. @a.ns.se. DNSKEY
;; ->>HEADER<<- opcode: QUERY, rcode: NOERROR, id: 37009
;; flags: qr aa rd ; QUERY: 1, ANSWER: 6, AUTHORITY: 0, ADDITIONAL: 0
;; QUESTION SECTION:
;; se. IN      DNSKEY

;; ANSWER SECTION:
se.      3600    IN      DNSKEY  256 3 5 AwEAAYUDNvoT6yfrNra/7d2c7ZoSBphCxjs5xrpPbAPi8F6aP/
oC2W9xPaXa5BXcEhneMwiabUBErF4LWFkSher5z2mrPN/3/YH/78IGwMMWV2wxOAtzbLkhuXWTh7cjH2u5sR8xvVeAQgAV0fGCE/ykvj6bF2pvj1r/
5KRY0izG7PAZ ;{id = 57240 (zsk), size = 1024b}
se.      3600    IN      DNSKEY  256 3 5 AwEAAdZgee2x1z9yHDWafJ3oqAnwoU/v/awp0bk6lCnxtoZ6ukq6+0xgY0dusS7qCHF
+LcBRCsAehpQJAWzL7c4xRrs2PT4/z4jZMtfA1EX6hN+s0ZXjxYwR7WdqVje4/Jtn2krpUvE+jjIyegQ+DKFkxbawGJ5pG3EgU3B0MSvEonMT ;{id =
12973 (zsk), size = 1024b}
se.      3600    IN      DNSKEY  257 3 5 AwEAAbaxTum9L7z1DmPiXPk0QZ2/qUM3to210Caey/ycZuvQ8Mh/
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YeyGUvq9Dan9PsvkwQejKN/7J0yCr2nYXfwGGCHkcBkkagv79EaRlZigUCp8= ;{id = 39547 (ksk), size = 2048b}
se.      3600    IN      DNSKEY  257 3 5 AwEAAb6IEZ2ETrgngbj0NAC10b4dRs/jD0MYPcMXRzQlo/eqo5AHXvqPaav+rgA3q
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VIn5stGU1W4juyDqF0= ;{id = 7649 (ksk), size = 2048b}
se.      3600    IN      RRSIG   DNSKEY  5 1 3600 20110409220345 20110403210548 7649 se. U8IIYExbbCcxYeTfQQGB/
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7649}
se.      3600    IN      RRSIG   DNSKEY  5 1 3600 20110410160345 20110404150548 39547 se.
E3uA3bUS0hBNTWoqARj6fSrFdxvaGDjsQRipT5Em+HUX4N09TR2/02tweeA8QKII3bKBRfZ1r56b1LK9nq0eLv3UhPgEbwHwmdpC9fHbRi9FCX
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JsMeKLIbgEvQIsu6zpaZvc1l8ng== ;{id = 39547}
```

Response:Query = 120:1

Average amplification factor: 30x

DNSCurve

DNSCurve: Rationale

DNSSEC: 15 years in the making

DNSSEC: does not solve all the security issues

mathematician
cryptologist
programmer

DNSSCurve: History



DNSCurve: Proposed in 2008

DNSCurve: Objectives

Confidentiality



All DNS payload data is encrypted

IP, UDP, TCP headers are plaintext

Encryption Methods

Does **not** use standard **RSA**

Does use ECC*

* Elliptic-Curve Cryptography

Specifically: Curve25519

Curve25519 is Open/Free

Curve25519 is fast enough for real time
encryption

RSA-1024 requires $\sim 2^{80}$ operations to break

How large computationally speaking is 2^{80} ?



x 1 year == 2^{69}

2048

x 1 year == 2^{80}

2003: 1024-bit RSA deemed breakable

2003: RSA Labs recommends 2048-bit RSA
for the remainder of the decade

2005: NSA recommends ECC for all public-key cryptography and withdrawing previous recommendations of RSA.

2007: NIST recommends 2048-bit RSA

2010: US gov. recommends 2048-bit RSA

Curve25519 == 3000 bit RSA

ECC-256 requires 2^{128} operations

ECC-256: no attack degradation on 25 years



Confidentiality?



Integrity...



All DNS Queries and Responses are
authenticated, cryptographically

Authenticity is guaranteed as well as
non-repudiation

Uses nonces* in all communication to prevent replay attacks

* one time use number

Backwards compatible with regular
DNS services



Integrity?



Availability



Server: several authoritative servers

Client: Non-authenticated/Rogue DNS
servers are rejected

Client: some amplification is produced



Availability



Protocol Specification

Two data formats of are used

Streamlined Format



Custom & Efficient



Query Format

```

                                1 1 1 1 1 1
      0  1  2  3  4  5  6  7  8  9  0  1  2  3  4  5
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
| Q| 6| f| n| v| W| j| 8|
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|                                CLIENT PUBLIC KEY                                |
|                                                                                       |
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|                                CLIENT NONCE                                |
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
/                                                                                       /
/                                CRYPTOGRAPHIC BOX                                /
/                                                                                       /
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
```



Response Format

```

                                1 1 1 1 1 1
      0  1  2  3  4  5  6  7  8  9  0  1  2  3  4  5
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
| R| 6| f| n| v| W| J| 8|
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|                               CLIENT NONCE                               |
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|                               SERVER NONCE                               |
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
/                                                                           /
/                               CRYPTOGRAPHIC BOX                          /
/                                                                           /
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
```



TXT-Format



DNSCurve: Key Usage



No new RRs are introduced



DNSSEC: RRs provide the keys



DNSSCurve: No new RRs



How do we get the name server keys?



“If it were not for key management,
Cryptography would be easy!”



“note -- i think dnssec is terribly ugly
but i have come to terms with that and
am pushing forward with it because i
want what it can do for the world.”

Paul Vixie, BIND/DNSSEC author/architect



Keys are obtained from existing RRs!



Server pub keys are hidden in NS records



E.g. example.org



Regular DNS: ns1.example.org

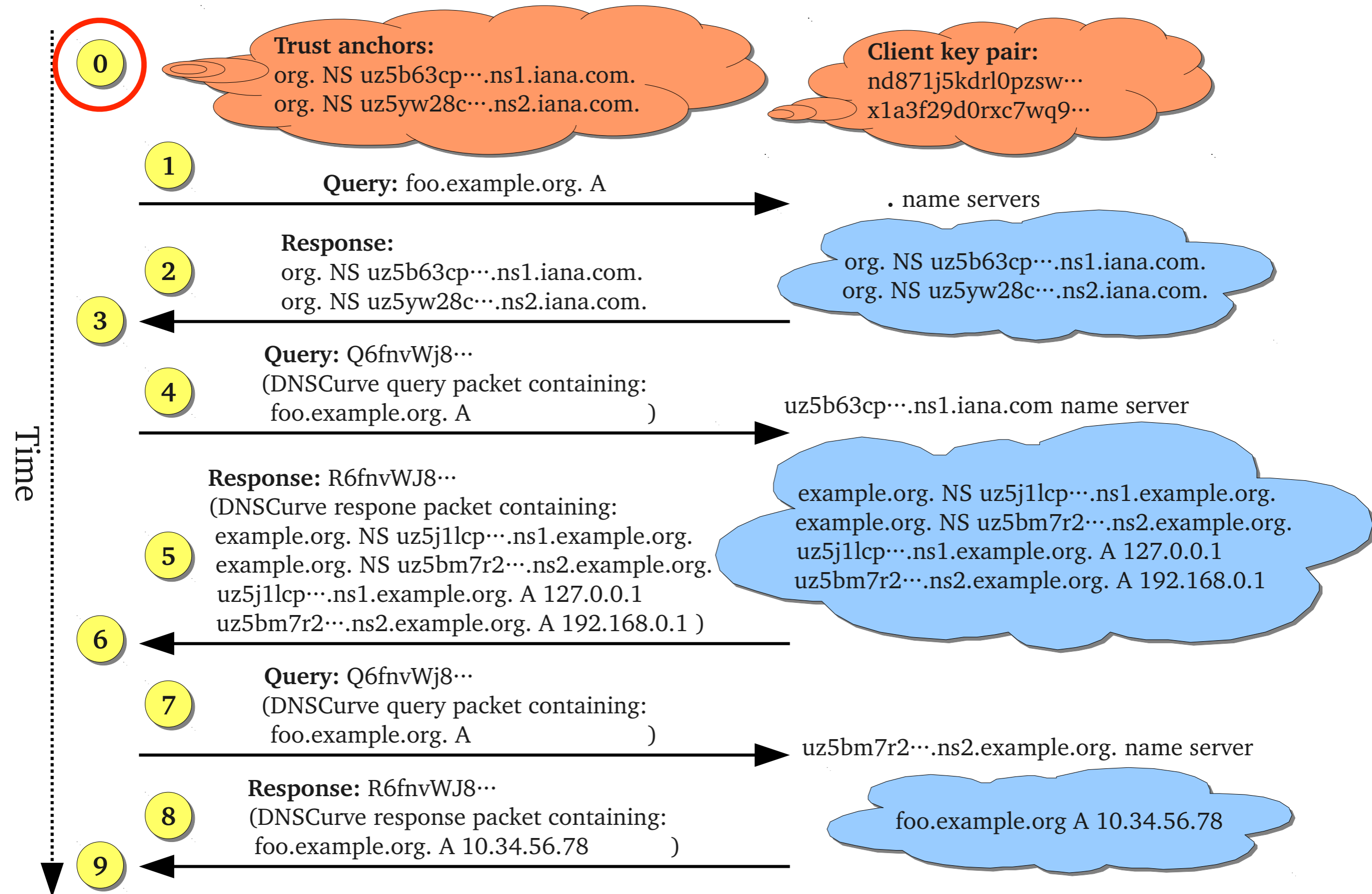


DNScurve:

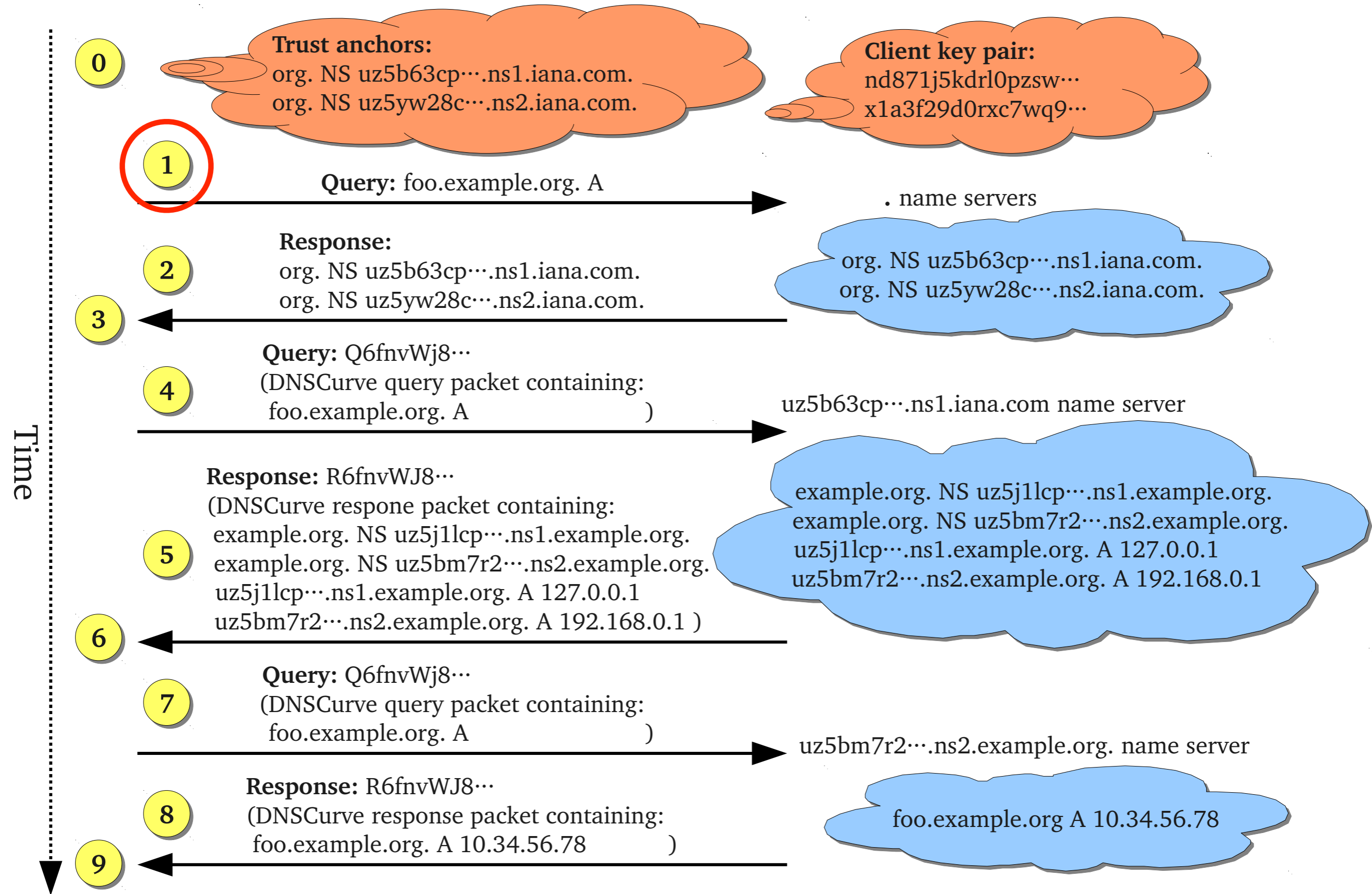
uz5dkhm9g380kyx9slmktyvmb1h0ck7whwzc5uqvl8f1cwfp8zl3ub.ns1.example.org



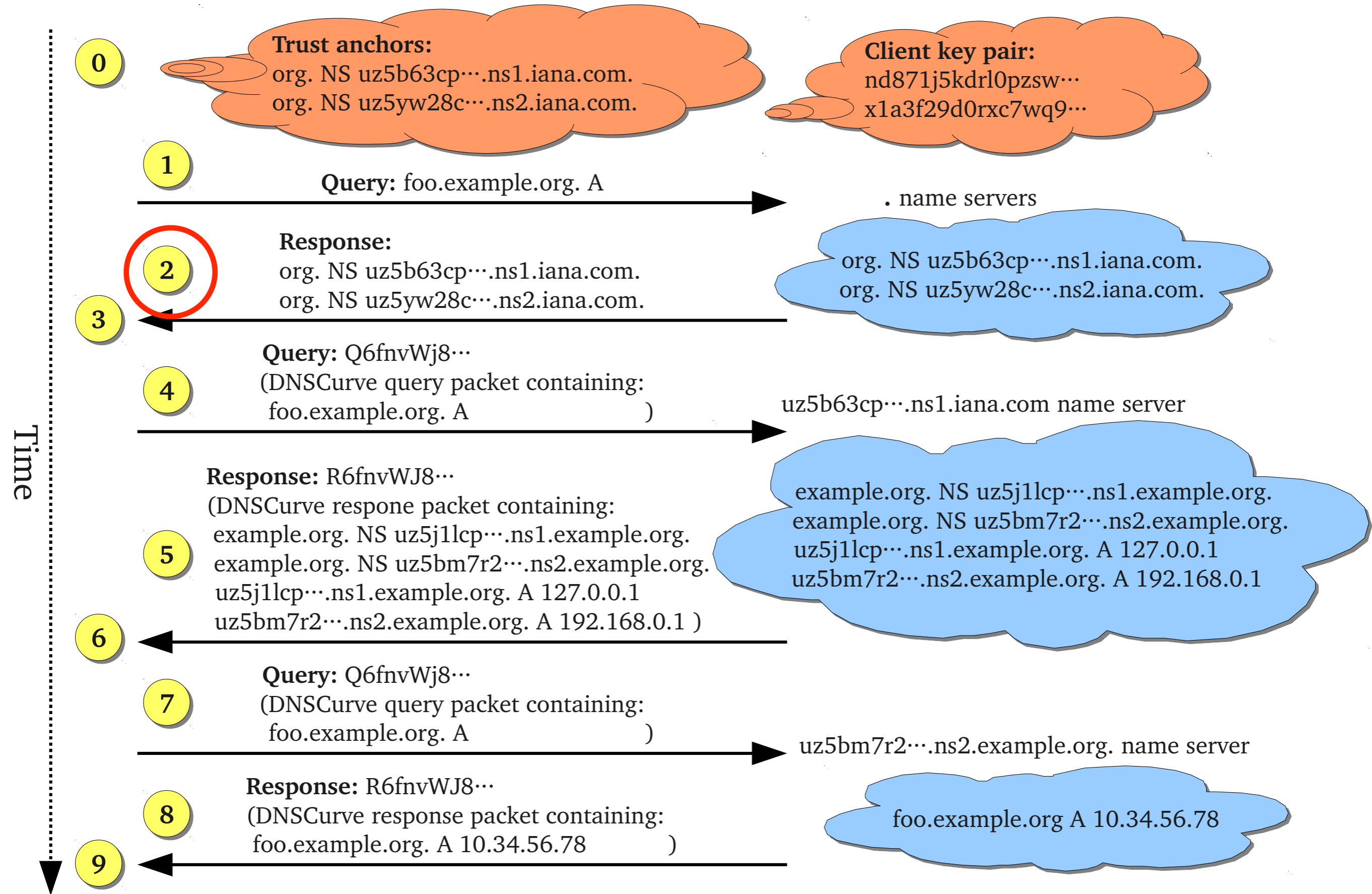
DNSTraversal: Traversal



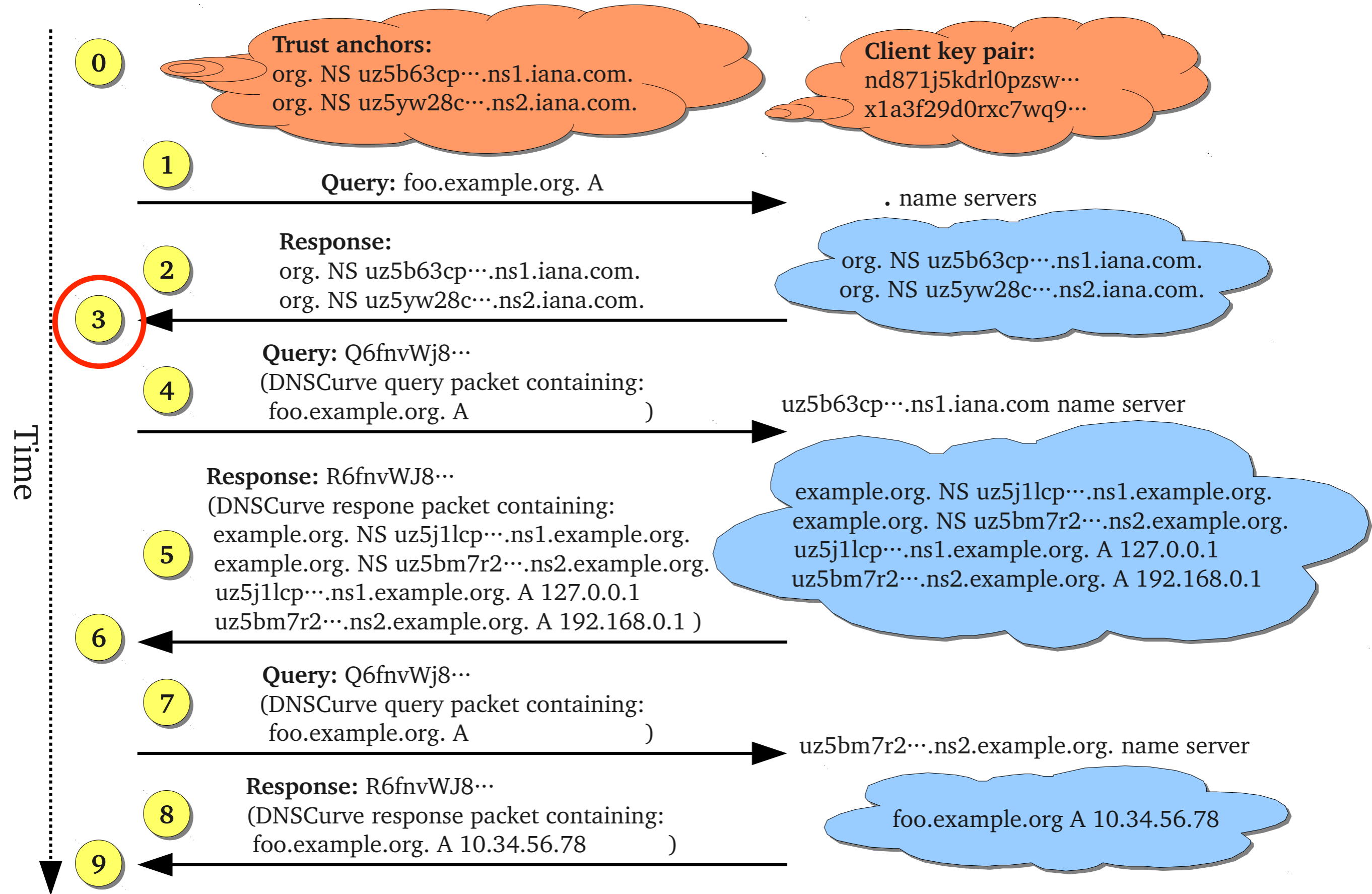
Graphical representation of a DNSCurve traversal



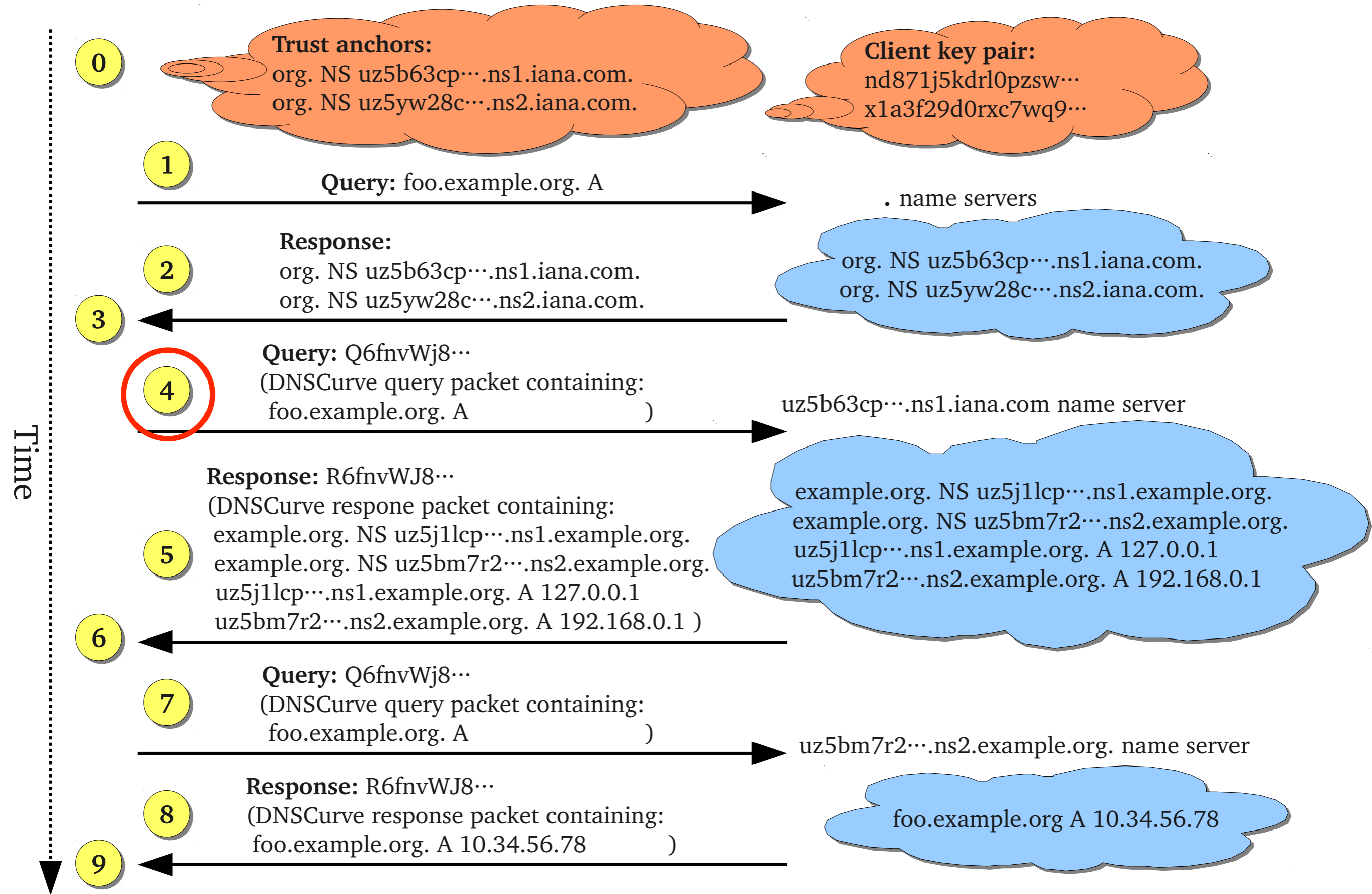
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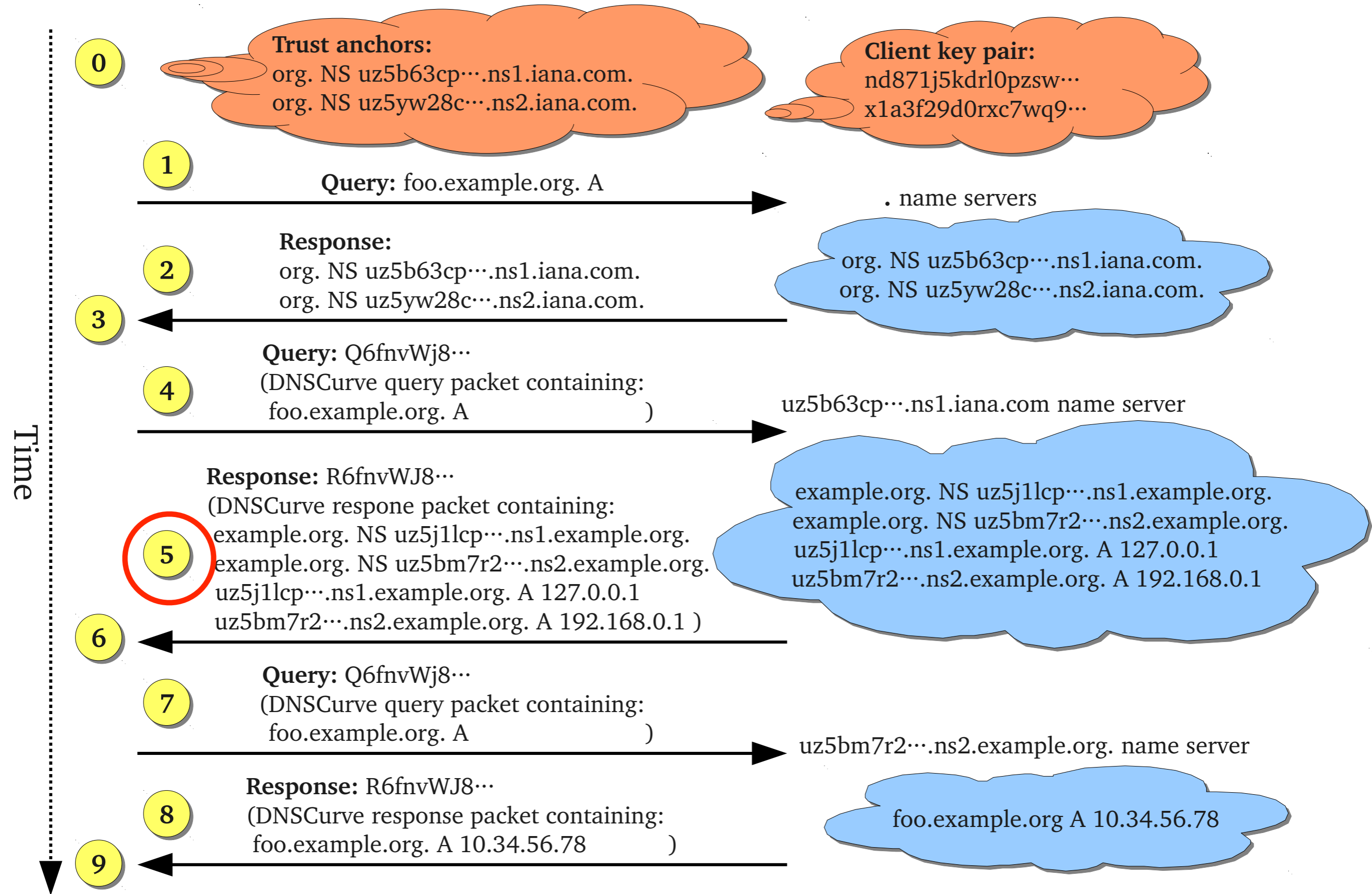
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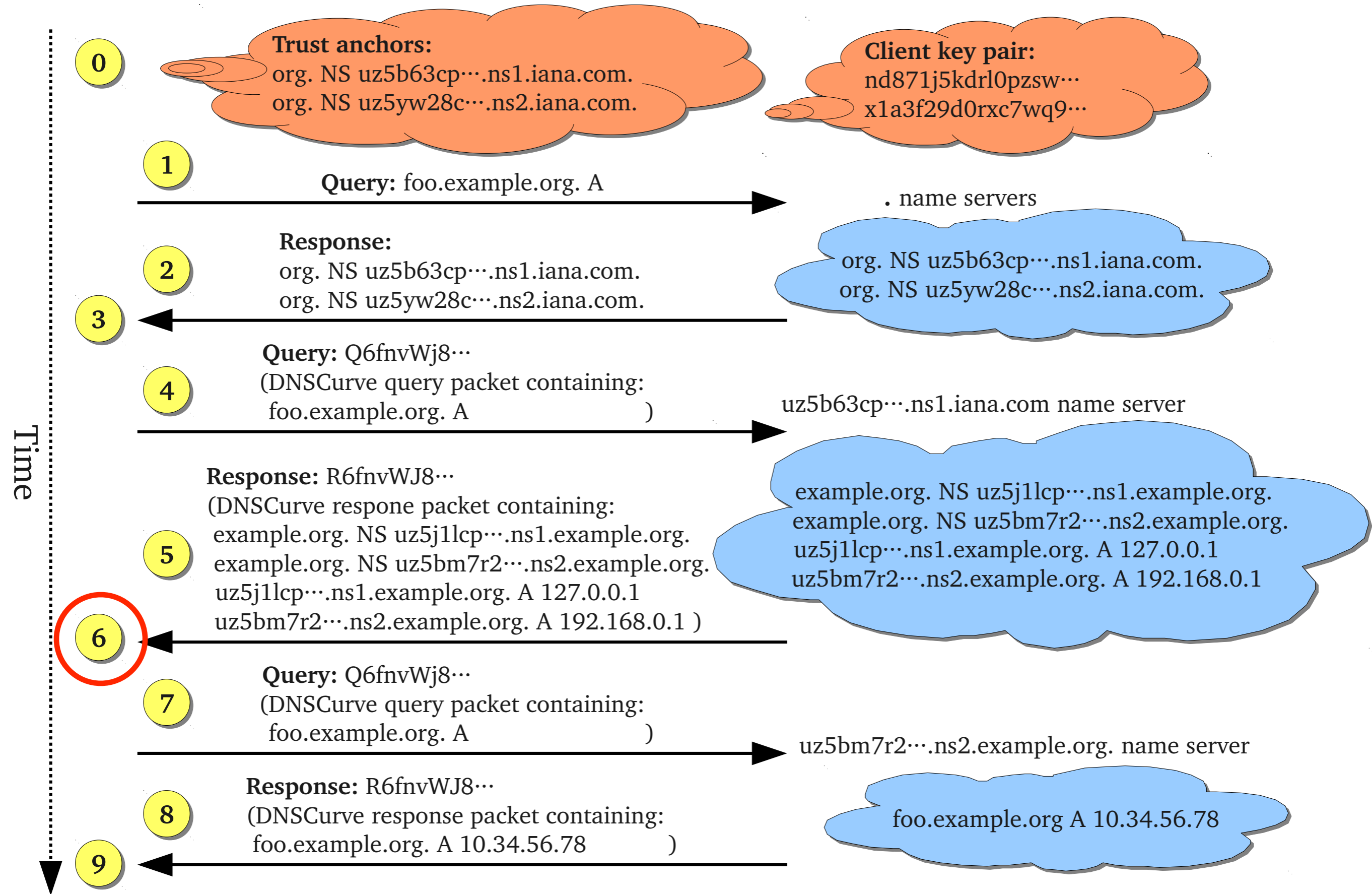
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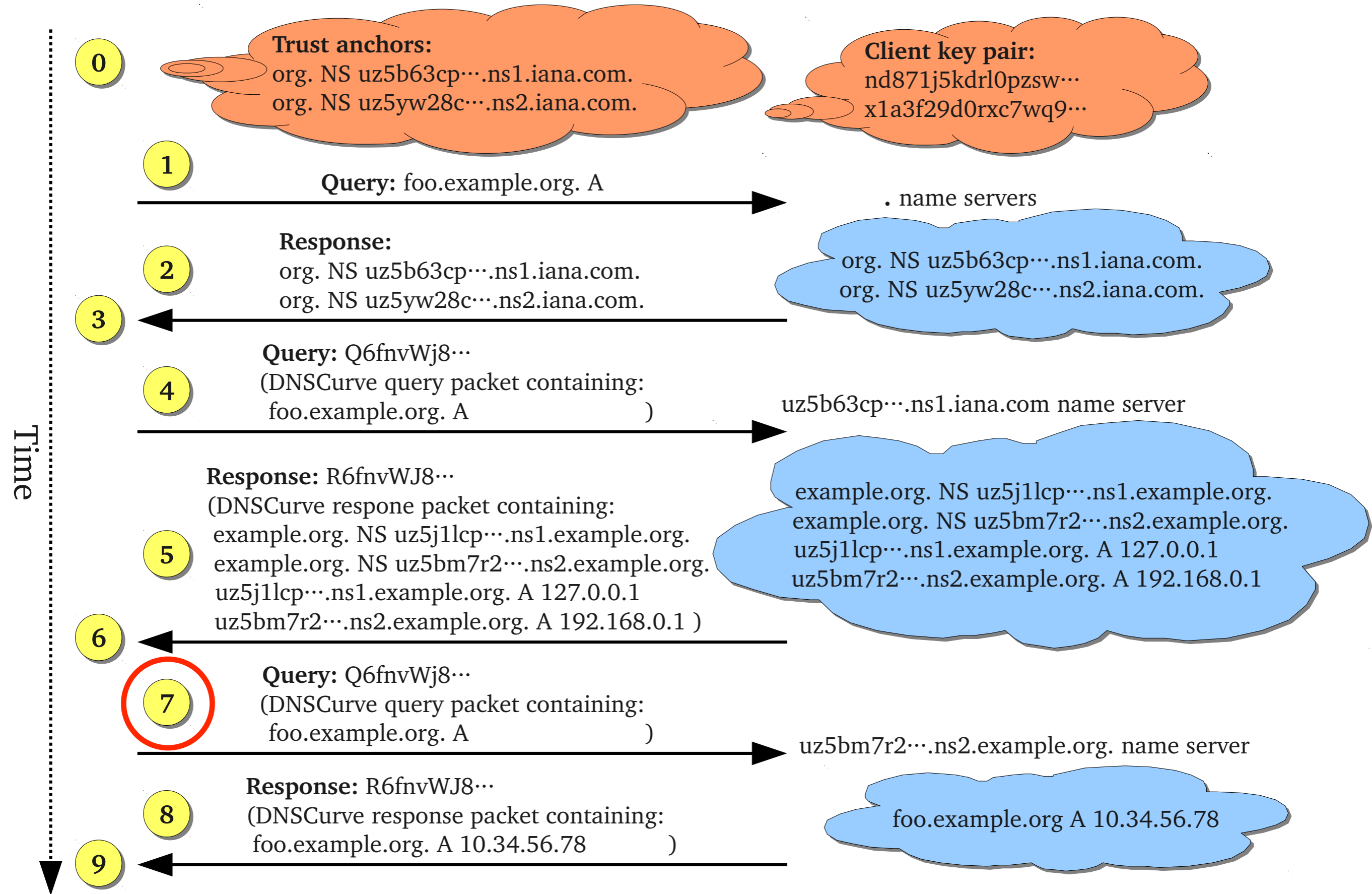
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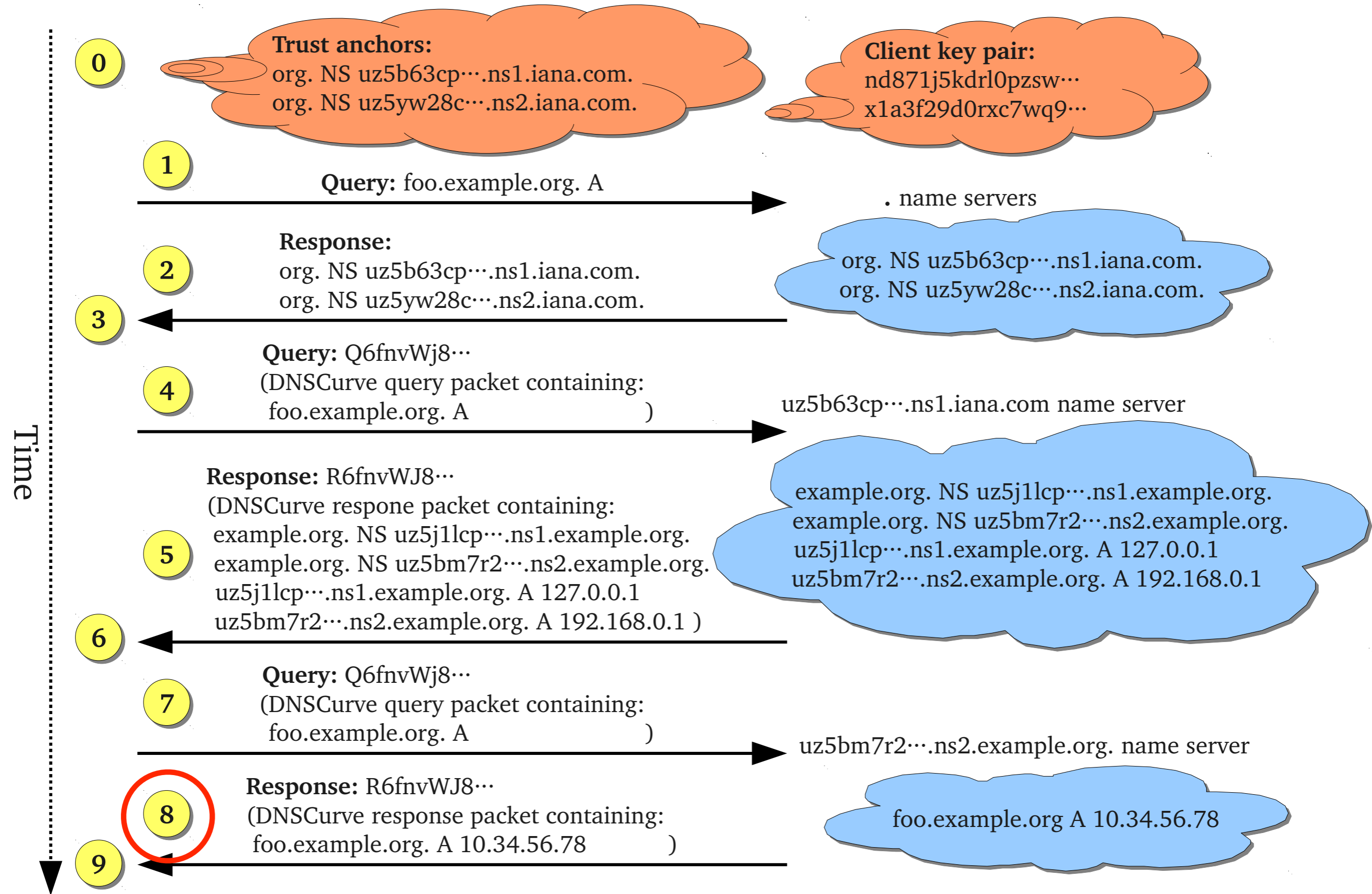
Graphical representation of a DNSCurve traversal



Graphical representation of a DNSCurve traversal

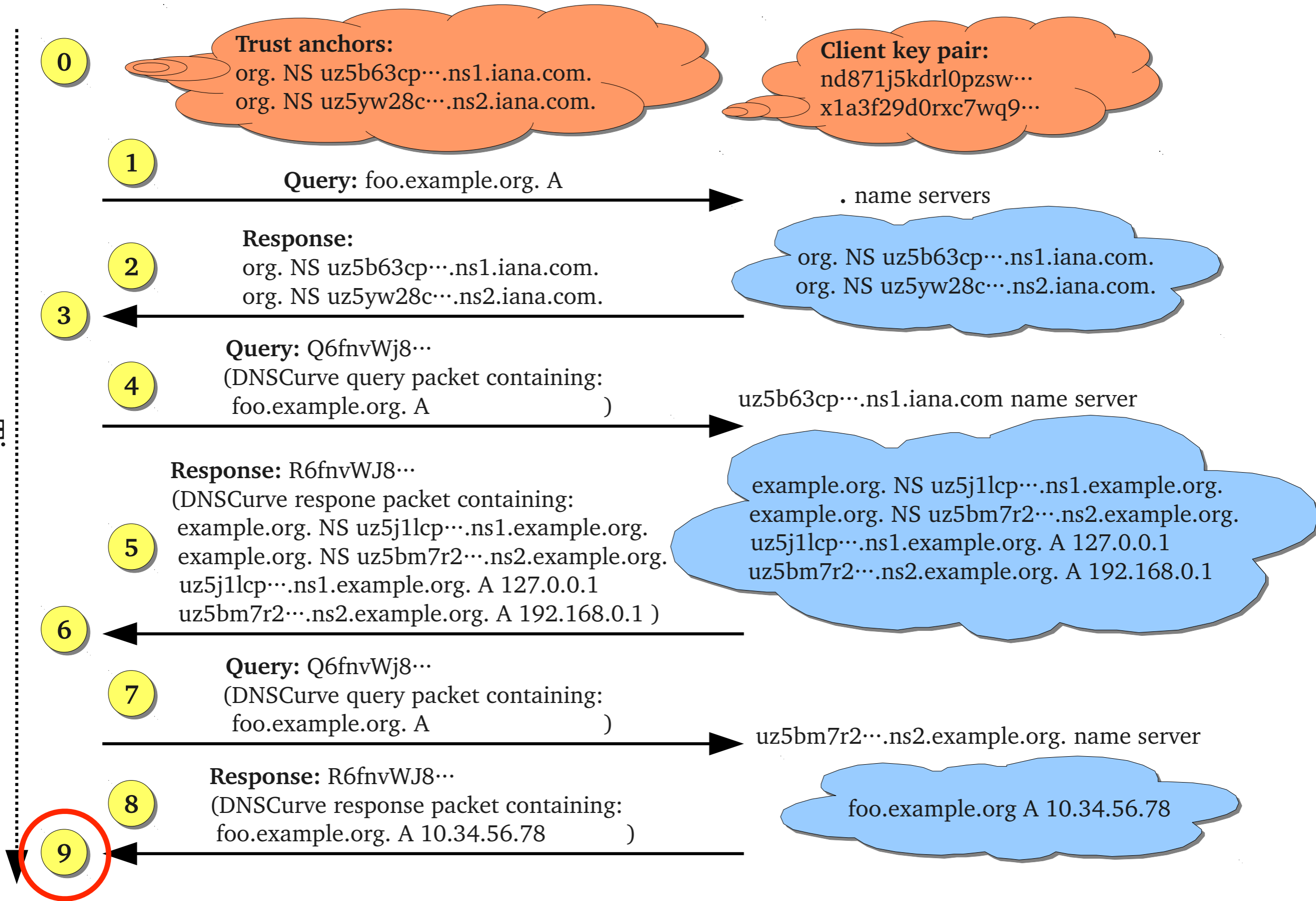


Graphical representation of a DNSCurve traversal



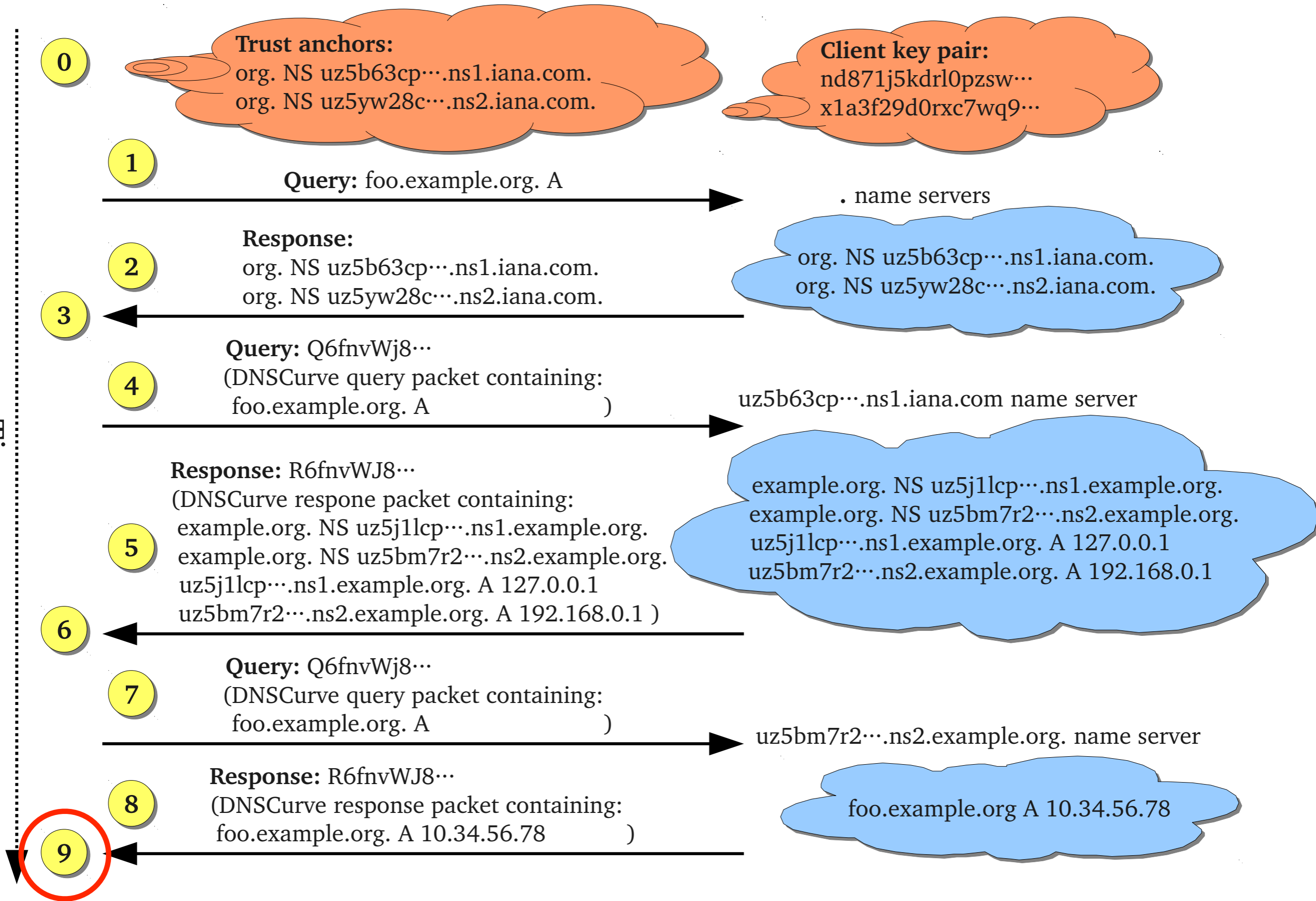
Graphical representation of a DNSCurve traversal

Time



Graphical representation of a DNSCurve traversal

Time



Graphical representation of a DNSCurve traversal

Conclusions...

DNSSEC is big and messy

DNSSEC solves some security problems

but creates some significant others

DNSSEC has problems with the last mile

Given HTTPS, what exactly does DNSSEC offer?

DNSTCurve is less messy

DNSSCurve solves more problems than
DNSSEC

DNSCurve is a more general solution

Both DNSSEC & DNSCurve need to be
tested and tried locally.